Looking good, behaving well

Michael Westergaard

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University of Aarhus

Example (1/4)



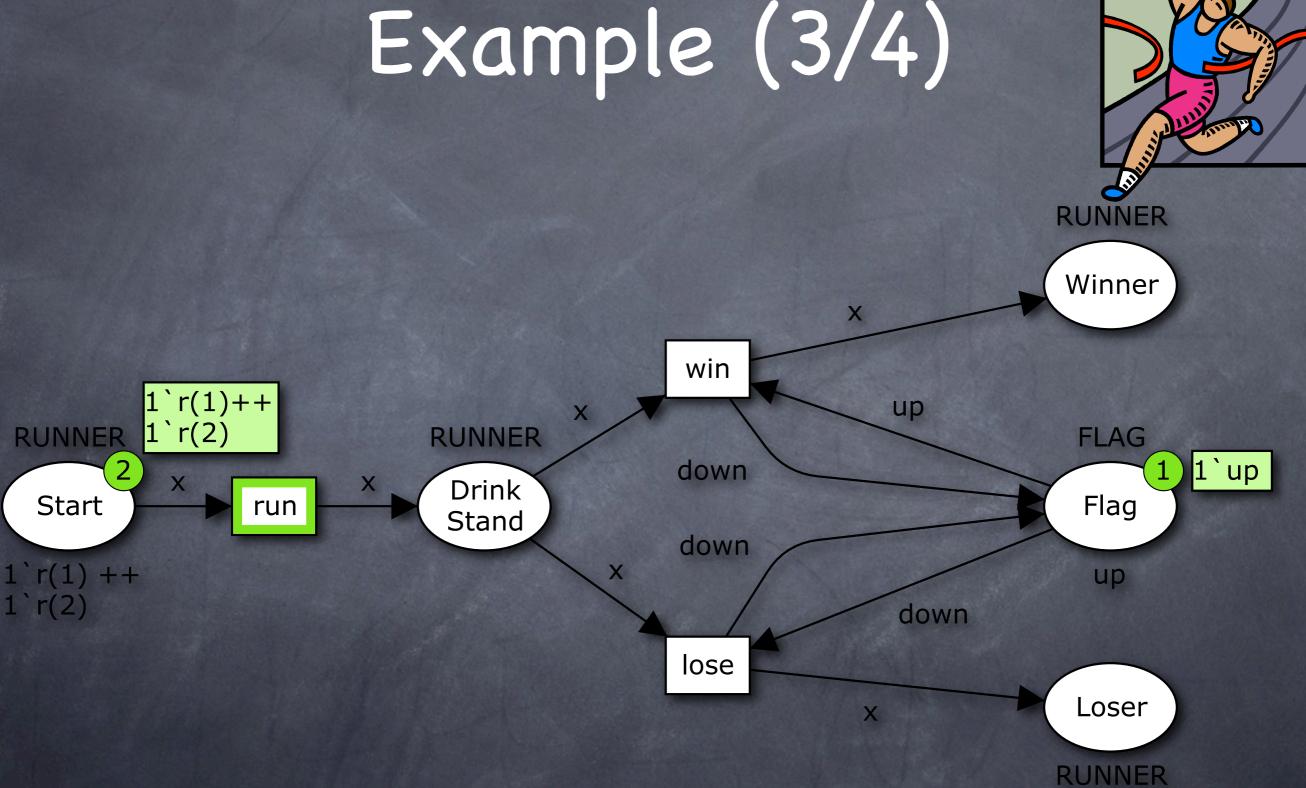
- 2 runners in a race, halfway through the race is a stand with water
- Either
 - orun: a runner runs to the drink stand,
 - win: a runner wins the race, or
 - o lose: a runner loses the race

Example (2/4)

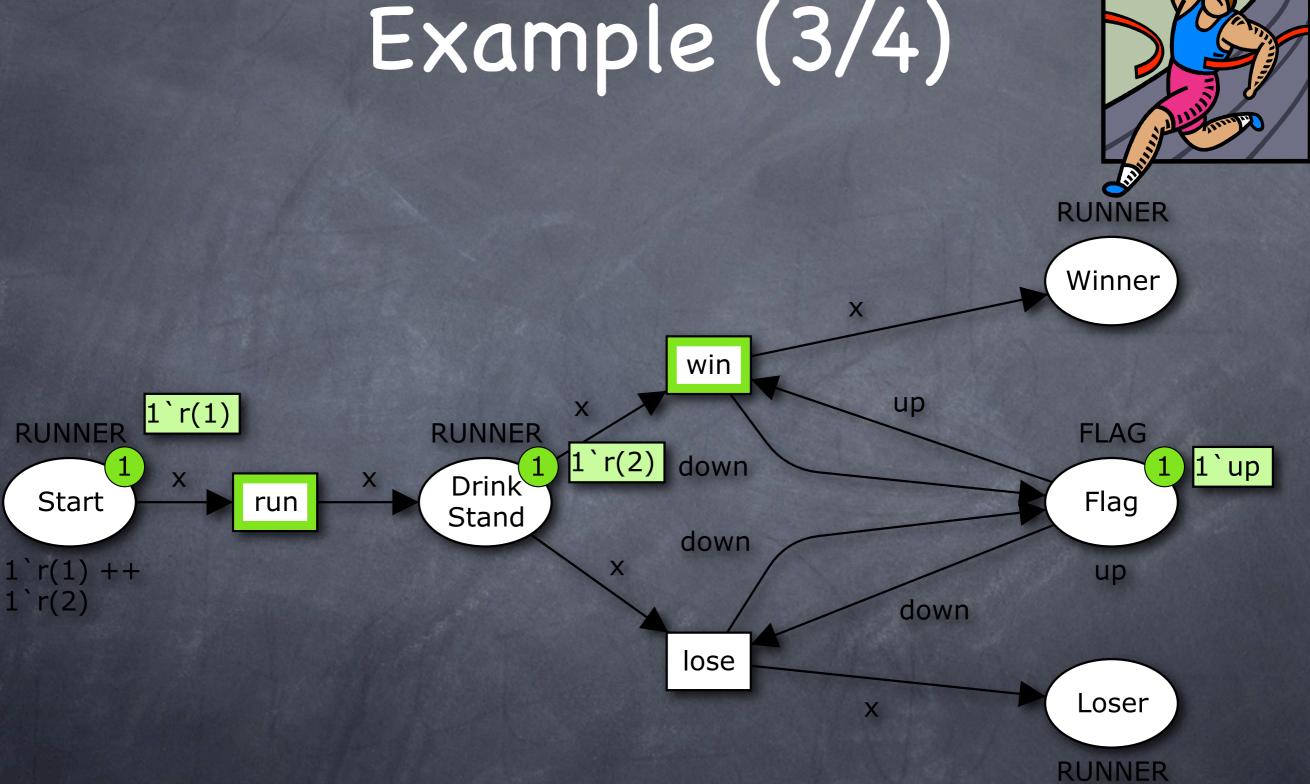


- Only one runner can win the race
- In the beginning neither of the runners have finished any laps
- We can model this using Timed Automata or Coloured Petri nets

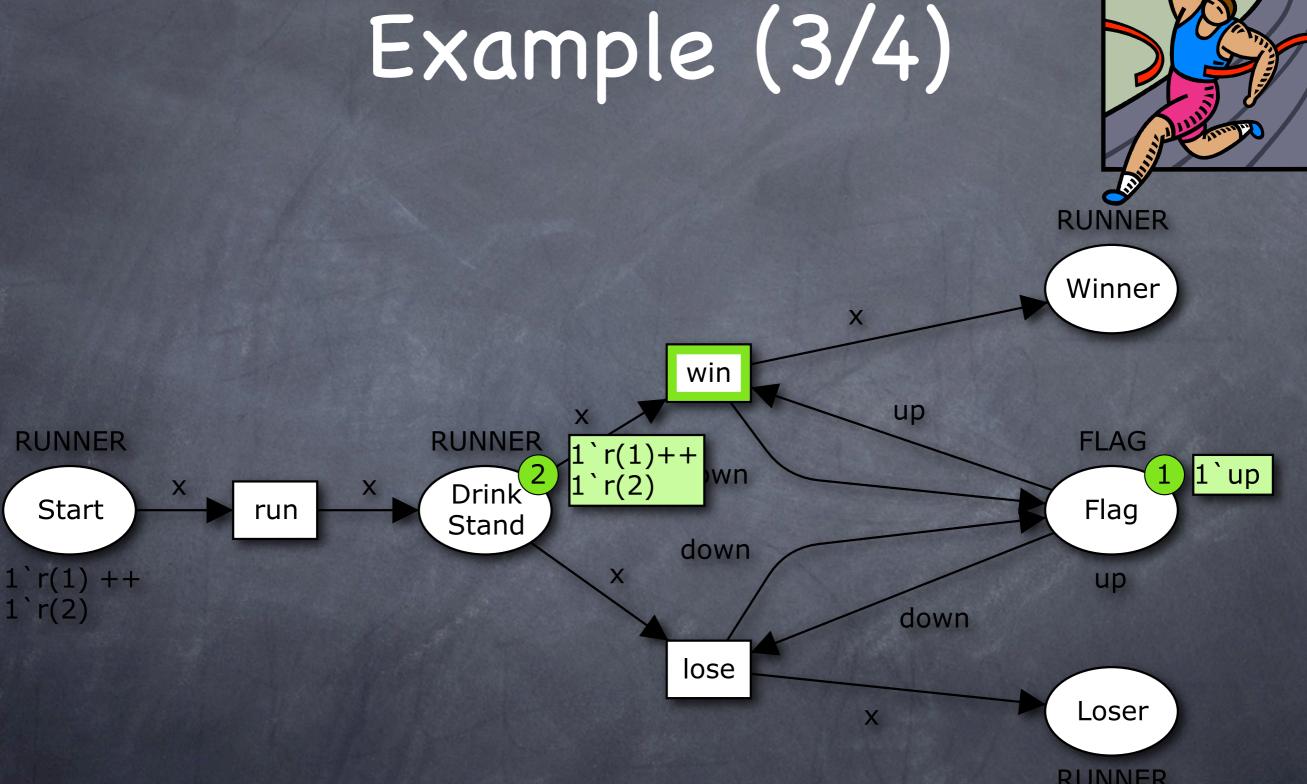
Example (3/4)



Example (3/4)



Example (3/4)



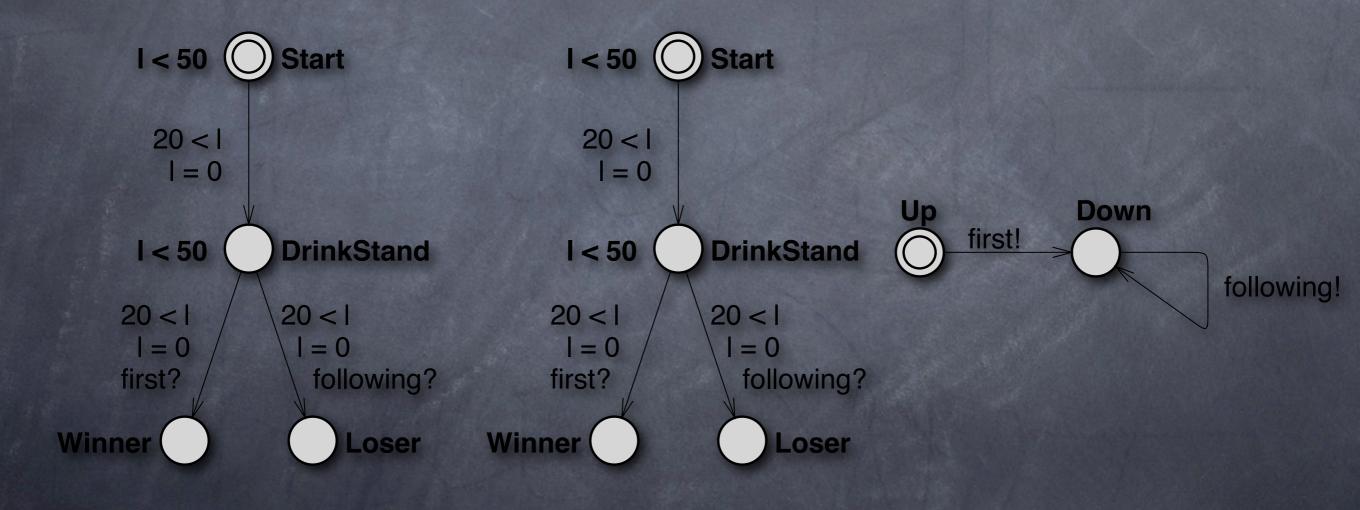
RUNNER

Example (3/4) RUNNER 1`r(2) Winner X win up RUNNER RUNNER **FLAG** 1`r(1) 1`down down Drink Start Flag run Stand down 1 r(1) ++ X up down lose Loser X RUNNER

Example (3/4) RUNNER 1`r(2) Winner X win up RUNNER RUNNER **FLAG** 1`down down Drink Start Flag run Stand down 1 r(1) ++ X up down 1`r(1) lose Loser X RUNNER

Example (4/4)





We want the model to

- look good (even to people not familiar with the modeling formalism)
- behave well (e.g. ensure only one runner can win the race)

Outline

- The BRITNeY animation tool
- A state space tool
 - Memory-efficient state storage using the sweep-line method
 - Memory-efficient state storage using hash compaction and backtracking

Looking Good

The BRITNeY animation tool



Domain expert



Domain expert

Figure 2 shows the approach taken to use CPN models to develop a prototype of the interoperability protocol. A CPN model (lower left of Fig. 2) has been developed by modelling the natural language protocol specification [22] (lower right) of the interoperability protocol. The modelling activity transforms the natural language specification into a formal executable specification represented by the CPN model. The CPN model captures the network architecture depicted in Fig. 1 and the protocol mechanisms of the interoperability protocol, e.g., the periodic transmission of advertisements, the dynamic updates of the DNS database, and traffic flows between hosts in the core network and nodes in the ad-hoc

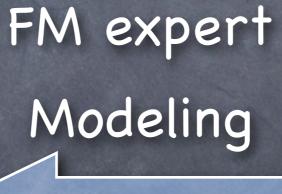




FM expert

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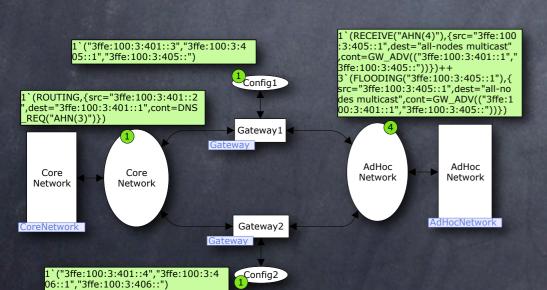


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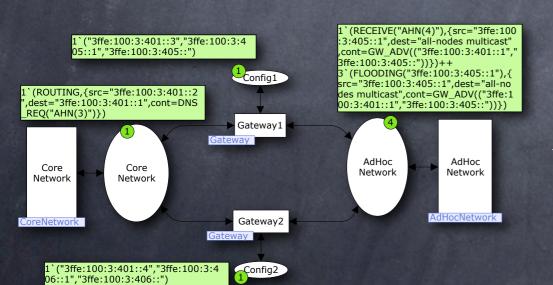
FM expert
Modeling

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Formal model



Domain expert



Modeling

Formal model

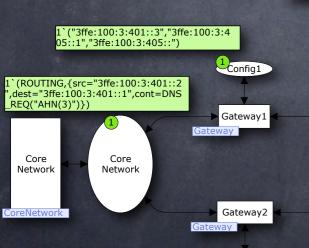
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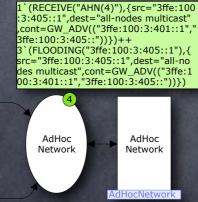
Validate

Domain expert

Figure 2 shows the approach taken to use CPN models to develop a



1`("3ffe:100:3:401::4","3ffe:100:3:4 06::1","3ffe:100:3:406::")



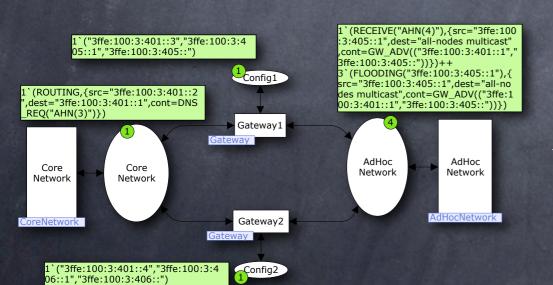
Modeling

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Formal model



Domain expert



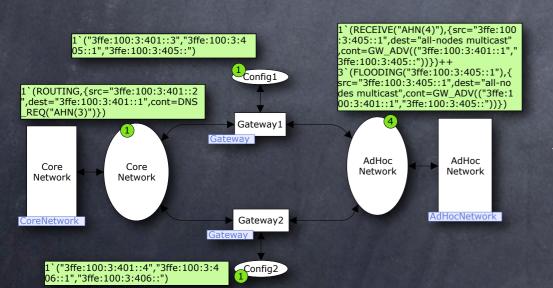
Modeling

Formal model

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Domain expert



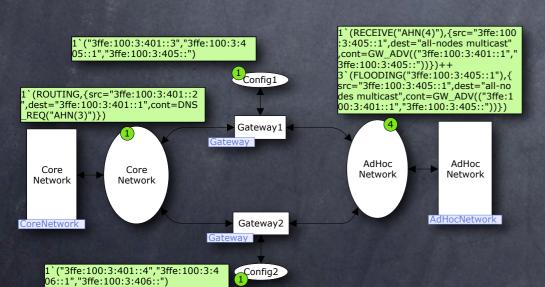
Modeling

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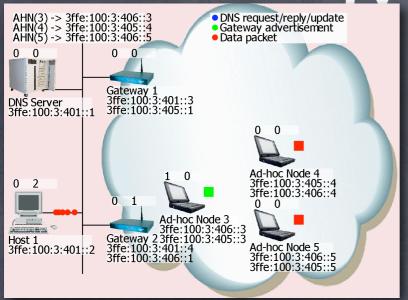




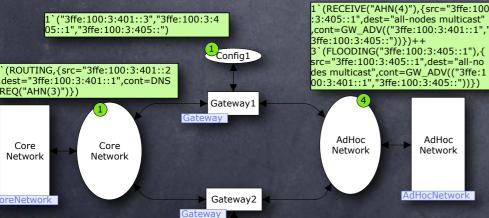
FM expert
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Formal model







cont=GW_ADV(("3ffe:100:3:401::1"," ffe:100:3:405::"))})++
'(FLOODING("3ffe:100:3:405::1"),{ rc="3ffe:100:3:405::1",dest="all-no les multicast",cont=GW_ADV(("3ffe:1 0:3:401::1","3ffe:100:3:405::"))})

FM expert Modeling

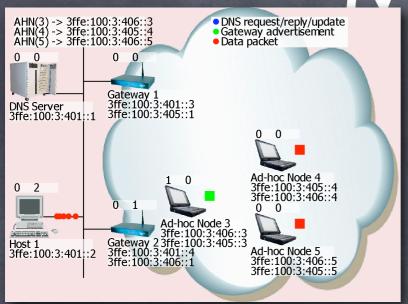


Domain expert

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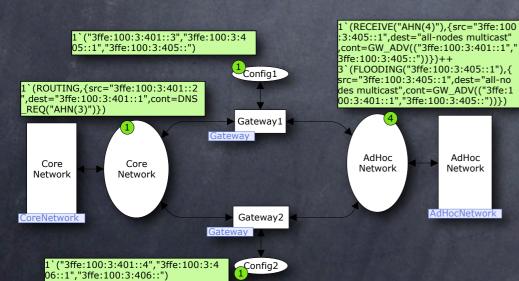
Formal model

1`("3ffe:100:3:401::4","3ffe:100:3:4 06::1","3ffe:100:3:406::")





Domain expert Animation



Modeling

Formal model

AdHoc

Network

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Domain expert



1`("3ffe:100:3:401::4","3ffe:100:3:4 06::1","3ffe:100:3:406::")

0 2

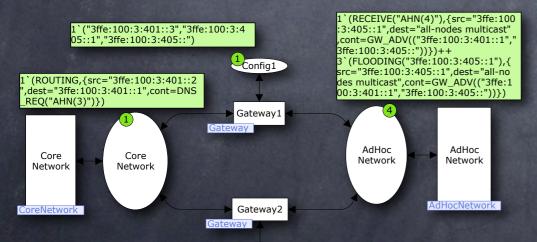
Host 1 3ffe:100:3:401::2 Gateway 1 3ffe:100:3:401::3 3ffe:100:3:405::1

Ad-hoc Node 3 3ffe:100:3:406::3 Gateway 2 3ffe:100:3:405::3 3ffe:100:3:401::4 3ffe:100:3:406::1

Animation

3ffe:100:3:405::4 3ffe:100:3:406::4

Ad-hoc Node 5 3ffe:100:3:406::5 3ffe:100:3:405::5

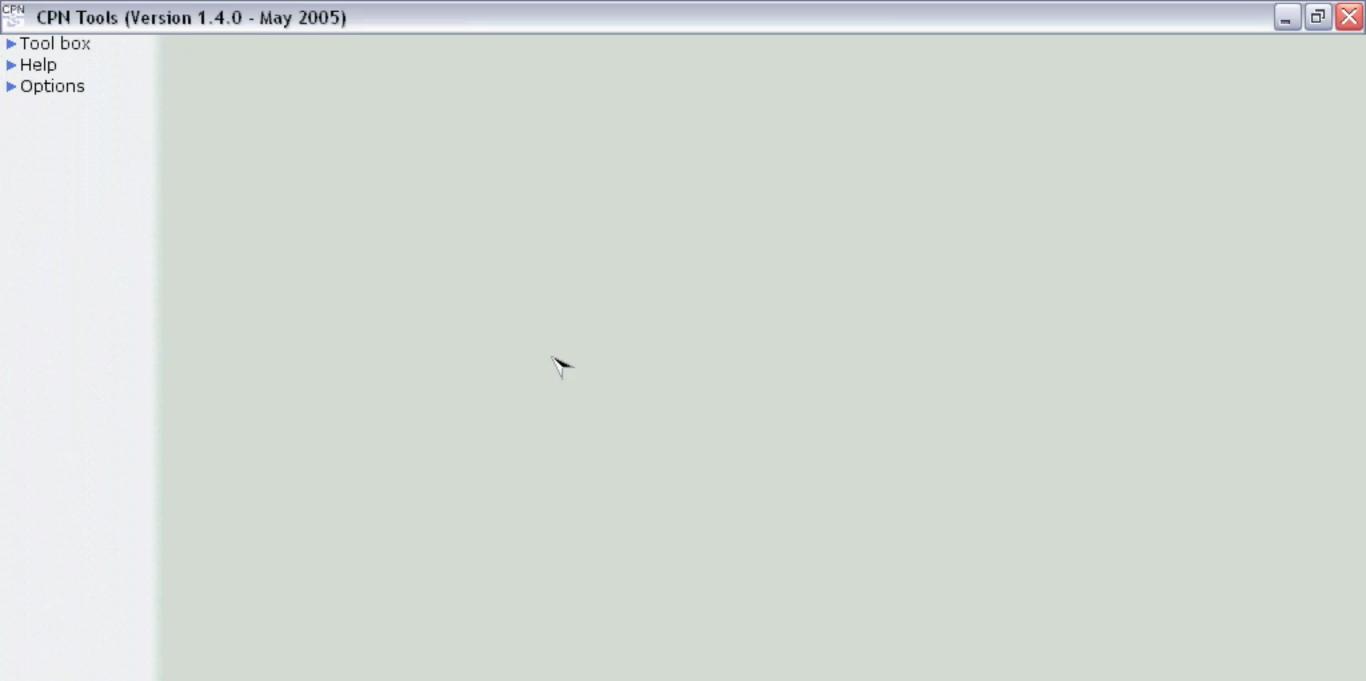


Modeling

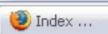
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Formal model

Good-looking Runners Executive Edition







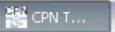














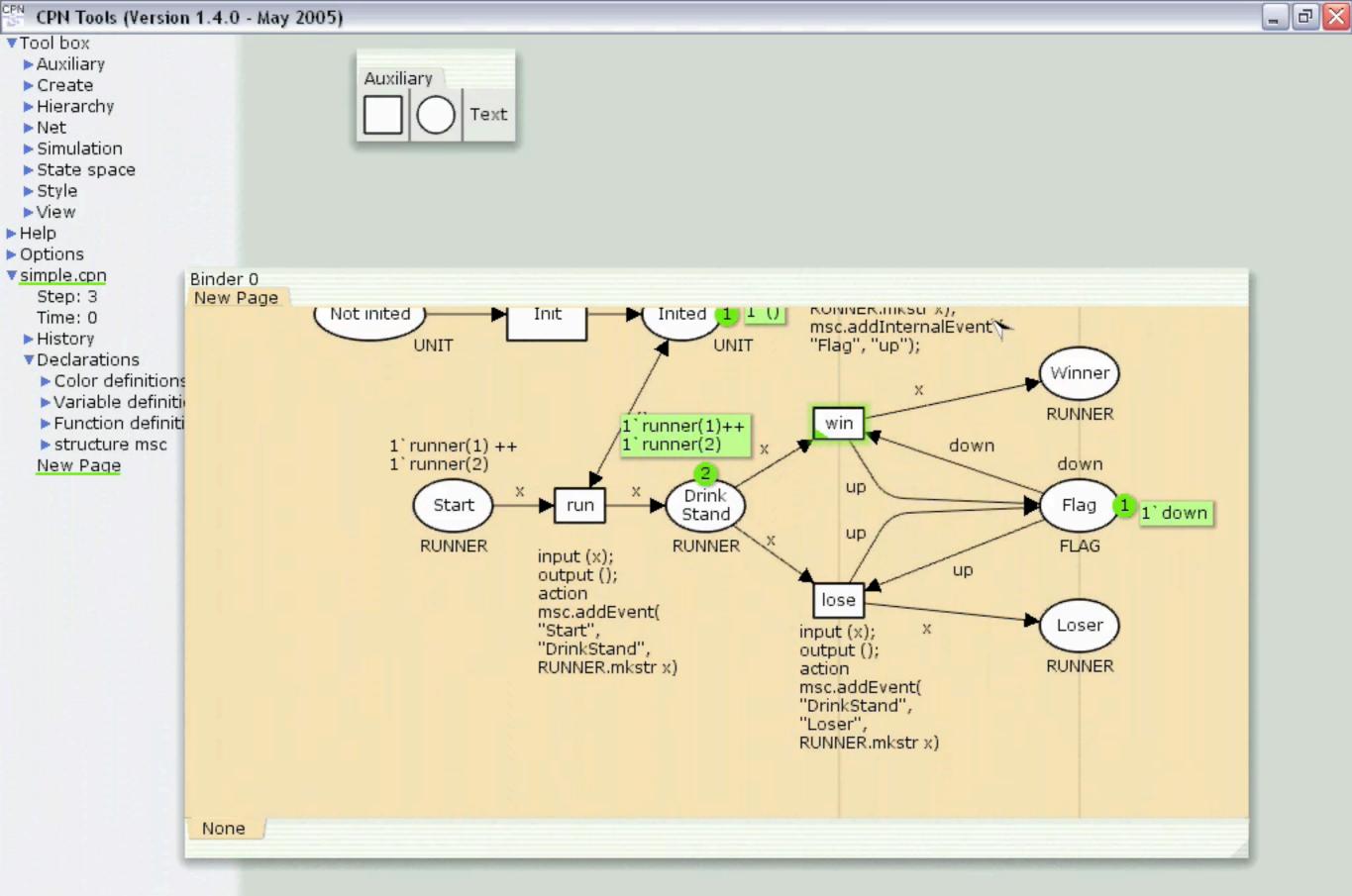
DA







Good-looking Runners Engineer Edition





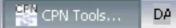










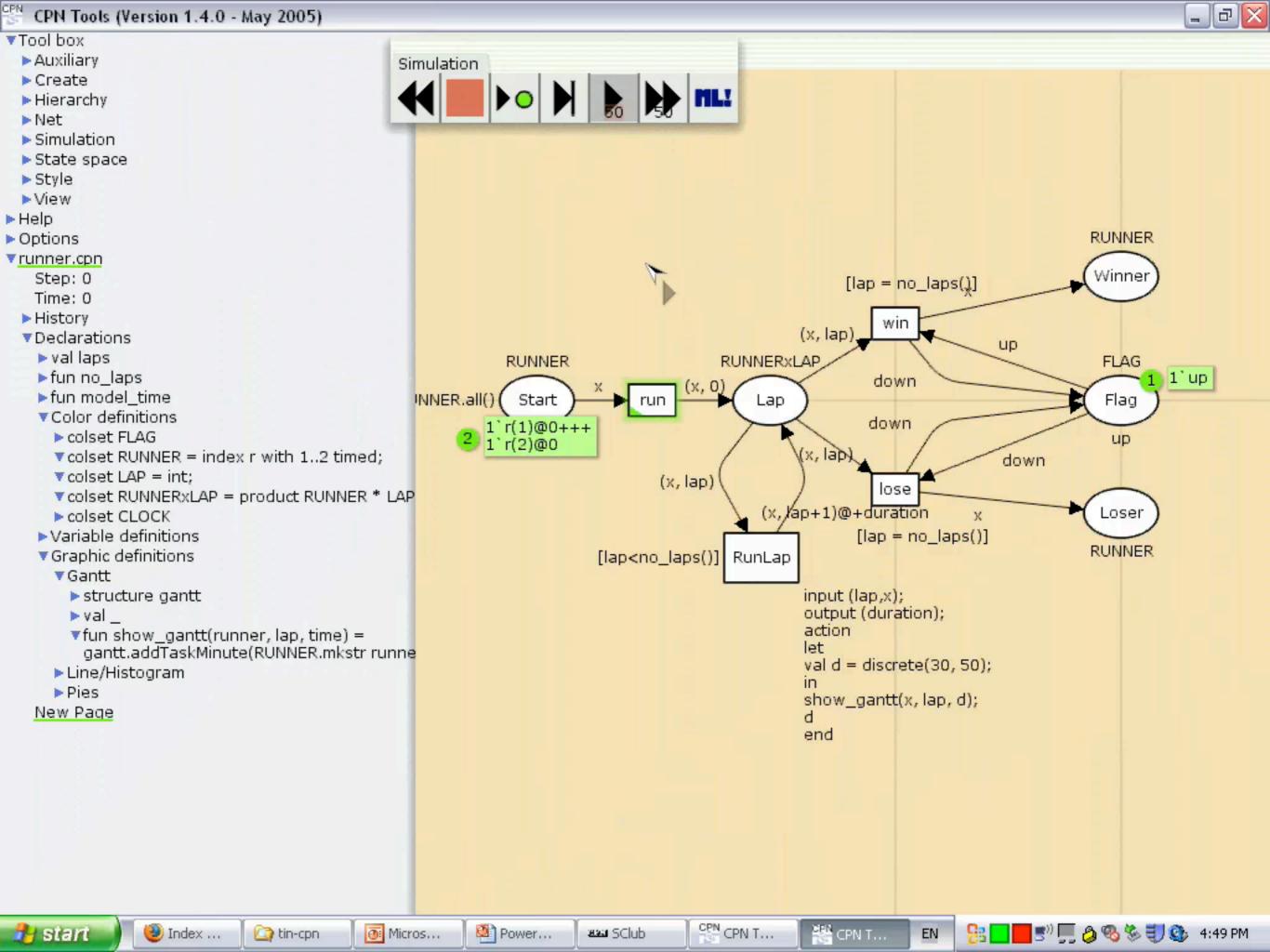




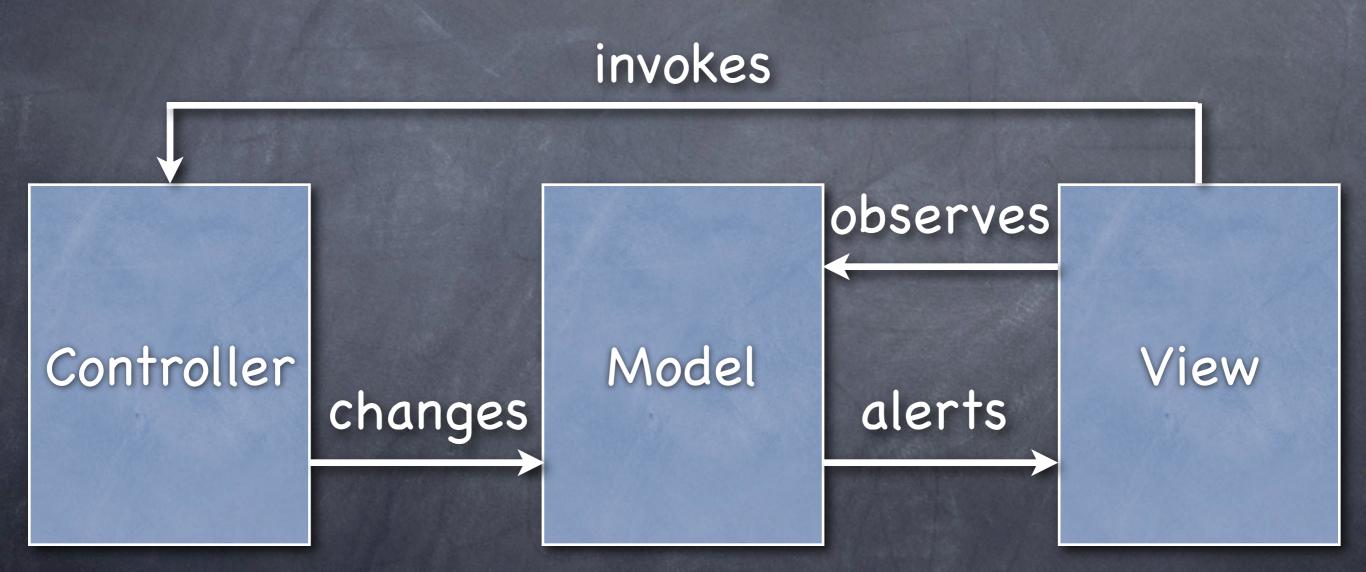




Good-looking Runners Performance Analyst Edition



BRITNeY animation (1/3)



BRITNeY animation (2/3)

invokes

Controller + Model

observes

alerts

View

Formal executable model

BRITNeY animation

BRITNeY animation (3/3) Animation

plugins

CPN Tools editor

CPN simulator

Animation stubs

XML-RPC client

BRITNeY animation

XML-RPC server

BRITNey animation (3/3) Animation

plugins

Formal executable model tool

Animation stubs

XML-RPC client

BRITNeY animation

XML-RPC server

More Information about BRITNeY animation

- Tool web-page: http://wiki.daimi.au.dk/tincpn
- Screen-cast from CPN Workshop 2005 tutorial: http://www.daimi.au.dk/~mw/local/demo/ BRITNeY animation/
- © Case study: L.M. Kristensen, M. Westergaard, and P.C. Nørgaard: Model-based Prototyping of an Interoperability Protocol forMobile Ad-hoc Networks, IFM 2005

Behaving Well

A state space tool

Reachability: Does any possible state of the system satisfy a given property?

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- Reachability: Does any possible state of the system satisfy a given property?
 - E.g., can more than one runner win?
- Analysis of CP-nets is impossible
- How do we do it anyway? Try all possible states
- Loops? Build reachability graph

State Space (w,l,d) for Runners lose 1 run 1 (s,w,d) (d,w,d) (l,w,d) win 2 win 2 lose 2 run 1 win 1 (r1,r2,flag) (d,d,u) (s,d,u) (w,d,d) $r1,r2 \in \{s,d,w,l\}$ run 2 run 2 run 2 $flag \in \{u,d\}$ run 1 win 1 (s,s,u) (w,s,d) (d,s,u)

Simple Algorithm for State Space Analysis

```
Queue.add(Simulator.get_initial state())
while !Queue.is_empty() do
  s := Queue.remove_first()
  Storage.add(s)
   process(s)
   forall s' in Simulator.get_next(s) do
     if !Storage.contains(s') then
        Queue.add(s')
     endif
  endfor
endwhile
```

Parametrizing the Algorithm

- The algorithm relies on 3 data-structures:
 - Simulator (get_initial_state, get_next)
 - Queue (add, is_empty, remove_first)
 - Storage (add, contains)
- By providing different implementations, we can control which formalism to use (Simulator), how to traverse the state space (Queue – waiting/unprocessed), and how to store data efficiently (Storage – passed/ processed)

Problem: The reachability graph is large, often even infinite

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- Problem: The reachability graph is large, often even infinite
 - Solution 1: Store only some of the graph
 - Solution 2: Store each node more efficiently
- A lot of so-called reduction methods exist and new reduction methods are found out every day

A Condensed Representation (1/2)

- We want to represent the entire state space
- A state of the system is (r1,r2,flag) with r1,r2∈{s,d,w,l} and flag∈{u,d}
- Only some (10) of the syntactically possible states (4·4·2=32) are reachable
- At least ceil(log(32))=5 bits are used to store each state, although ceil(log(10))=4 bits would suffice

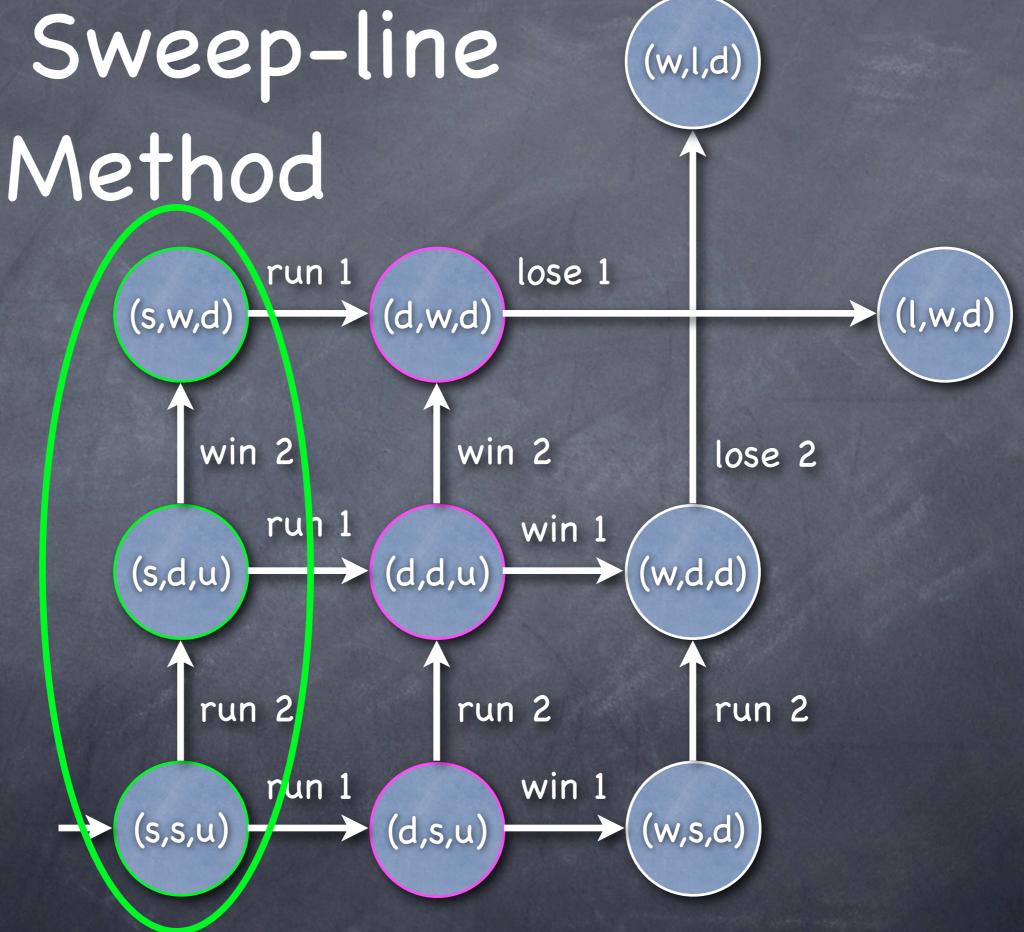
A Condensed Representation (2/2)

- In realistic examples, the number of syntactically possible states is much larger than the number of reachable states, so distinguishing only between reachable states yields a good reduction
- Alas, we first know the number of reachable states, when we have constructed the reachability graph

The Sweep-line (w,l,d) Method lose 1 run 1 (d,w,d) (s,w,d) (l,w,d) win 2 win 2 lose 2 run 1 win 1 (s,d,u) (d,d,u) (w,d,d) run 2 run 2 run 2 win 1 run 1 (d,s,u) (s,s,u) (w,s,d)

The Sweep-line

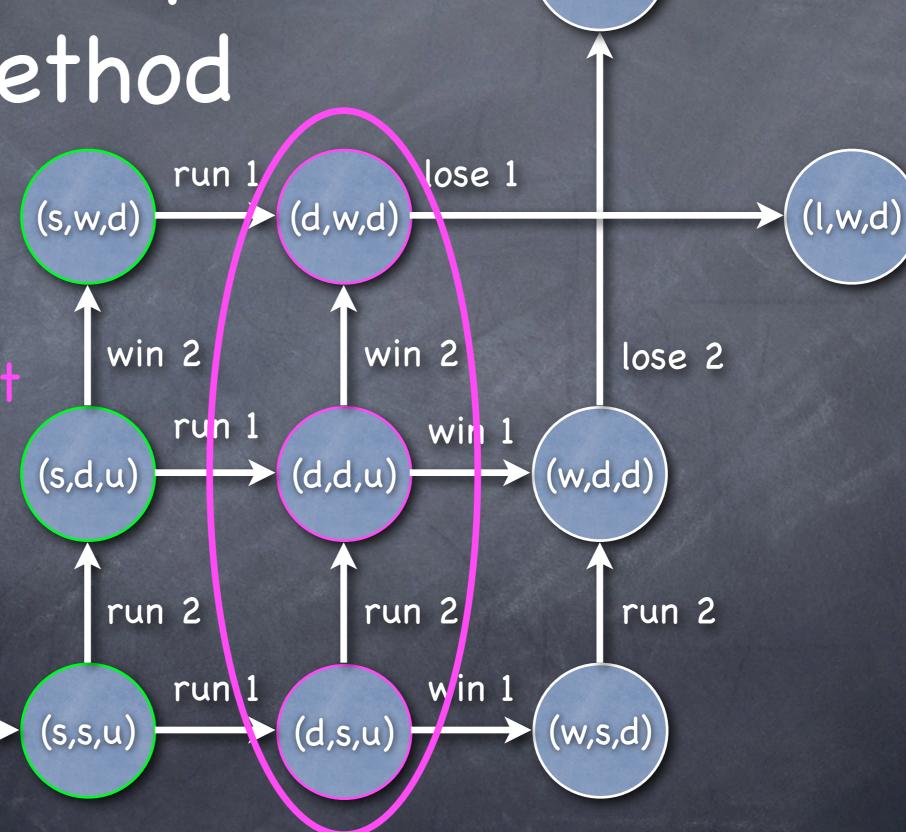
Already processed



The Sweep-line Method

Already processed

Discovered but not yet processed



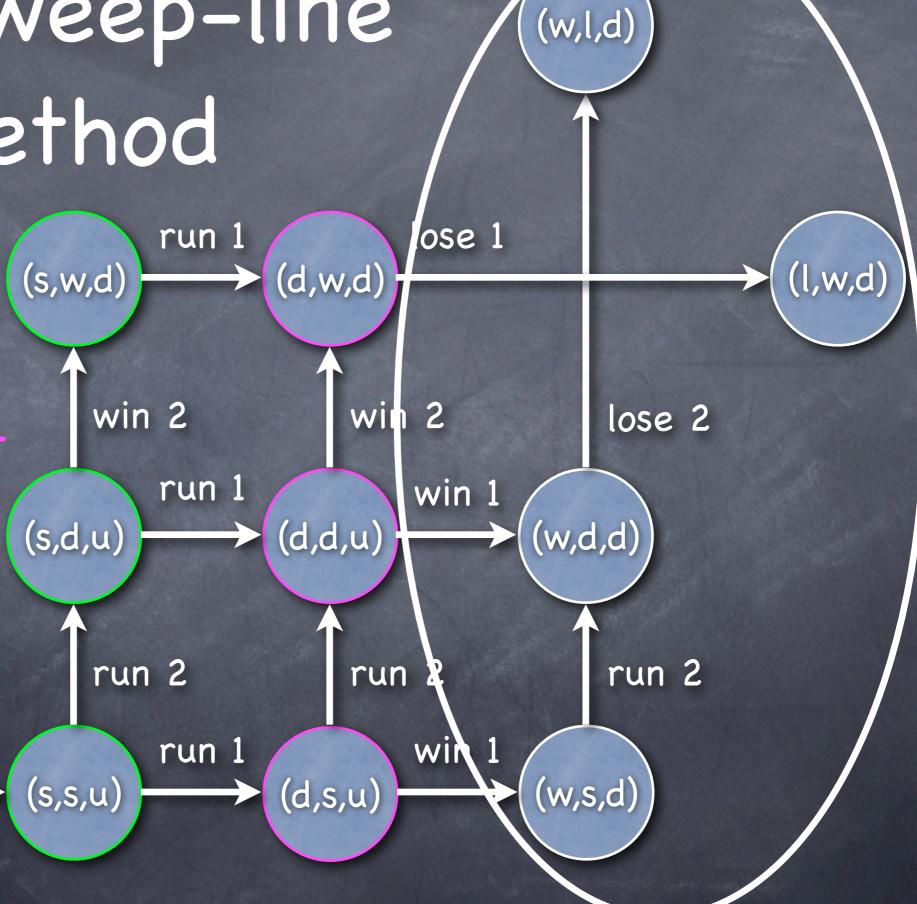
(w,l,d)

The Sweep-line Method

Already processed

Discovered but not yet processed

Not yet discovered

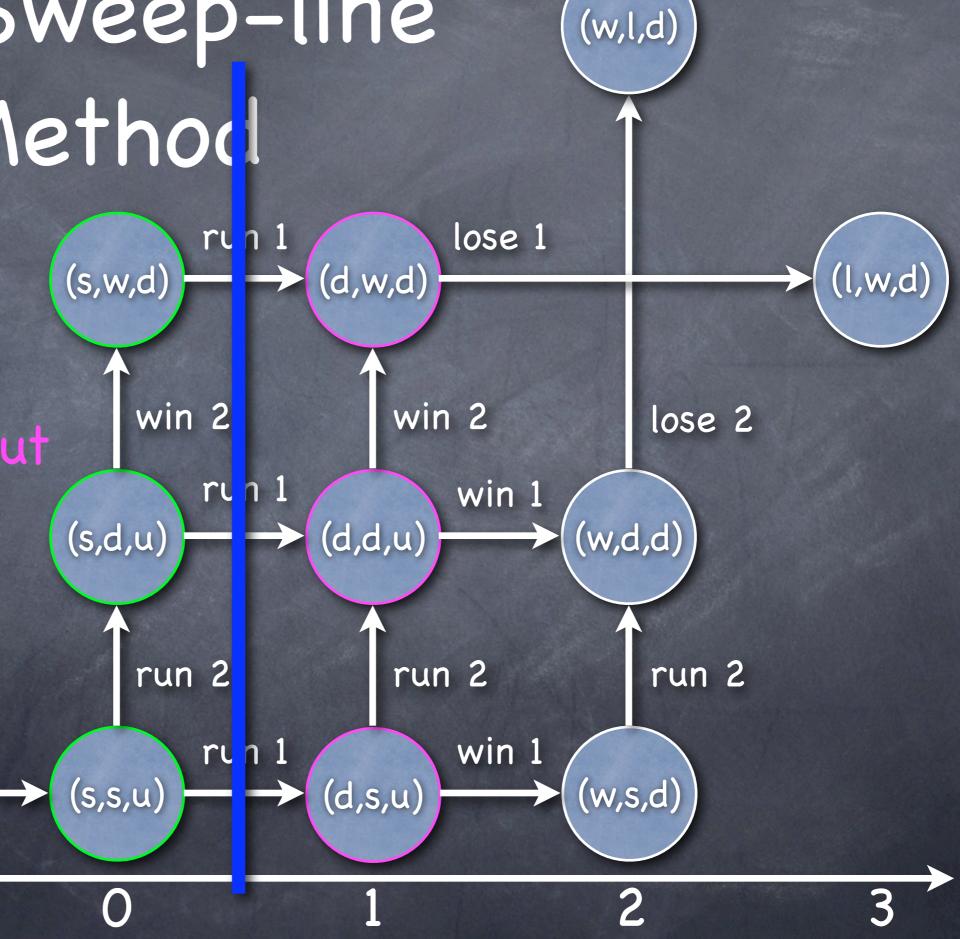


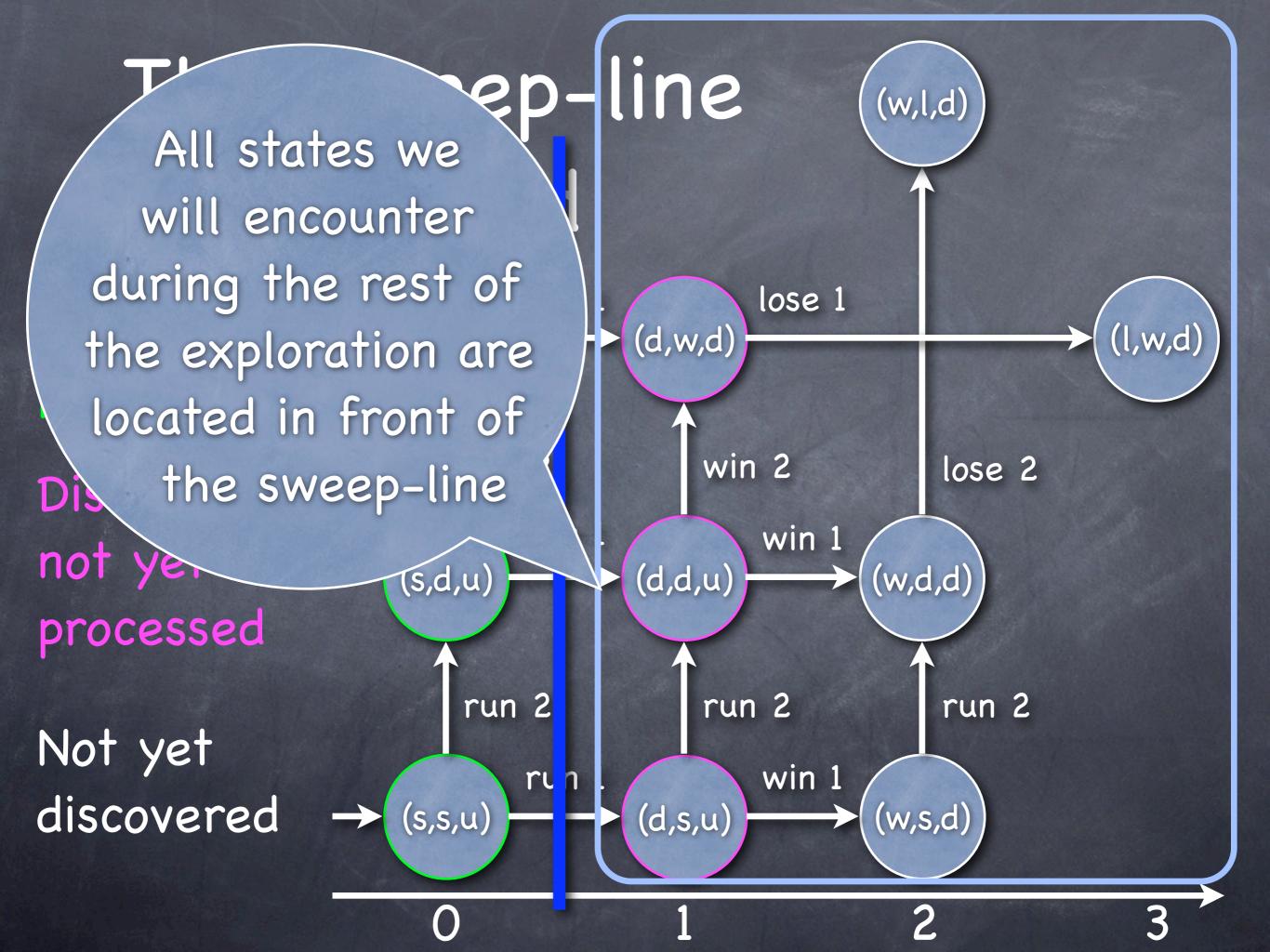
The Sweep-line Method

Already processed

Discovered but not yet processed

Not yet discovered



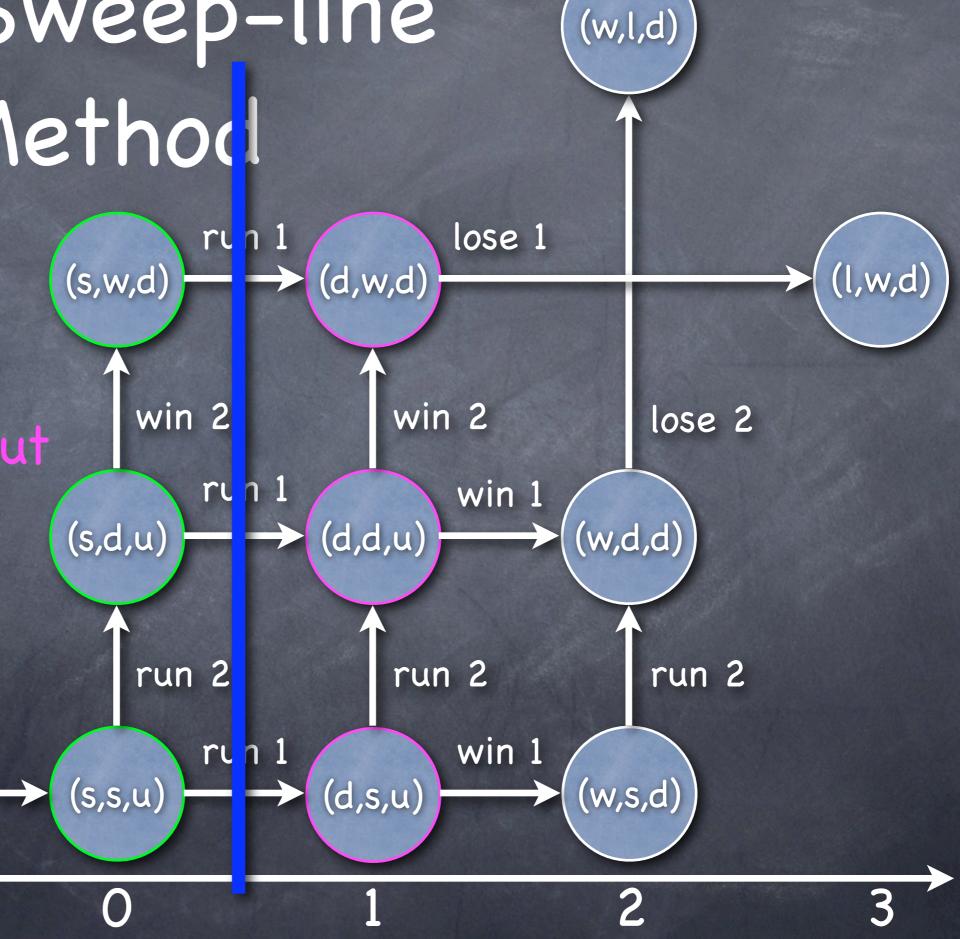


The Sweep-line Method

Already processed

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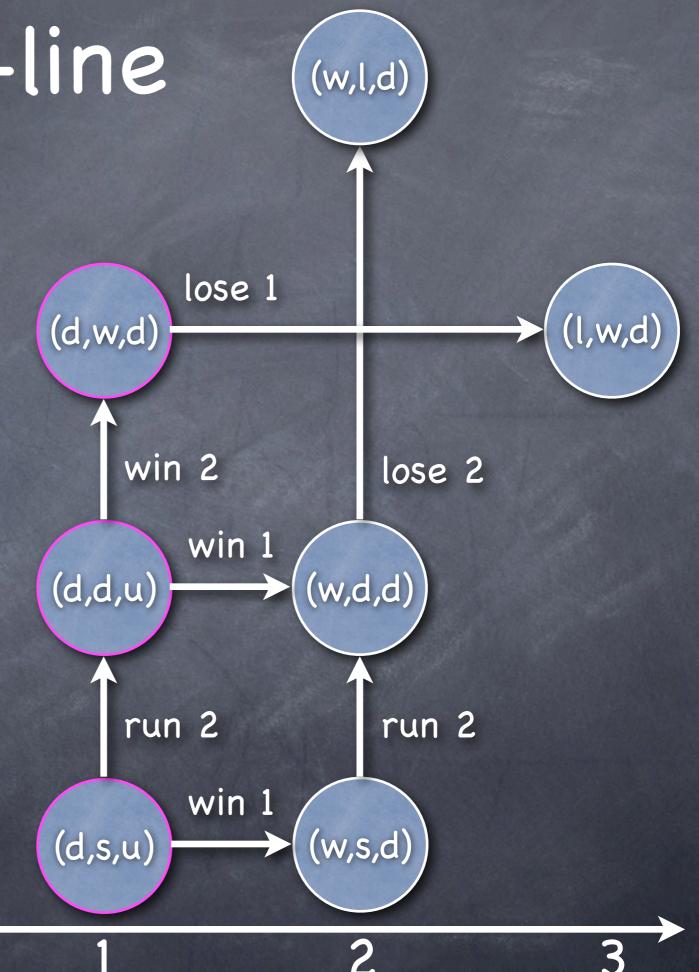


The Sweep-line Method

Already processed

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Not yet discovered

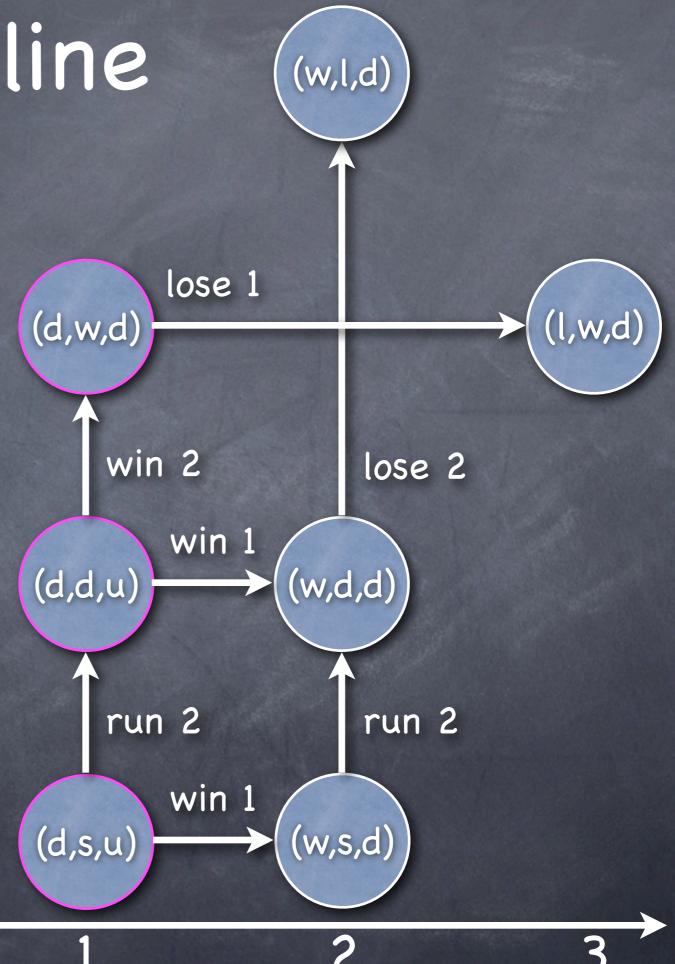


The Sweep-line Method

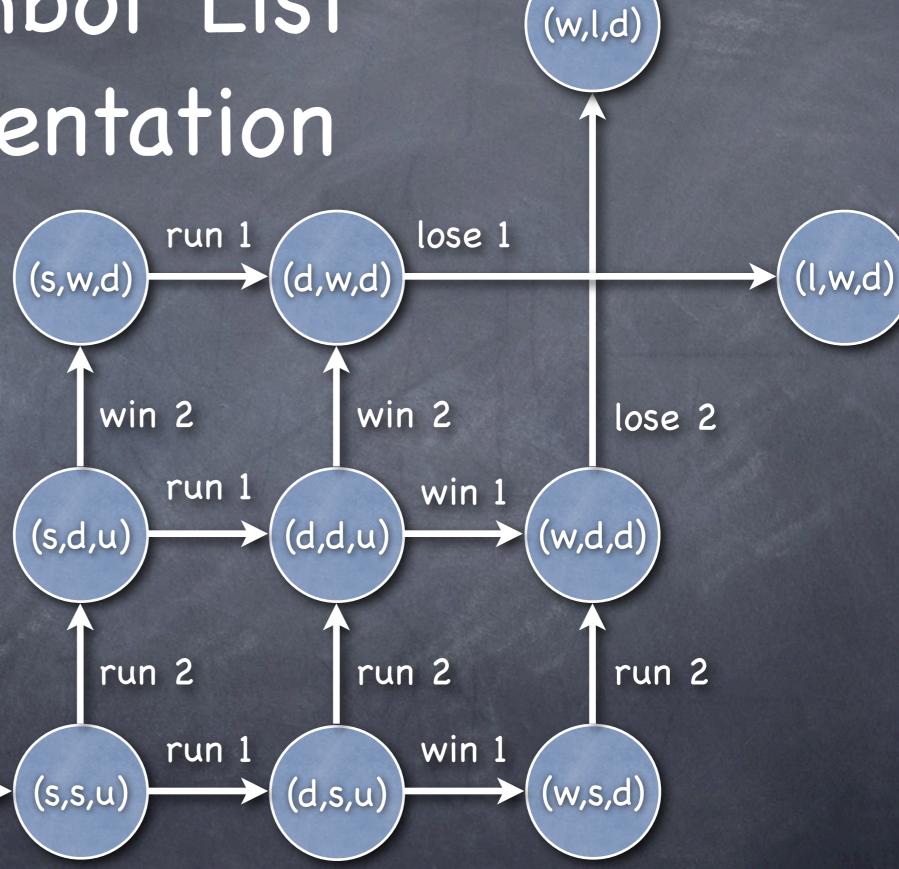
Already processed

Discovered but not yet processed

Not yet discovered



A Neighbor List Representation



Assume we can enumerate all transitions:

0: run 1

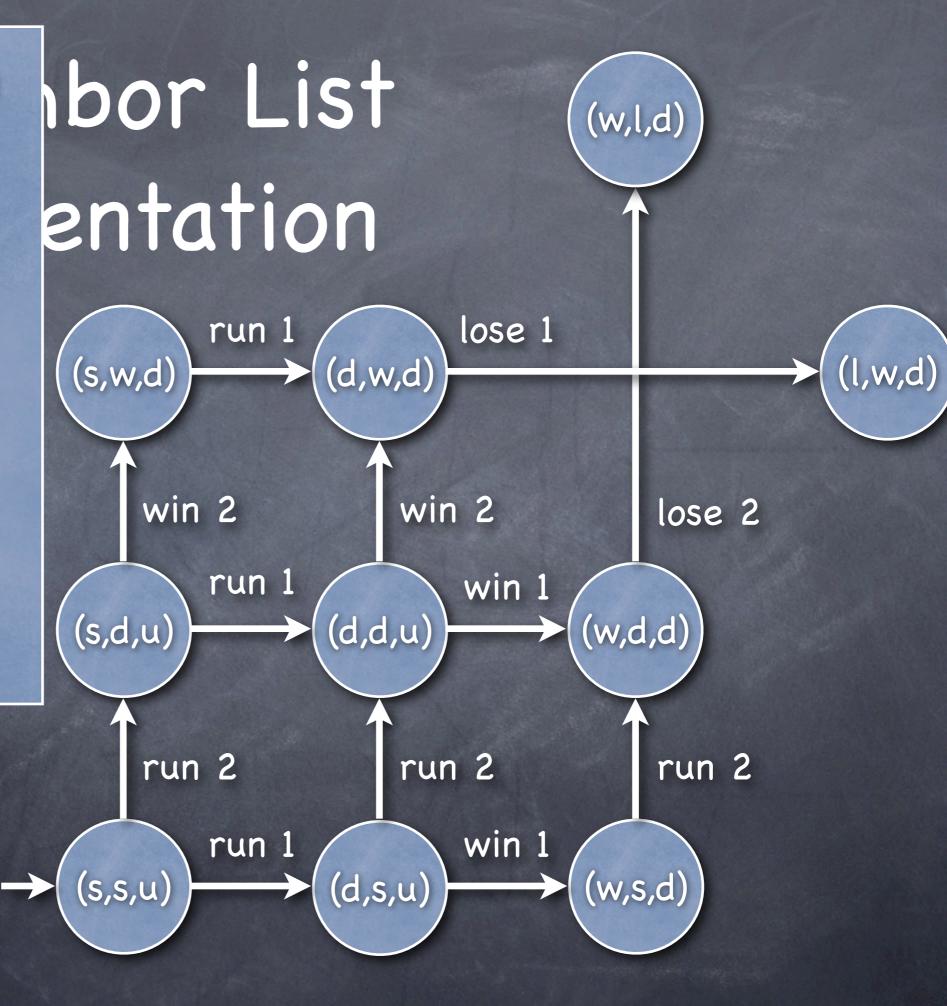
1: run 2

2: win 1

3: win 2

4: lose 1

5: lose 2



Assume we can enumerate all transitions:

0: run 1

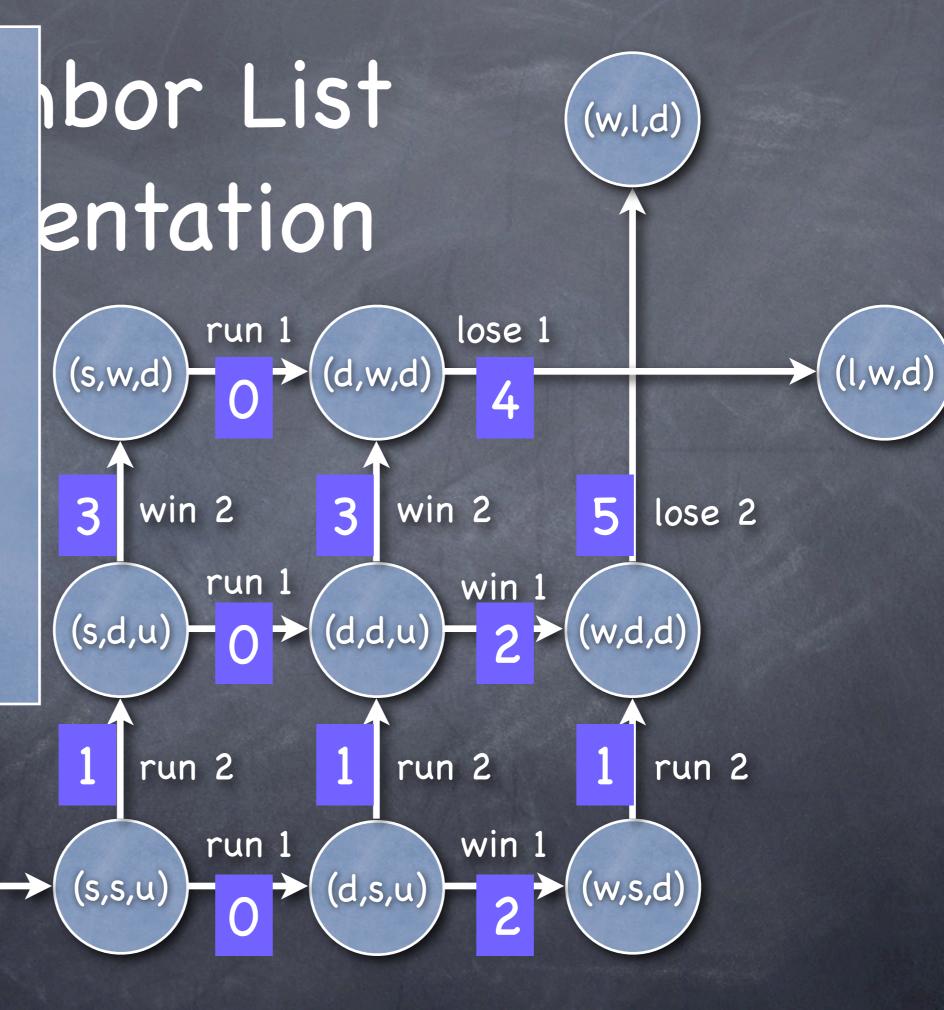
1: run 2

2: win 1

3: win 2

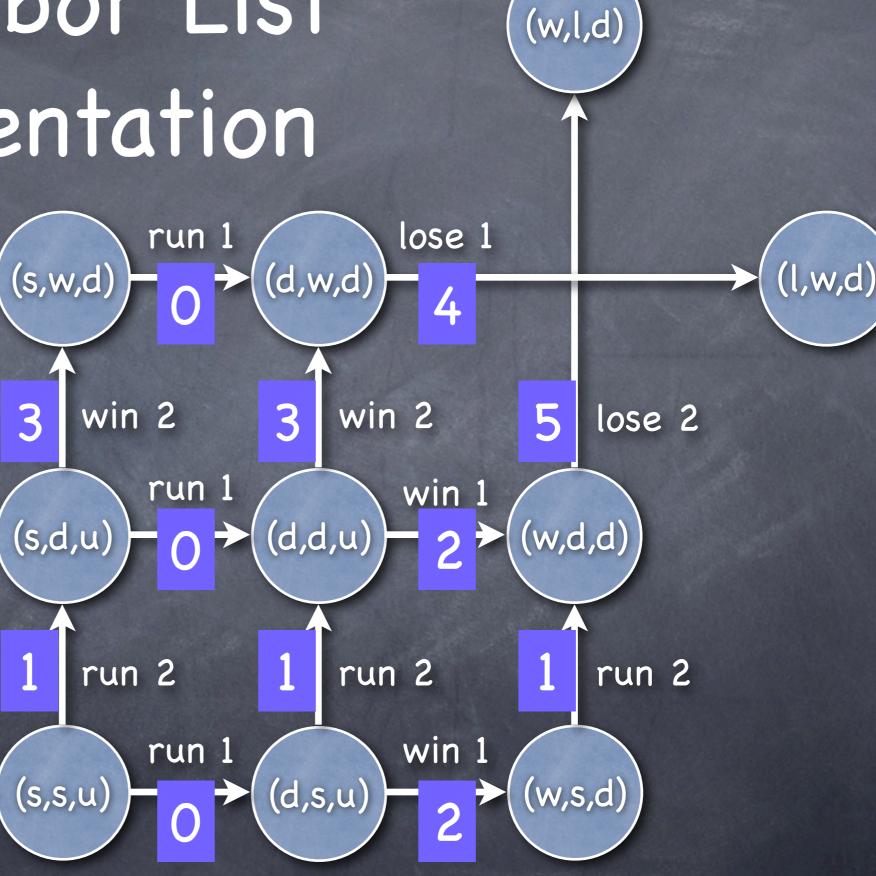
4: lose 1

5: lose 2



Assign a unique bor List number, 0...9, to each state

entation



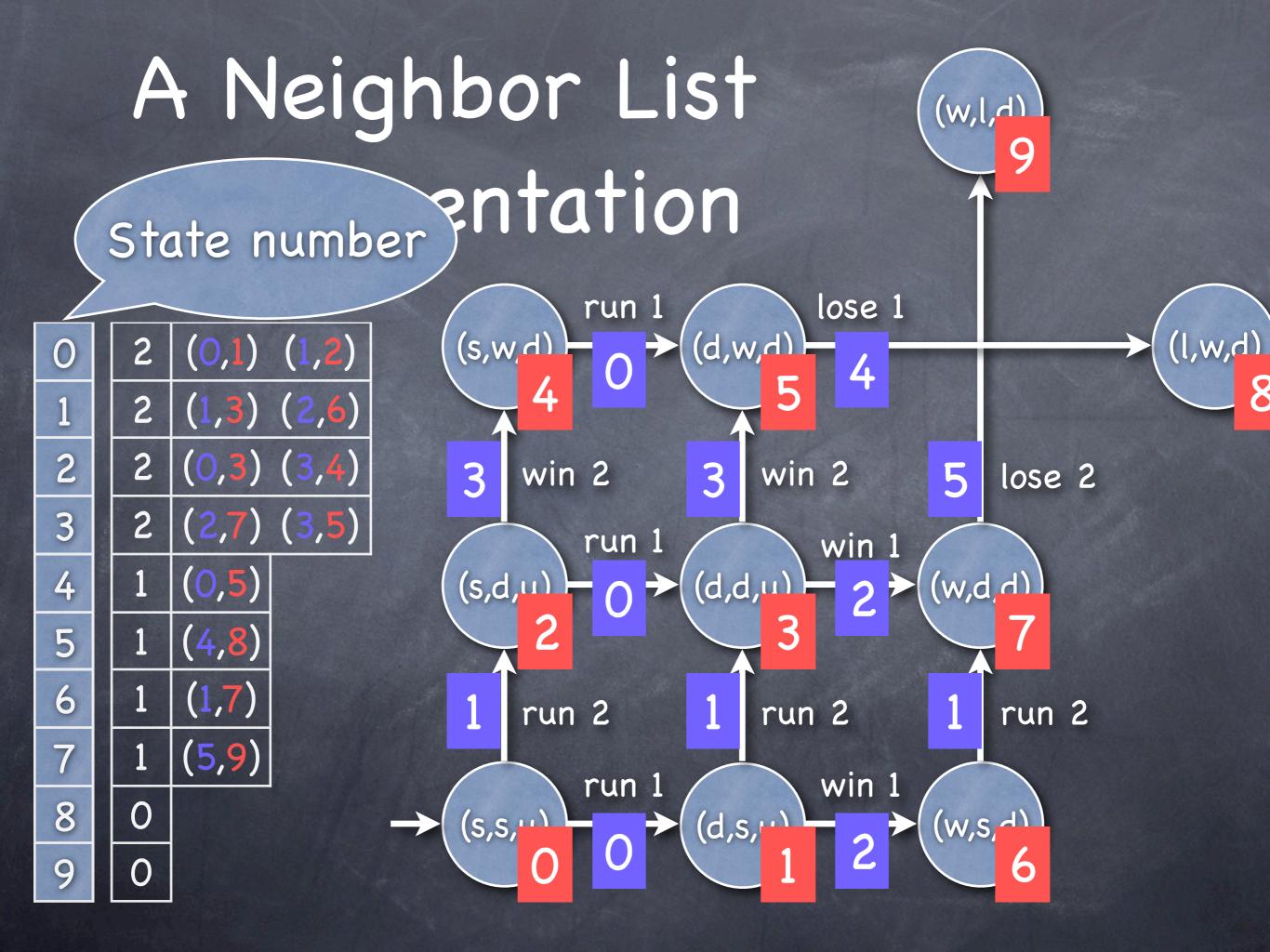
Assign a unique bor List (w,l,d) number, 0...9, to entation each state lose 1 run 1 (s,w,d) (d,w,d) (l,w,d)win 2 win 2 5 lose 2 run 1 win 1 (s,d,u) = 0 (w,d,d)(d,d,u) run 2 run 2 run 2 win 1 run 1 (d,s,1) (w,s, d)

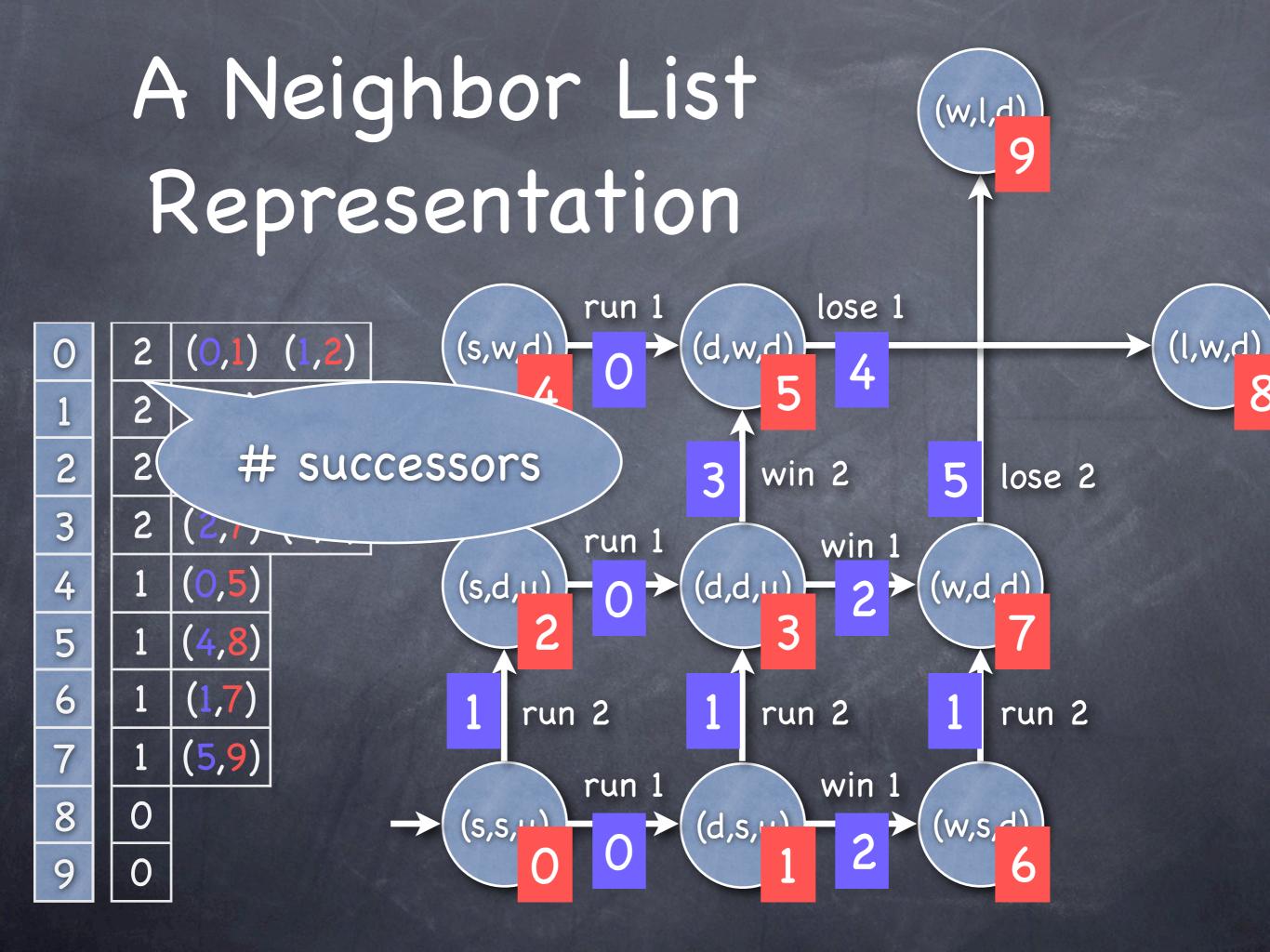
A Neighbor List (w,l,d)Representation run 1 lose 1 (l,w,d)(d,w,d) win 2 win 2 lose 2 run 1 win 1 (w,d,d)(d,d,u) run 2 run 2 run 2 win 1 run 1 (d,s,,,)

A Neighbor List Representation



(w,l,d)







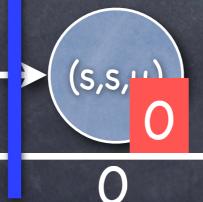




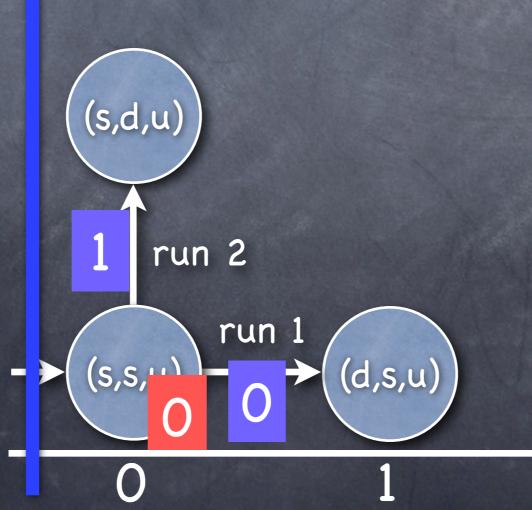
0 1 2



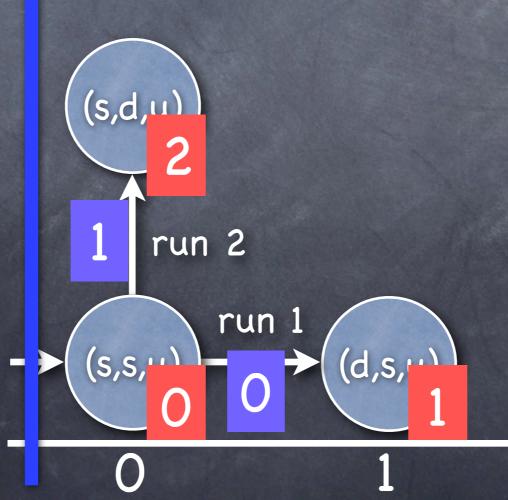


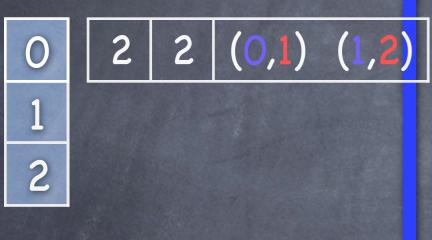


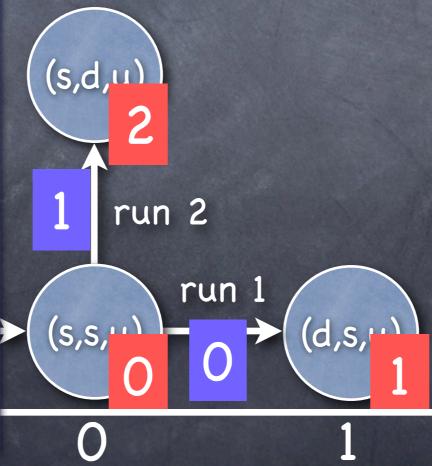


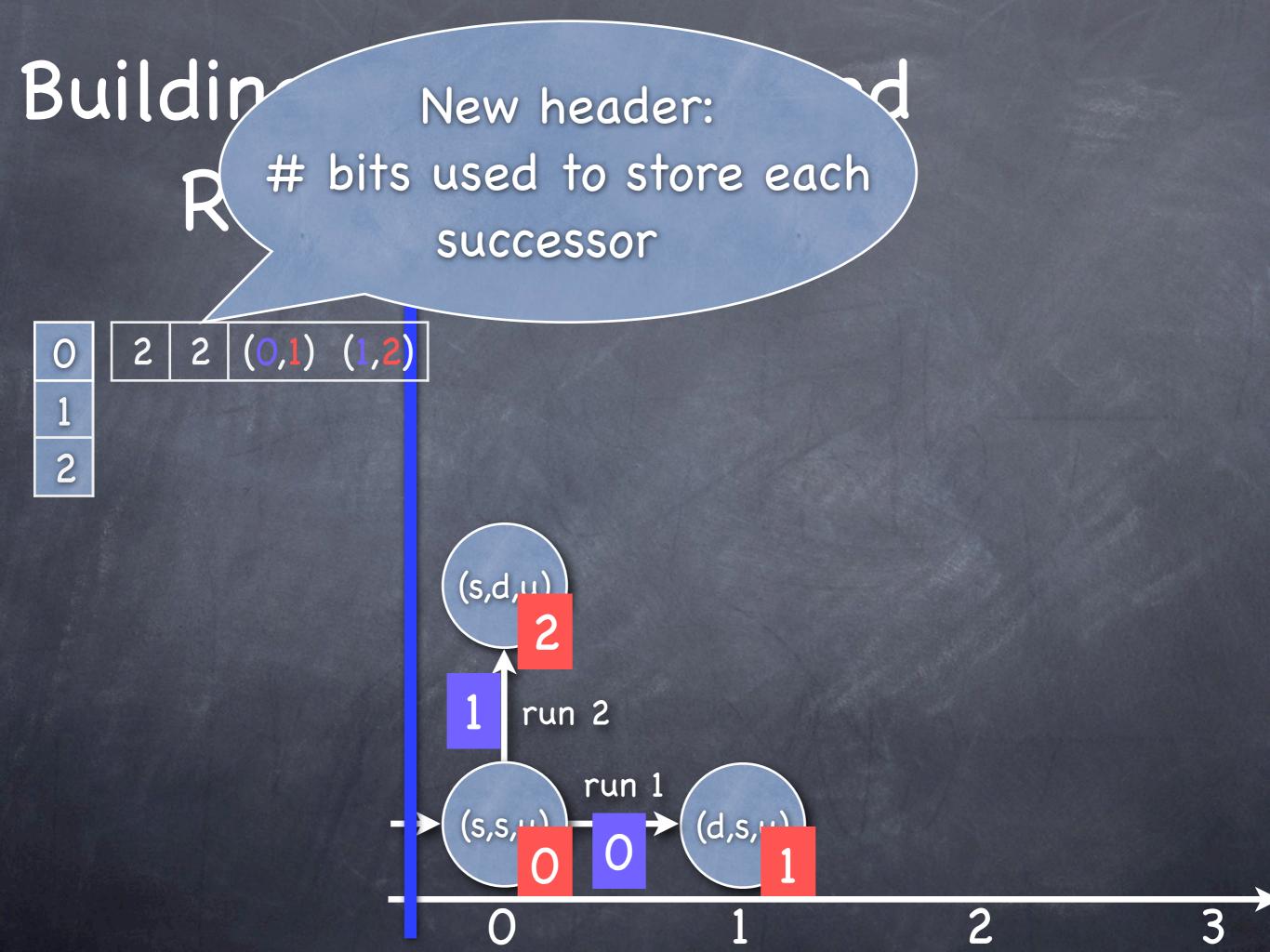


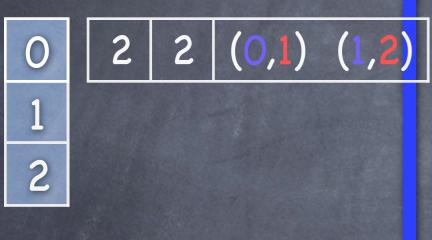


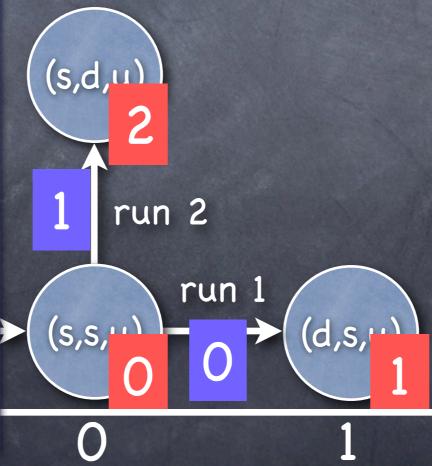


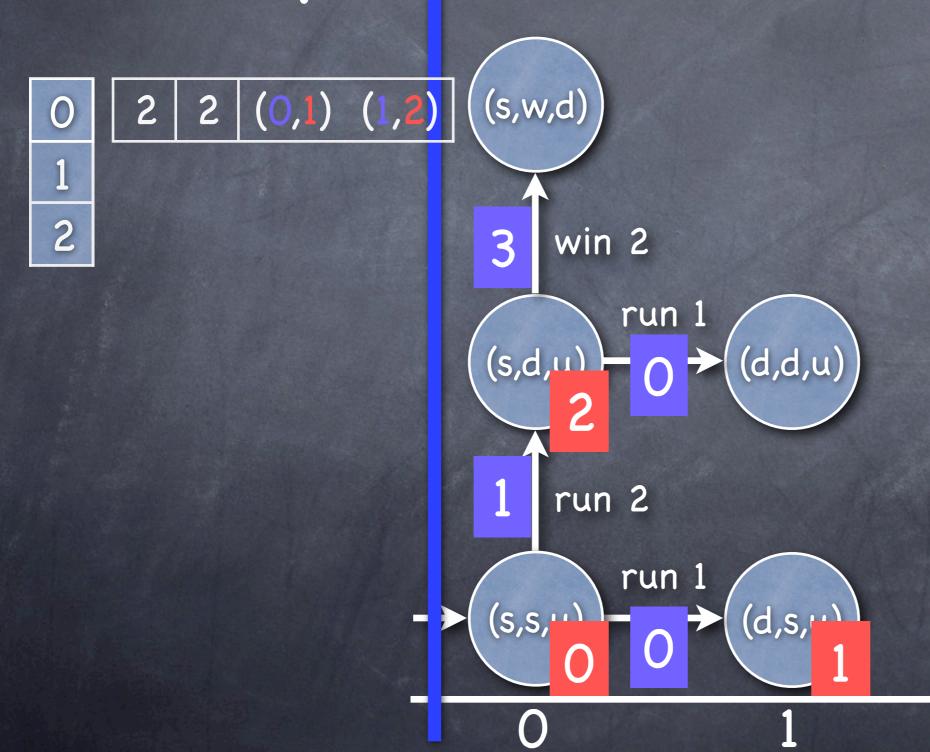


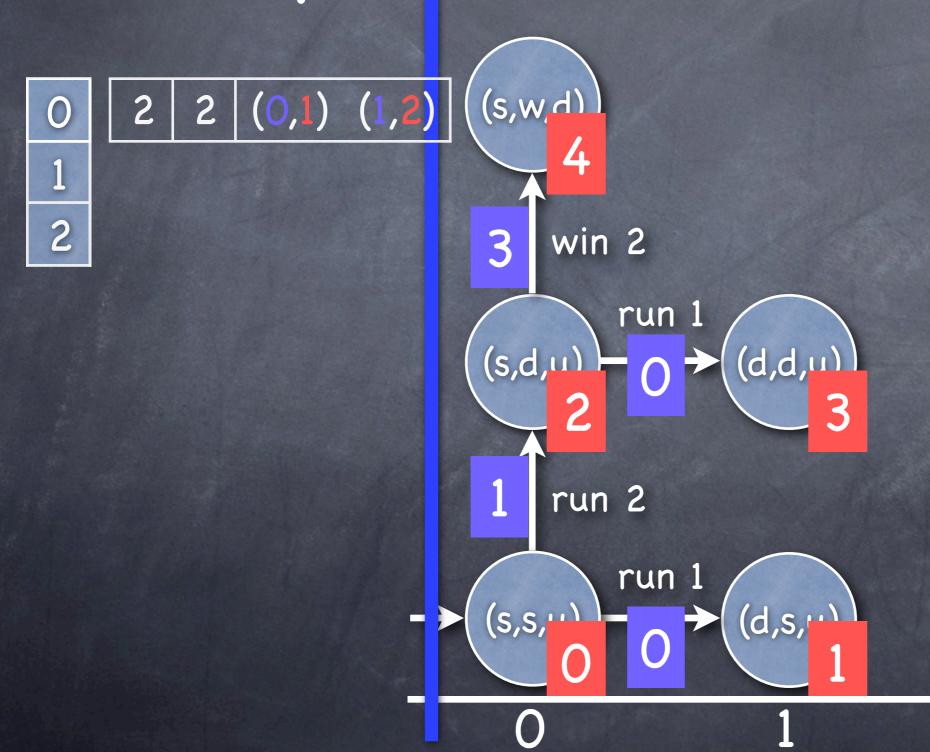


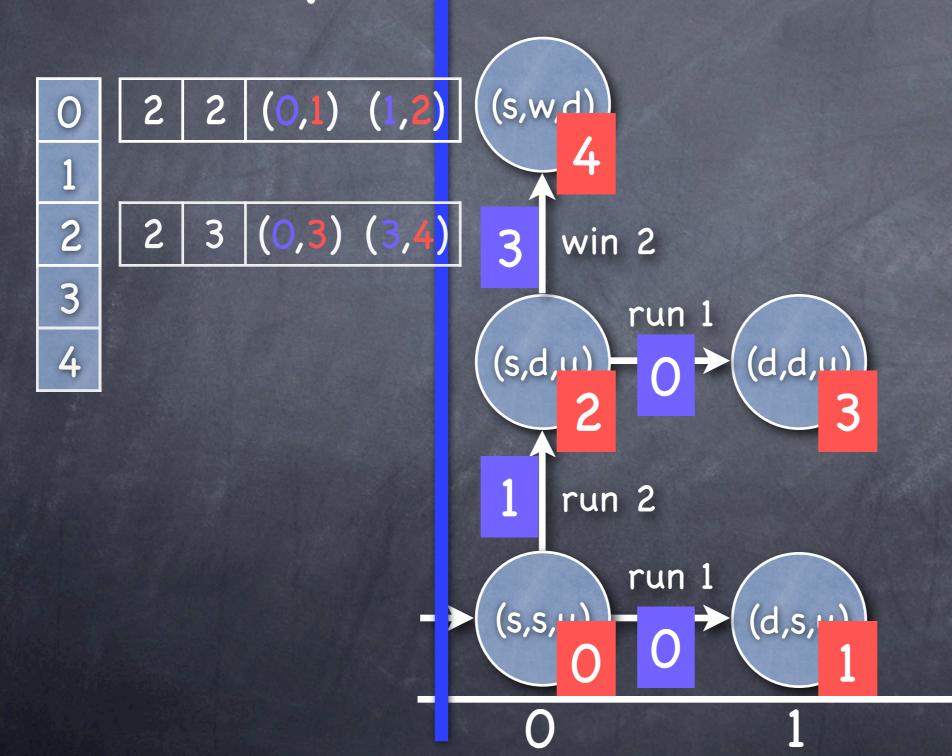


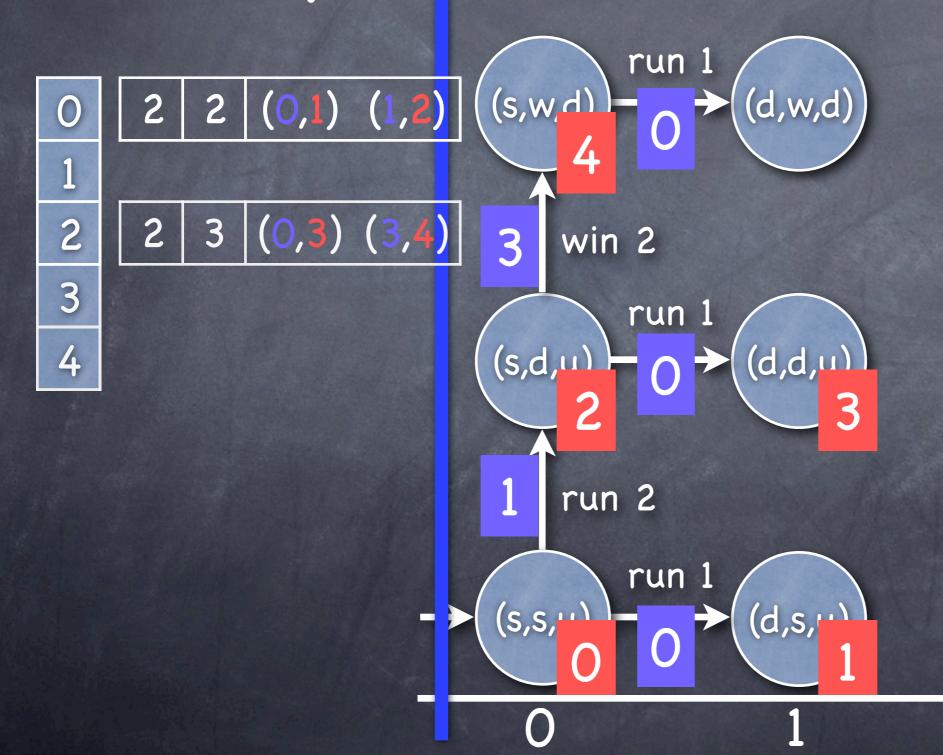


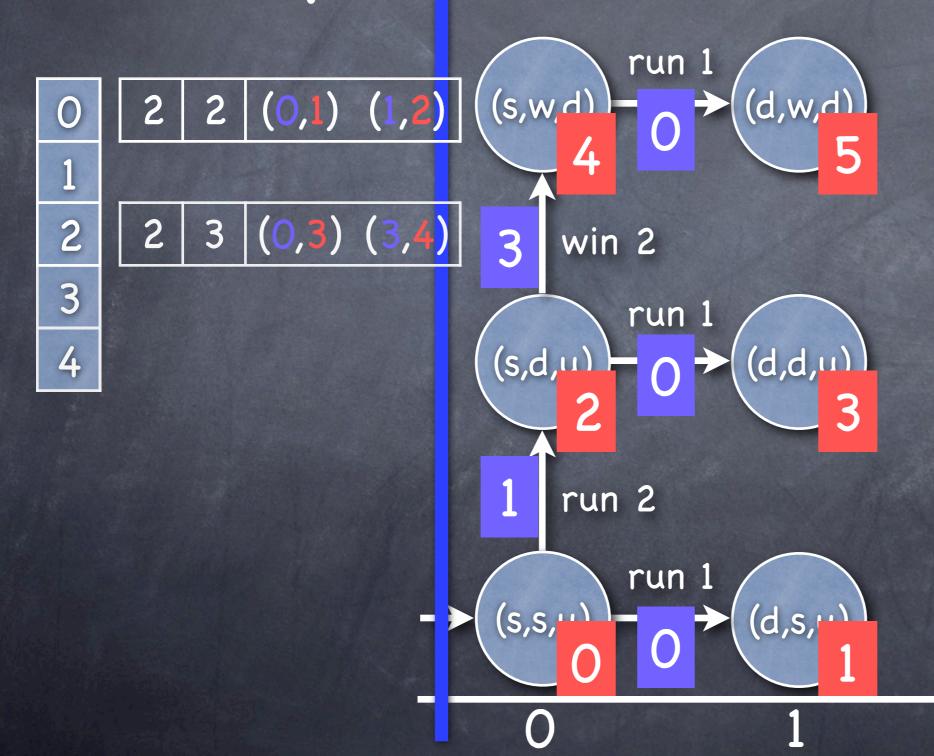


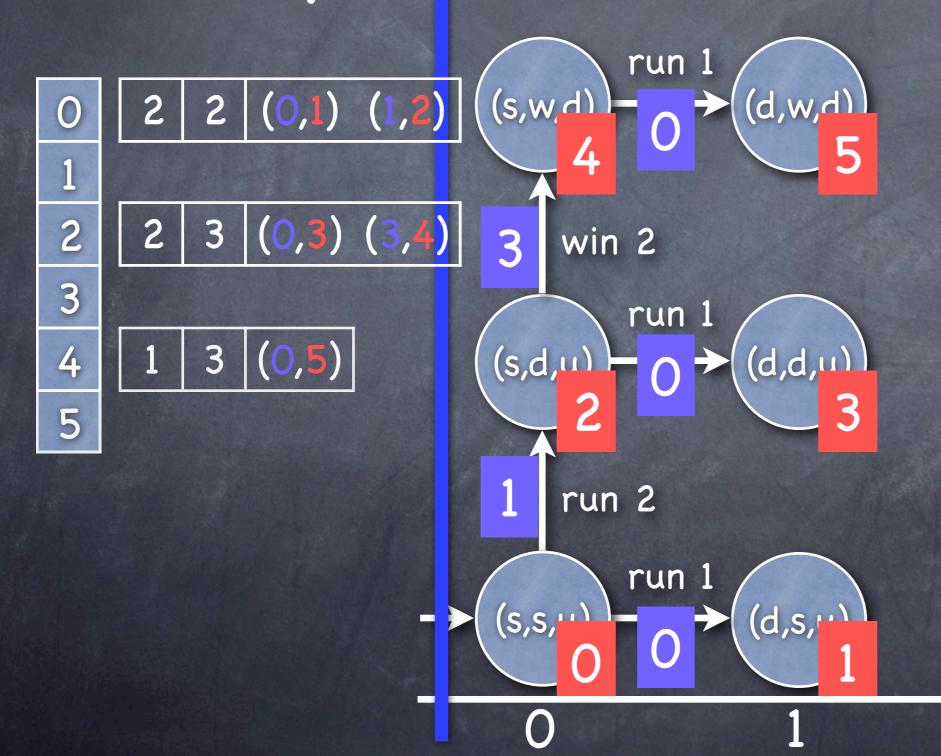


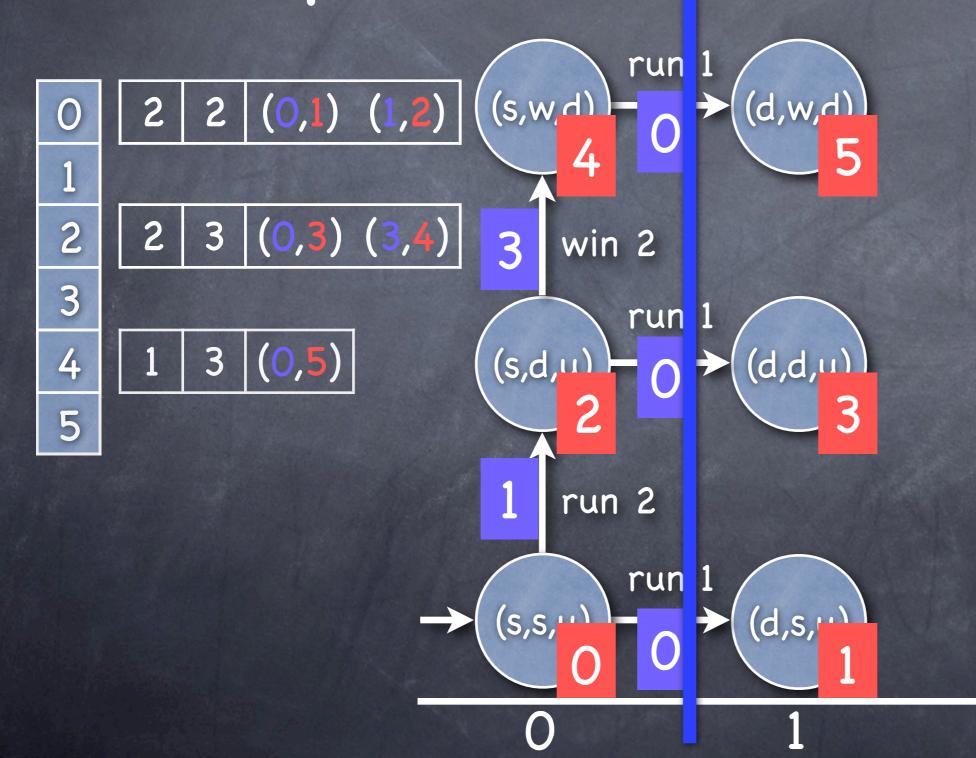


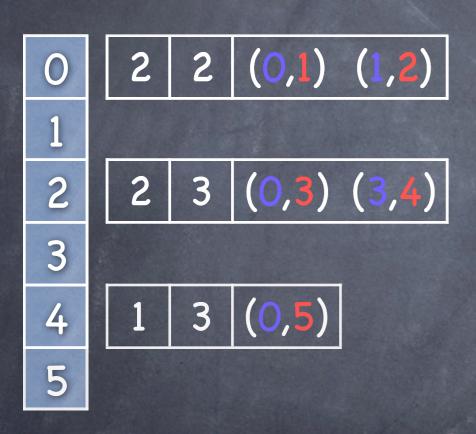








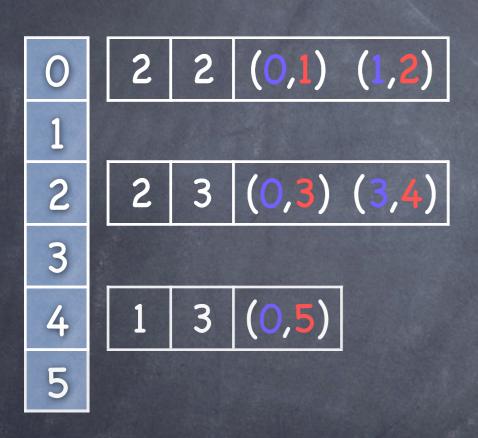




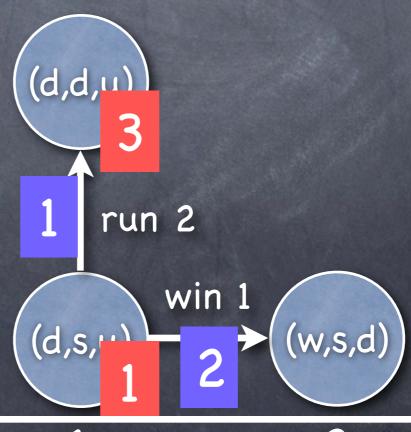




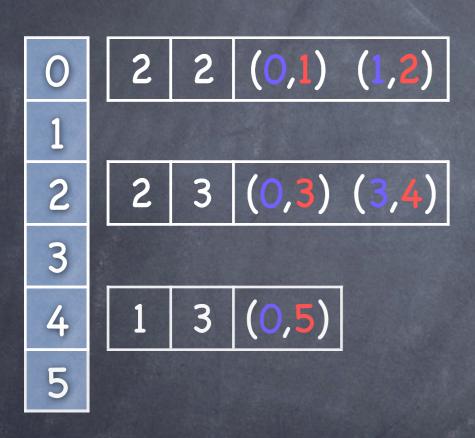




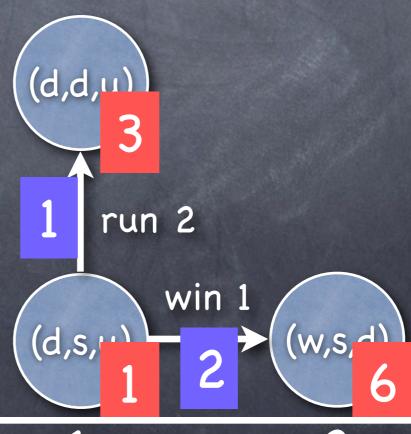




2



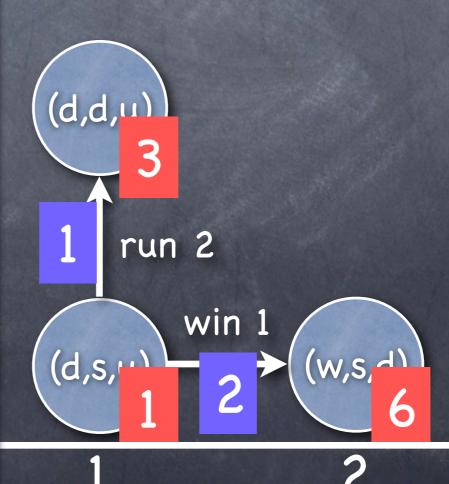


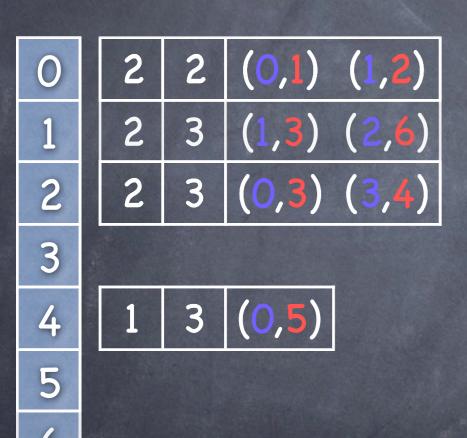


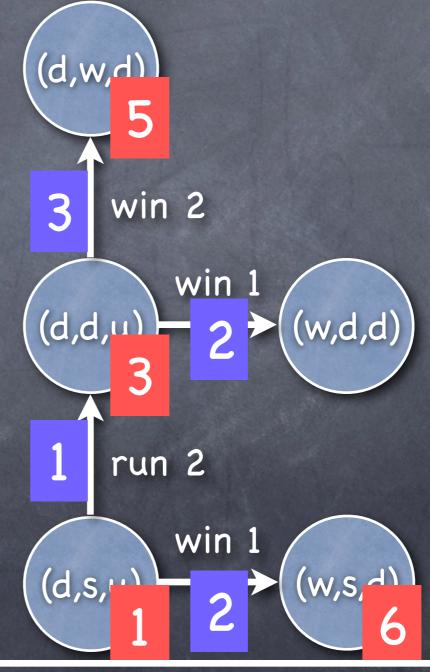
0







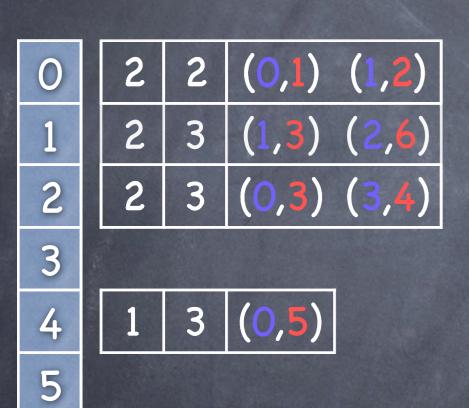




C

1

2





 0
 2
 2
 (0,1) (1,2)

 1
 2
 3
 (1,3) (2,6)

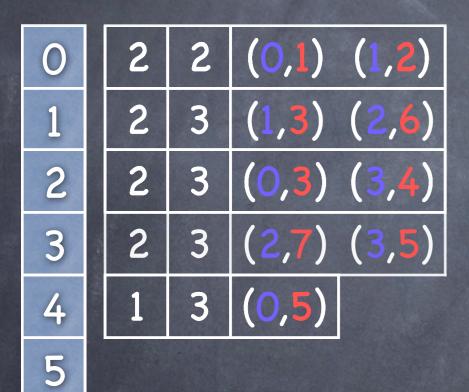
 2
 3
 (0,3) (3,4)

 3
 2
 3 (2,7) (3,5)

 4
 1
 3 (0,5)

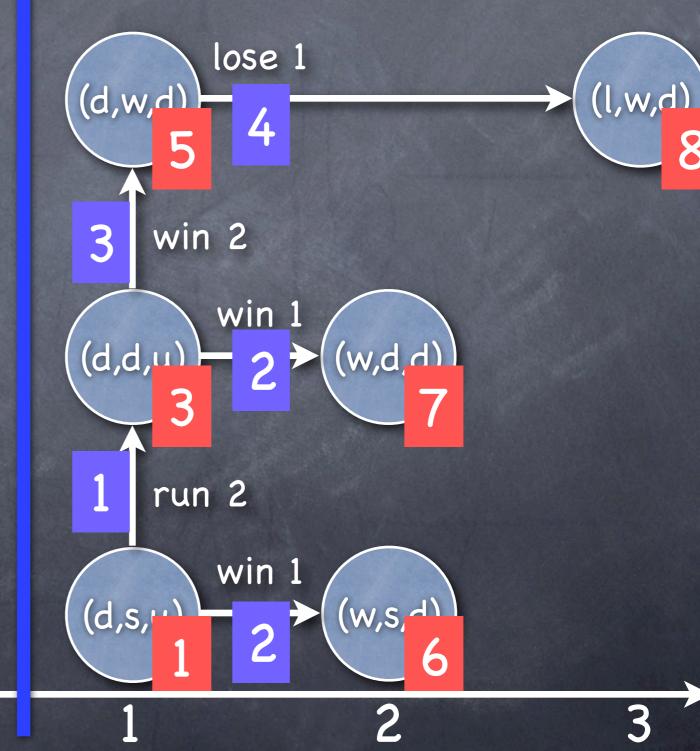
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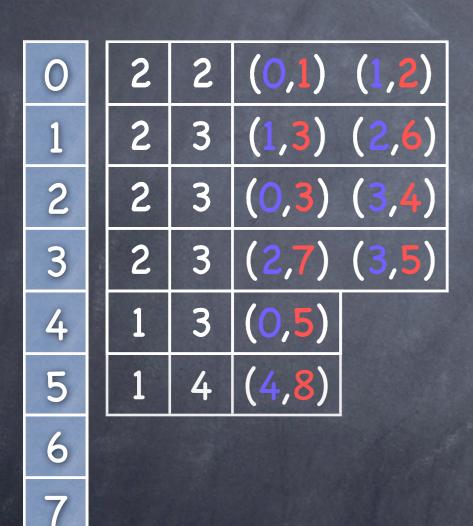




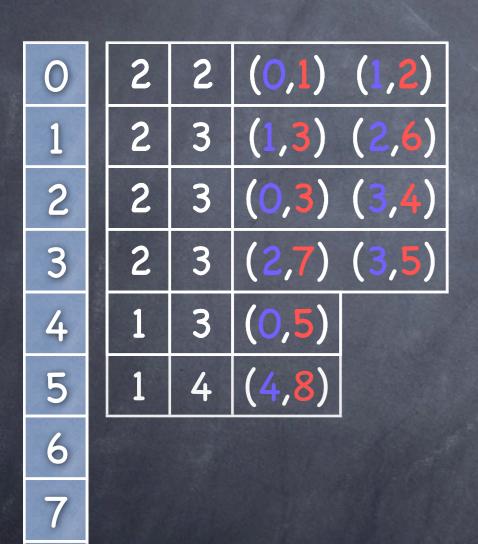


0	2	2	(0,1)	(1,2)
1	2	3	(1,3)	(2,6)
2	2	3	(<mark>0,3</mark>)	(3,4)
3	2	3	(2,7)	(3,5)
4	1	3	(0,5)	by
September 1			A BUSINESS	

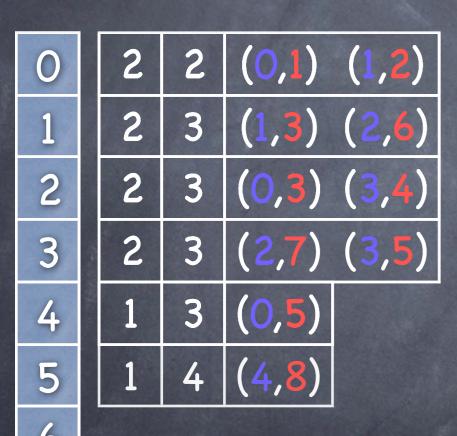








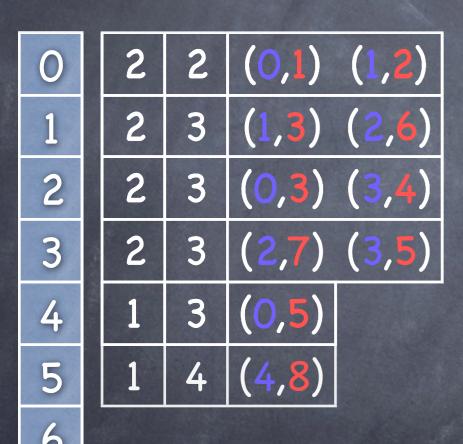




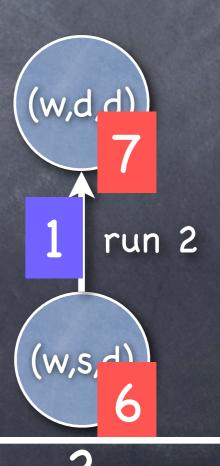


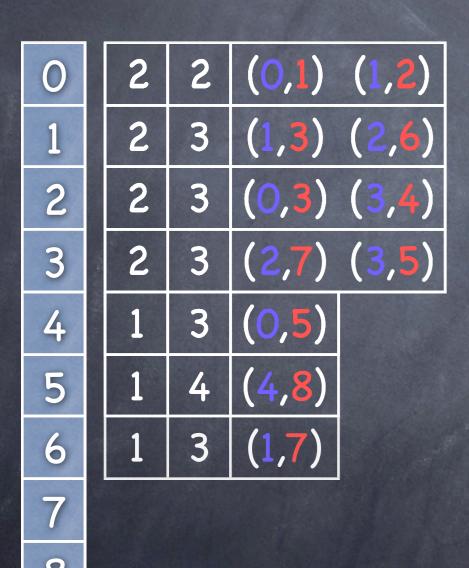






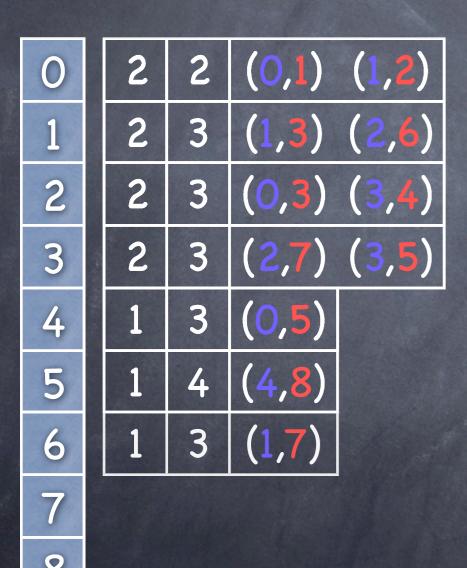


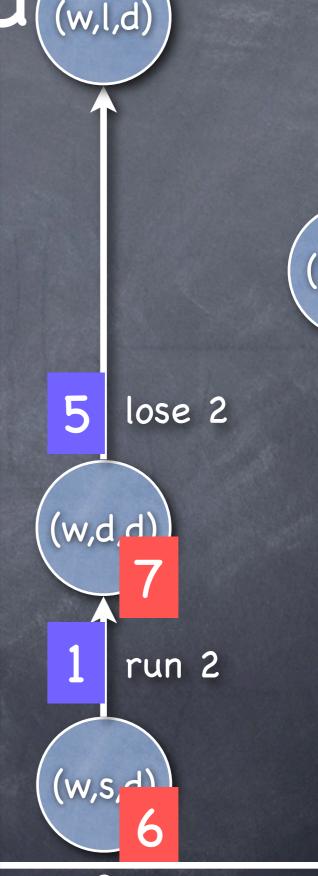


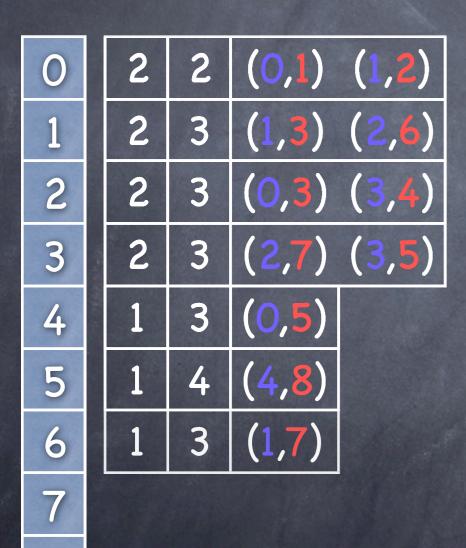




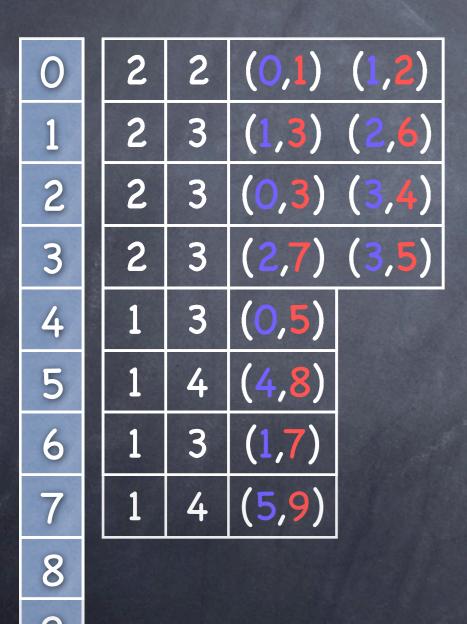




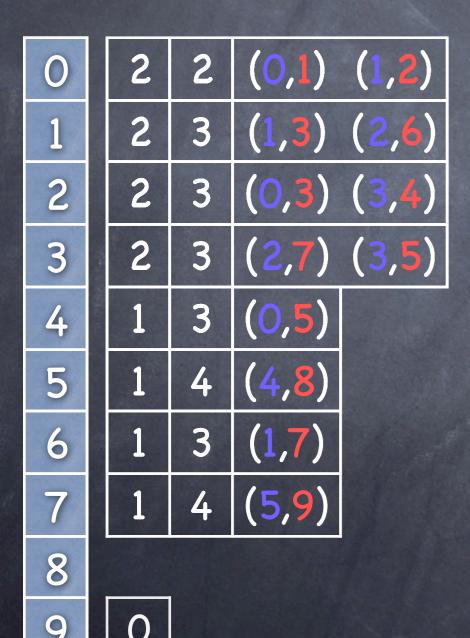




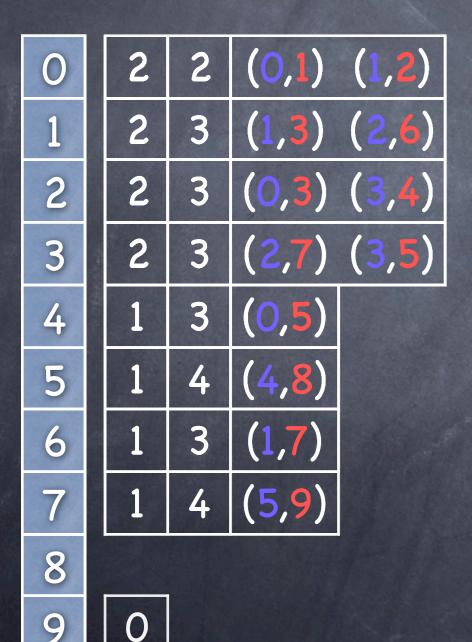




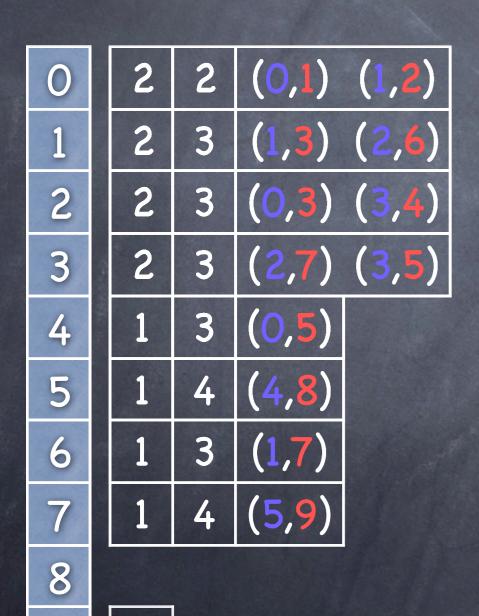




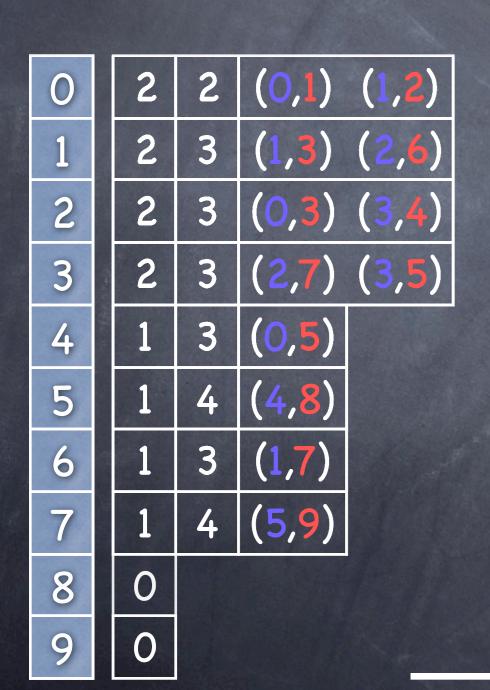




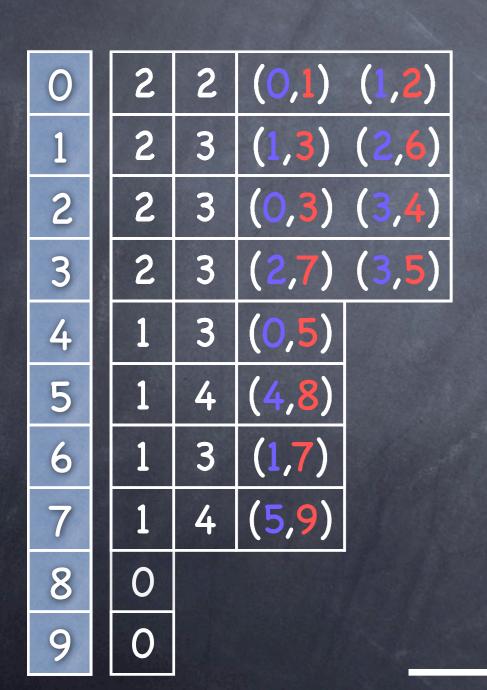




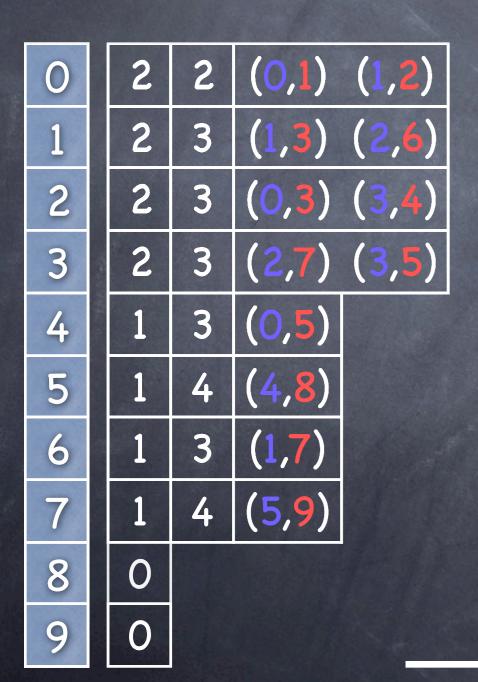












Evaluation of the Algorithm

- We only store a few (6) actual states
- The condensed representation uses $4 \cdot |R| \cdot w + |E| \cdot (\log |T| + \log |R|) + |F| \cdot \log |S|$ bits
 - R: all reachable states
 - w: size of machine word
 - © E: all reachable arcs
 - T: all transitions
 - S: all syntactically possible states
 - F: states on the front
- Efficient standard representation |R|·(3·w + log|S|) + |E|·log|S|

More Information about the Algorithm

T. Mailund, M. Westergaard: Obtaining Memory-Efficient Reachability Graph Representations Using the Sweep-Line Method, TACAS 2004

We need to support the operations add(s) and contains(s) but not any get operation

- We need to support the operations add(s) and contains(s) but not any get operation
- Idea of hash compaction: Store hash value for each state only

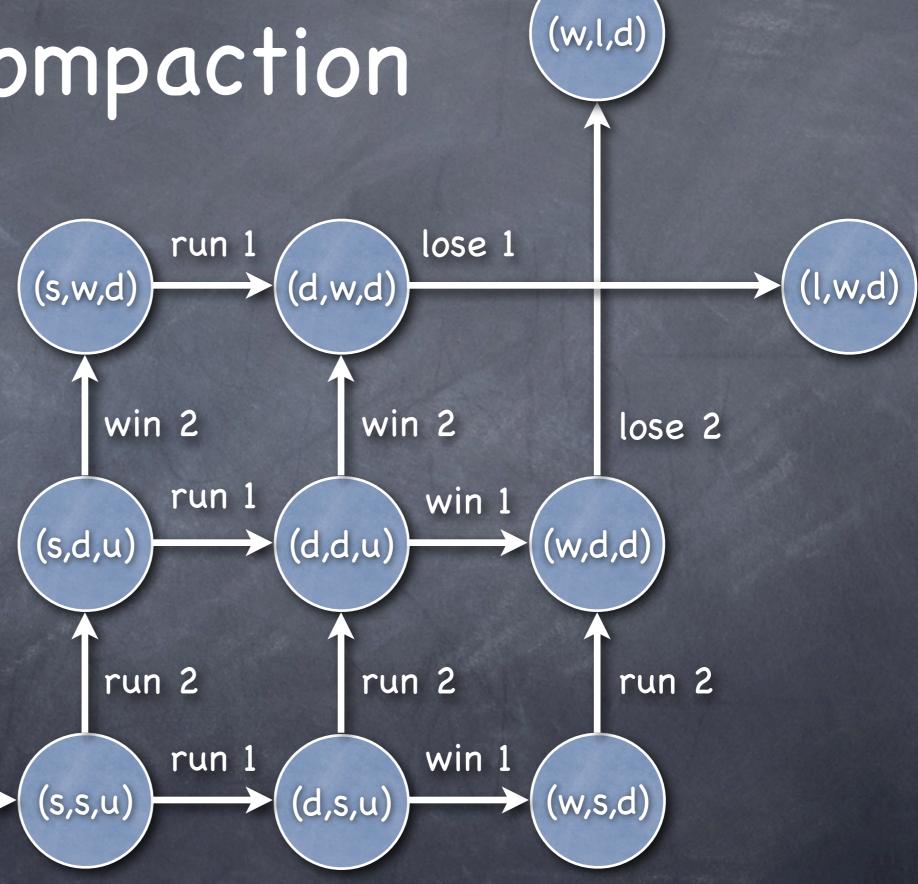
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 - Hash compaction fails

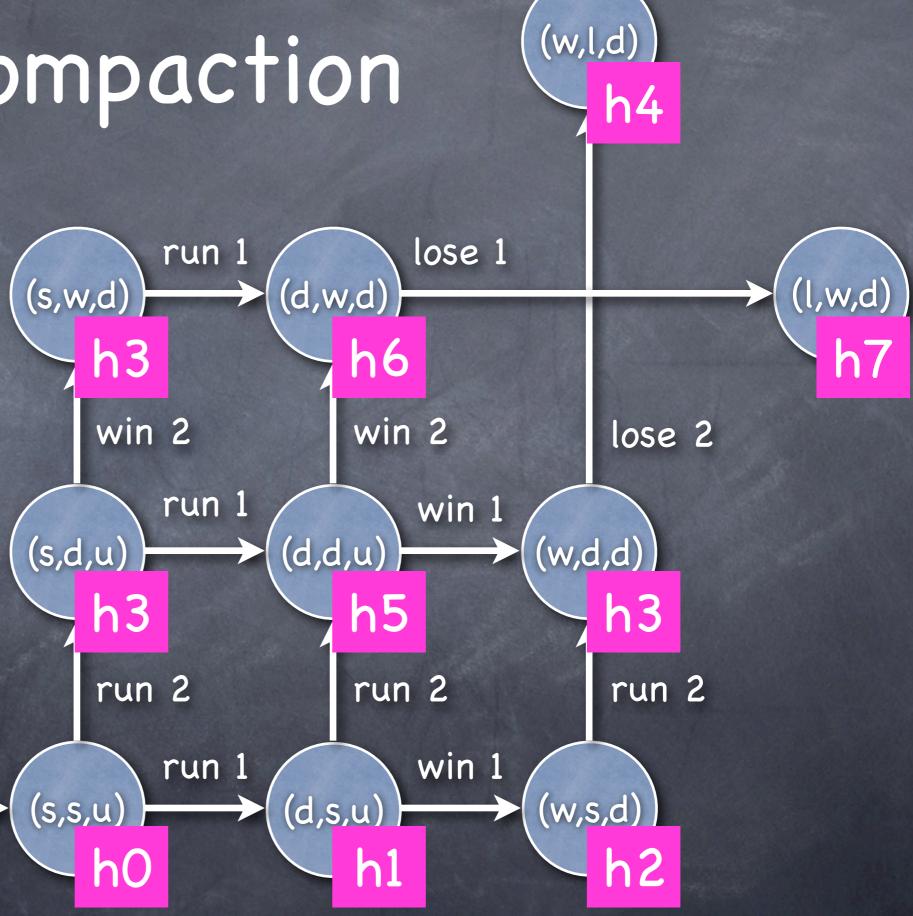
- We need to support the operations add(s) and contains(s) but not any get operation
- Idea of hash compaction: Store hash value for each state only
- Hash collision?
 - Hash compaction fails
 - We will store for each state a predecessor and trace from the initial state

(w,l,d) Hash Compaction lose 1 run 1 (s,w,d) (d,w,d) (l,w,d) win 2 win 2 lose 2 run 1 win 1 (s,d,u) (d,d,u) (w,d,d) run 2 run 2 run 2 win 1 run 1 (w,s,d) (s,s,u) (d,s,u)

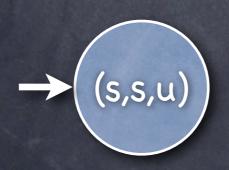
We assume a hash function, assigning to each state a hash-value



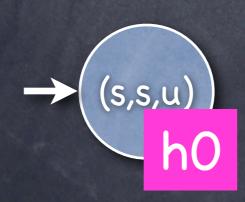
We assume a hash function, assigning to each state a hash-value



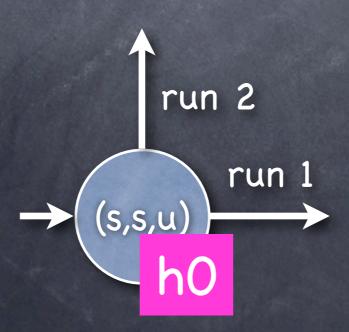




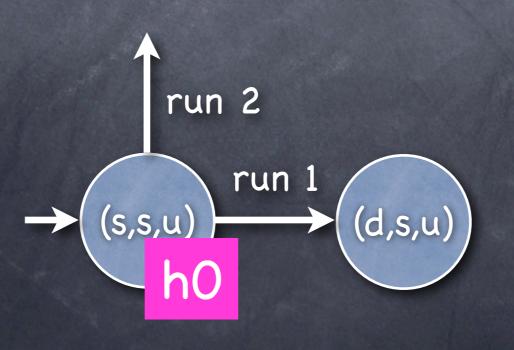
hO



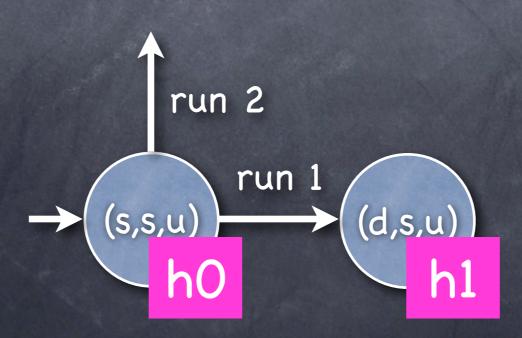
hO



hO



h0 h1



(w,l,d) h4

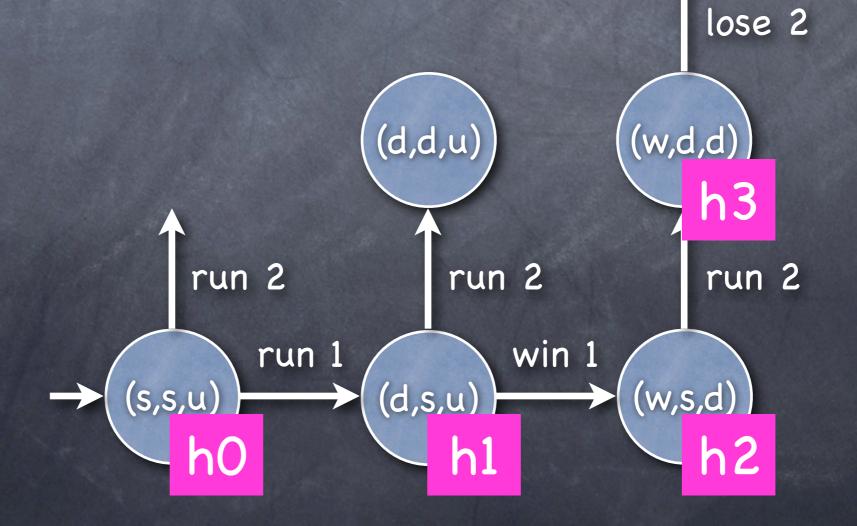
lose 2

h0
h1
h2
h3

run 2 run 2 run 2 run 2 (w,d,d)h3
run 2 (u,d,d)h3
run 2 (u,d,d)h0 (u,d,d)h3
run 2 (u,d,d)h1 (u,d,d)h2

(w,l,d) | h4

h0 h1 h2 h3



(w,l,d) | h4

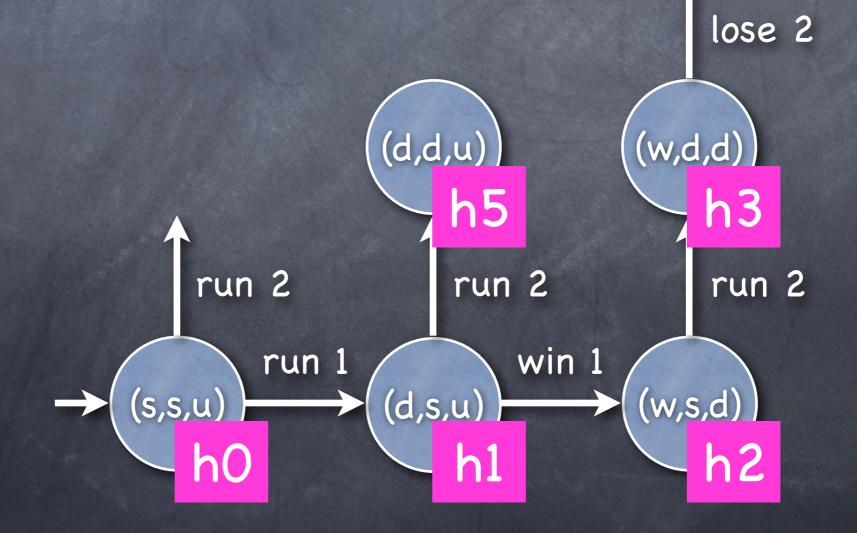
h0 h1

h2

h3

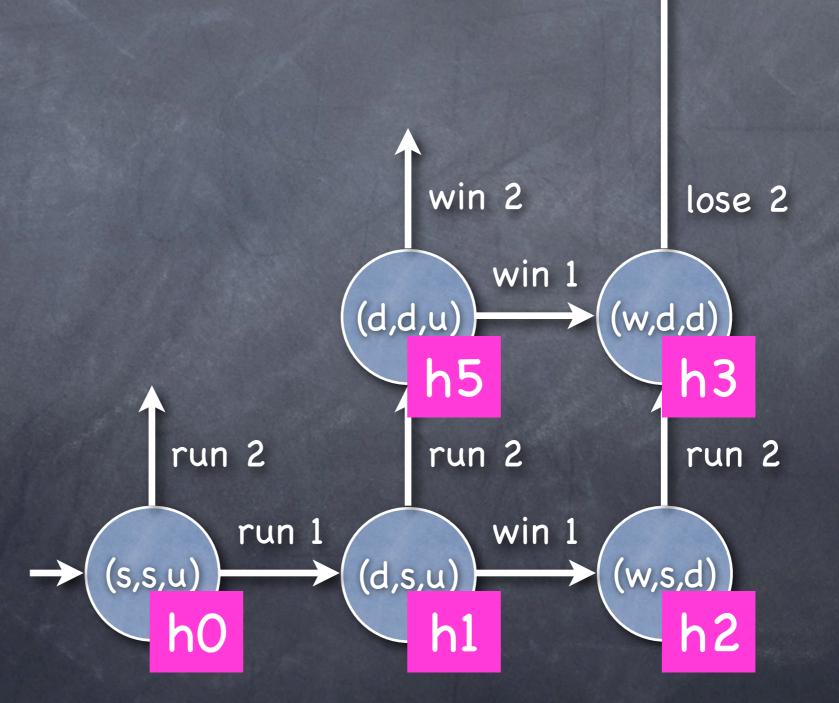
h4

h5



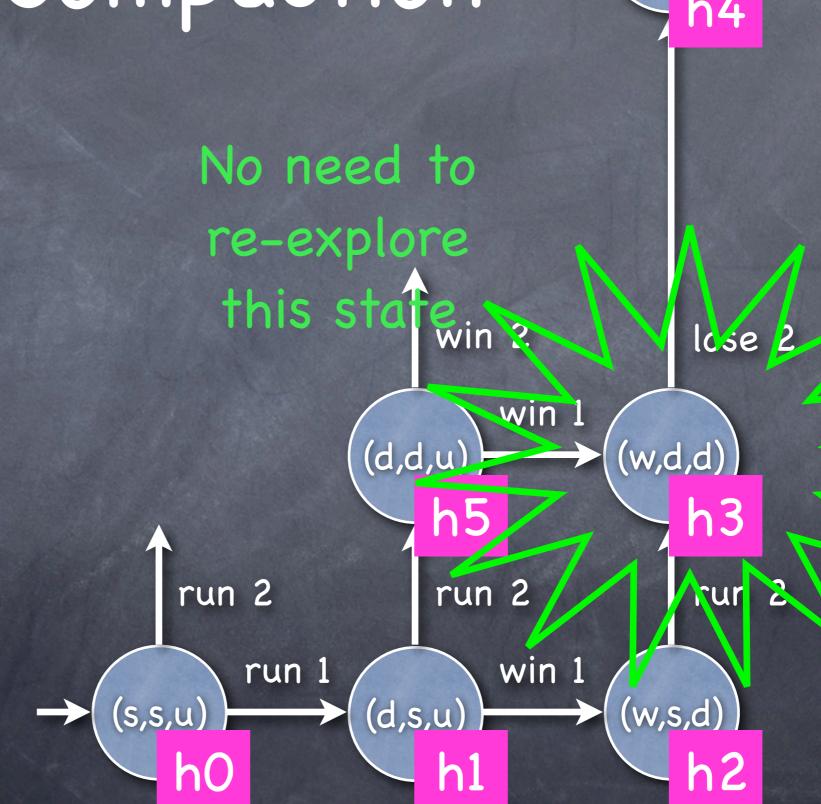
(w,l,d) | h4

h0 h1 h2 h3 h4



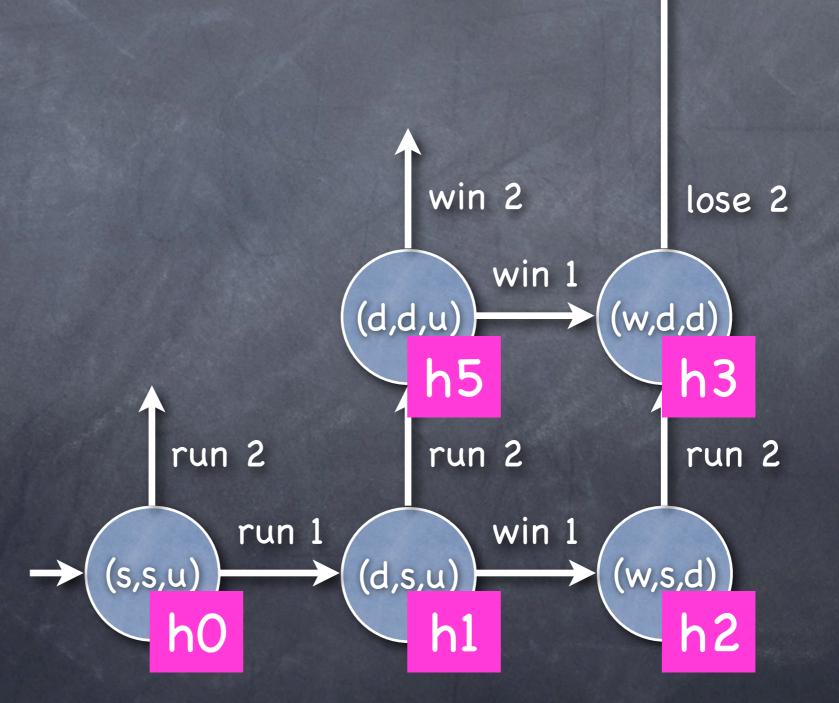
(w,l,d) \\ h4

h0
h1
h2
h3
h4



(w,l,d) | h4

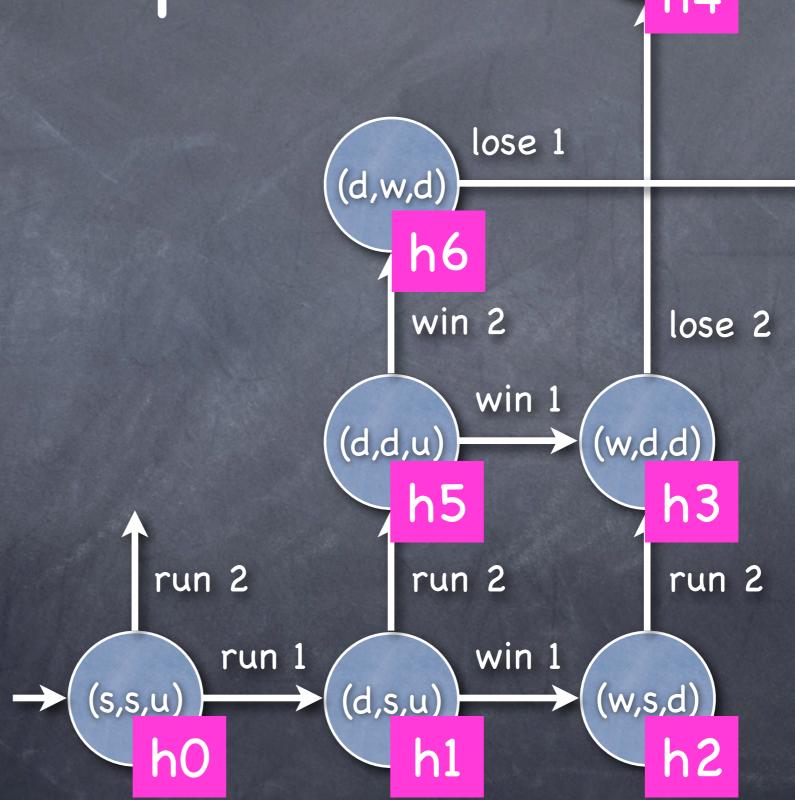
h0 h1 h2 h3 h4



(w,l,d) h4

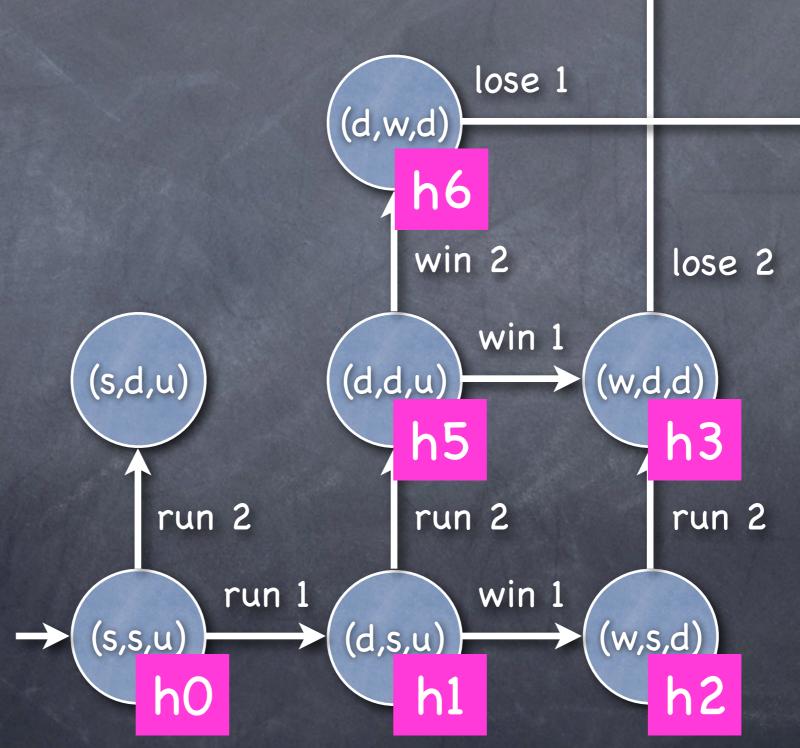
(l,w,d)





(w,l,d) h4 ((l,w,d)) h7

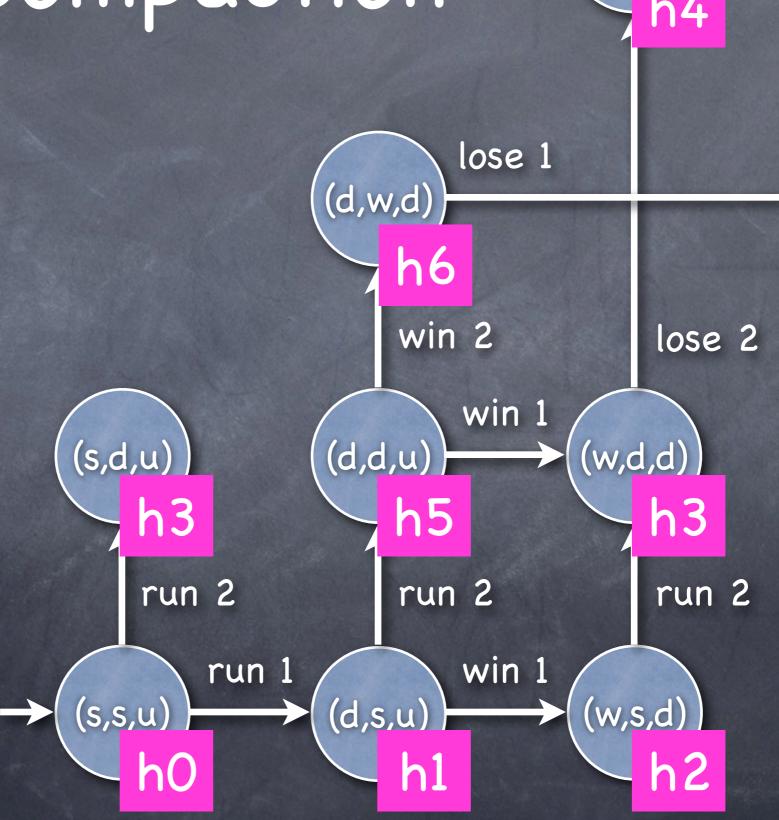
hO h1 h2 h3 h4 h5 h6



(w,l,d) h4

(l,w,d)



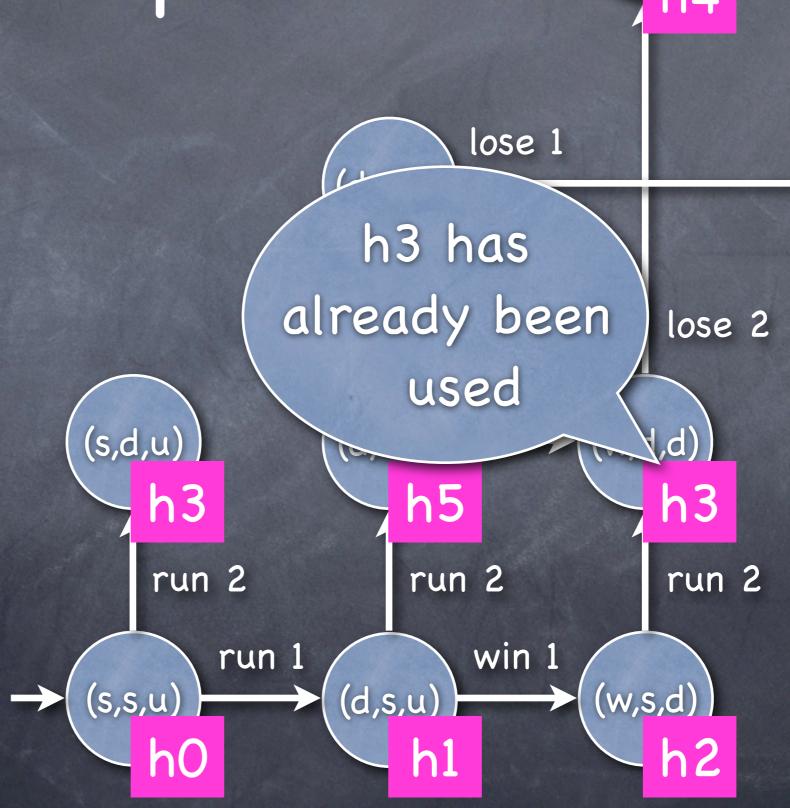


(w,l,d) h4

(l,w,d)

h7

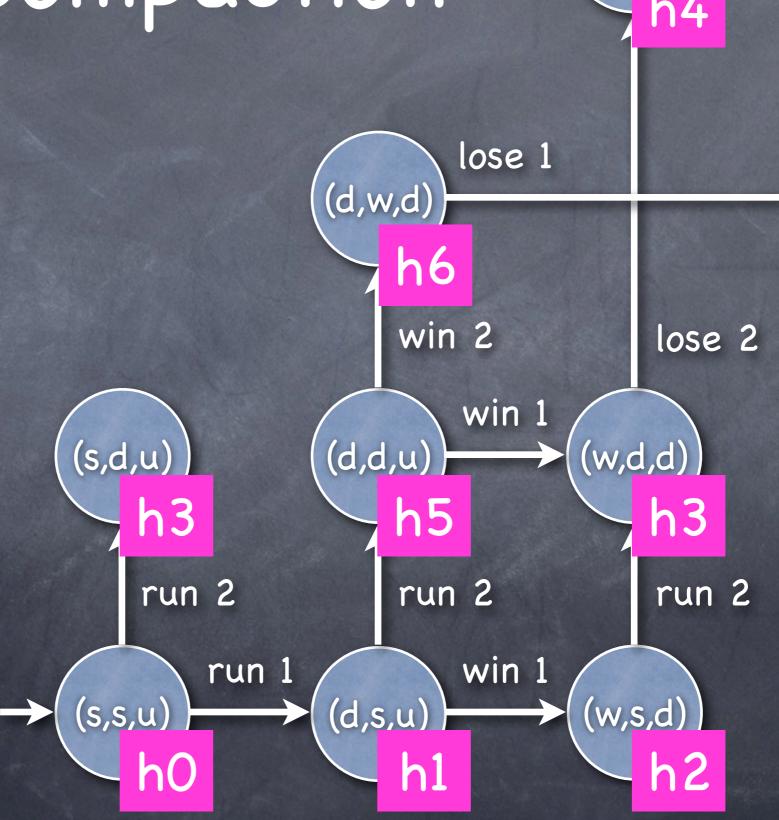
hO h1 h2 h3 h4 h5 h6

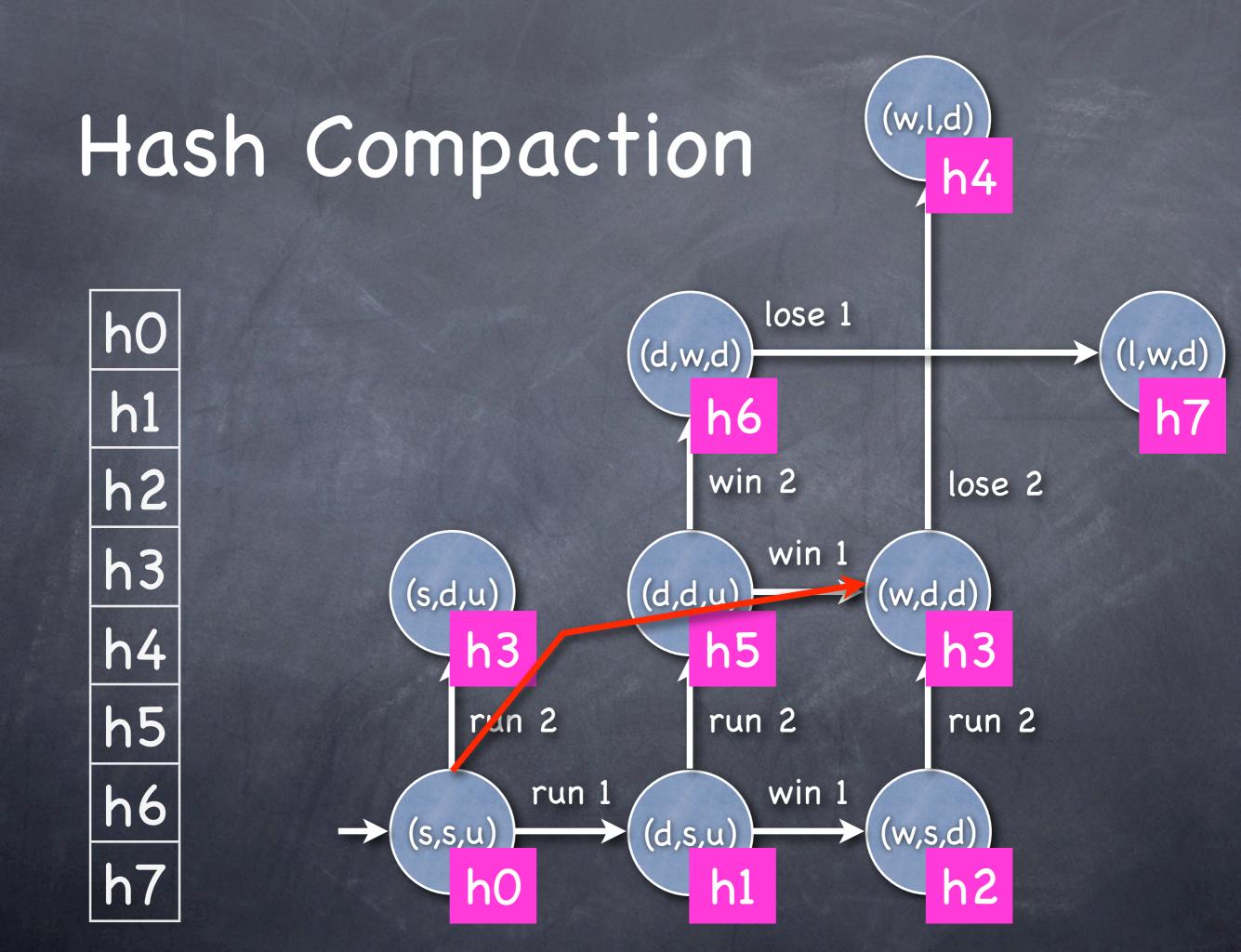


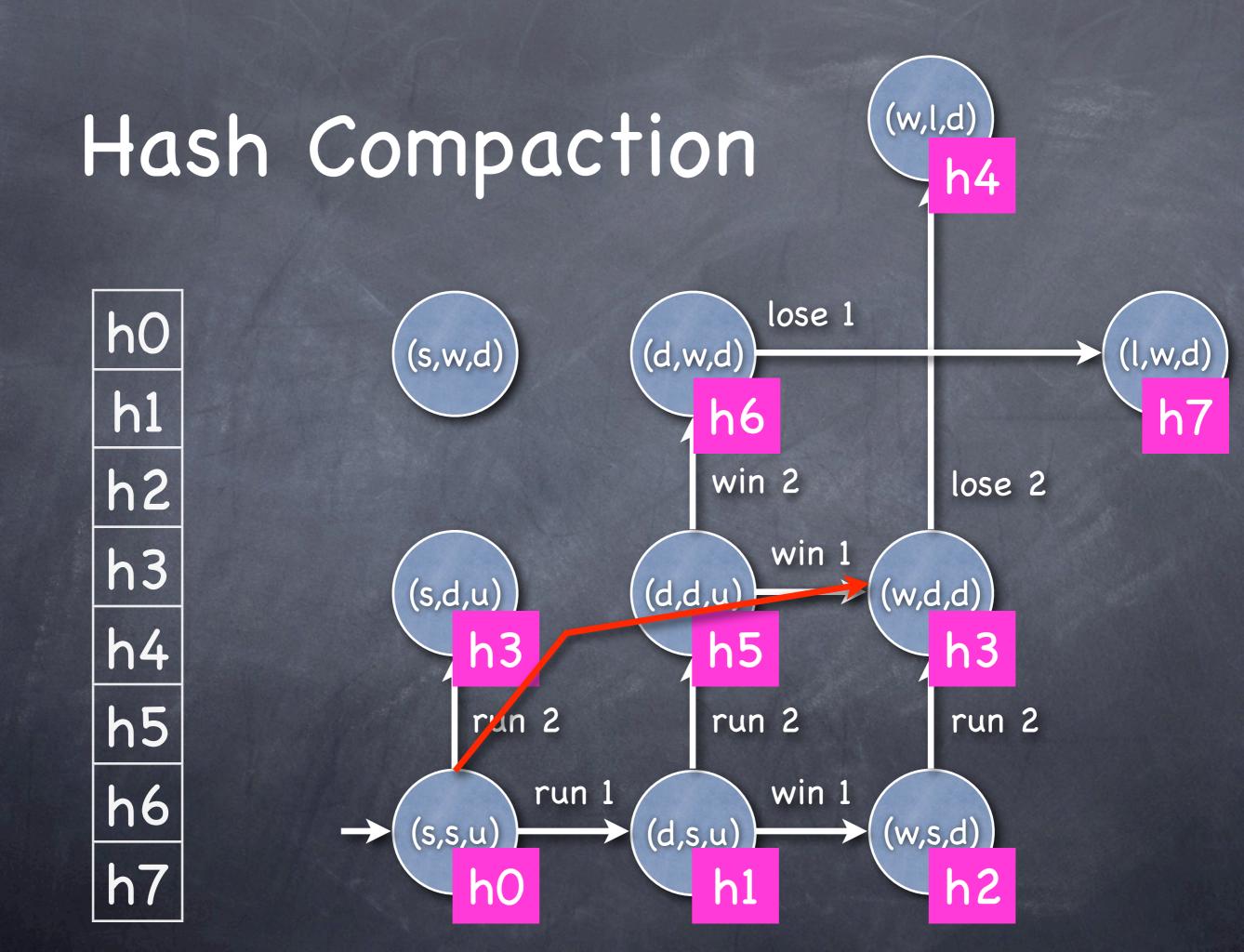
(w,l,d) h4

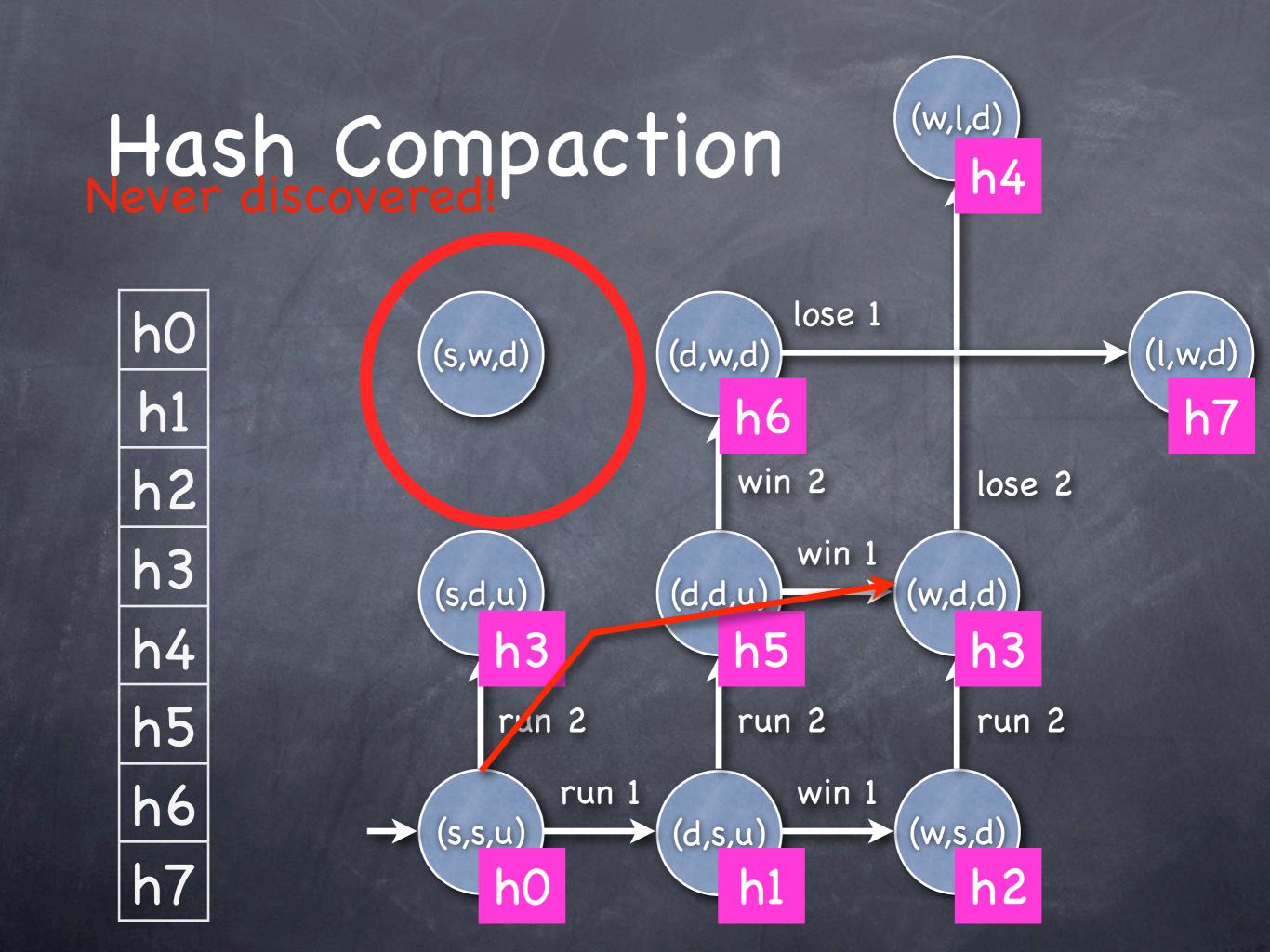
(l,w,d)





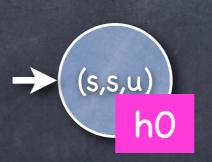




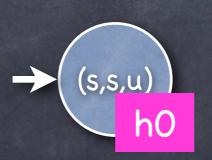




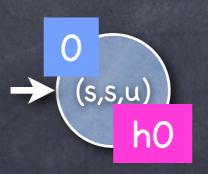




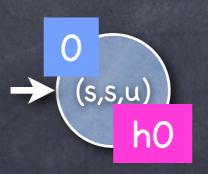






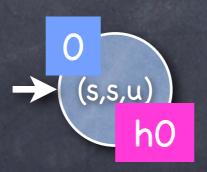






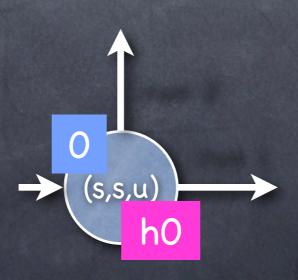




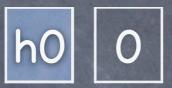


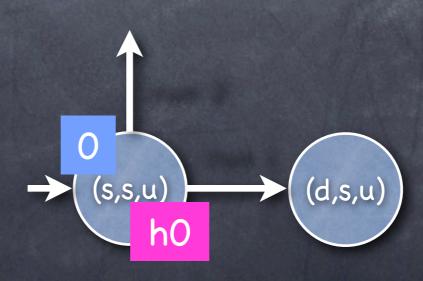






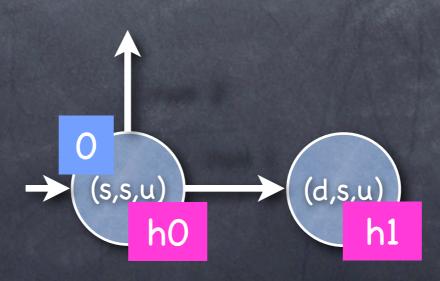




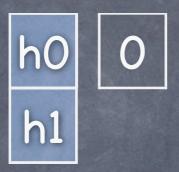


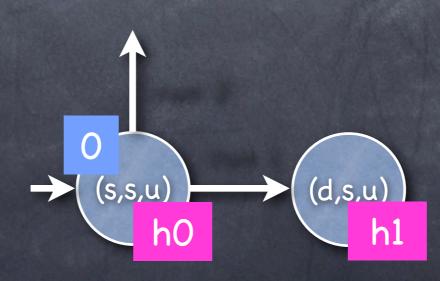




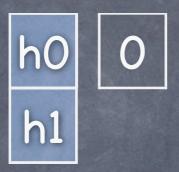


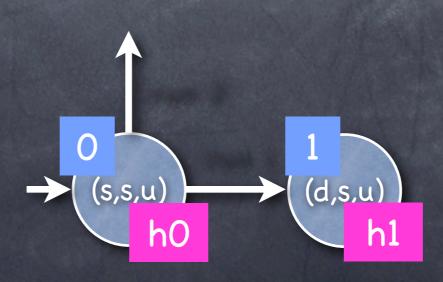




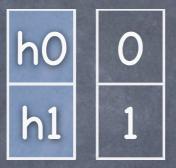


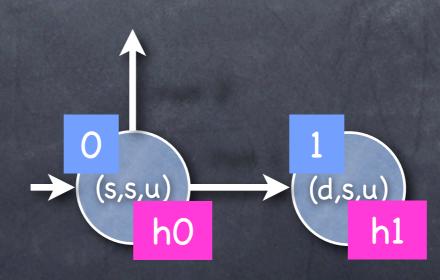




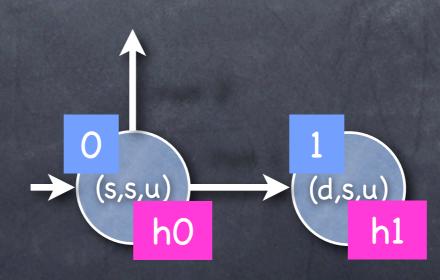


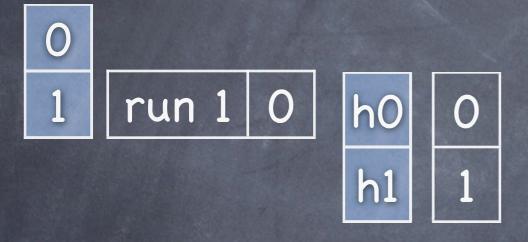


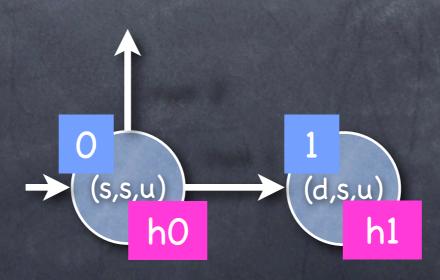


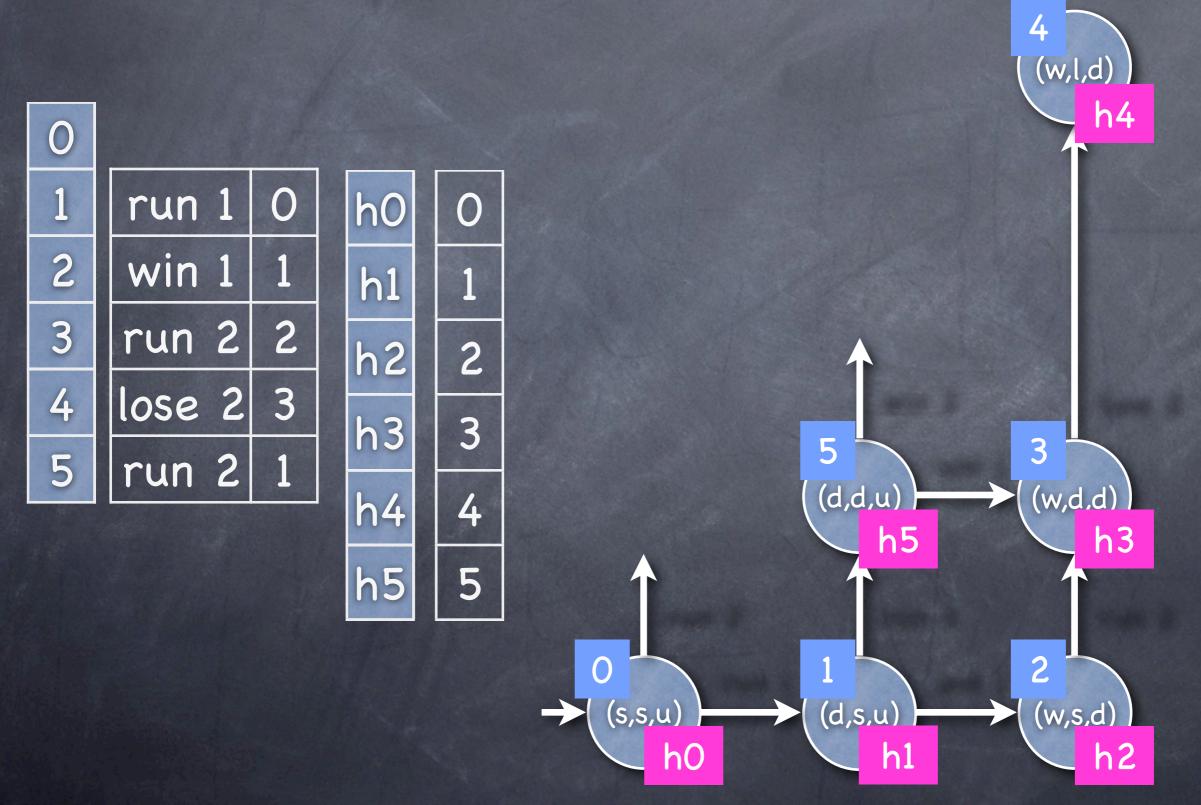




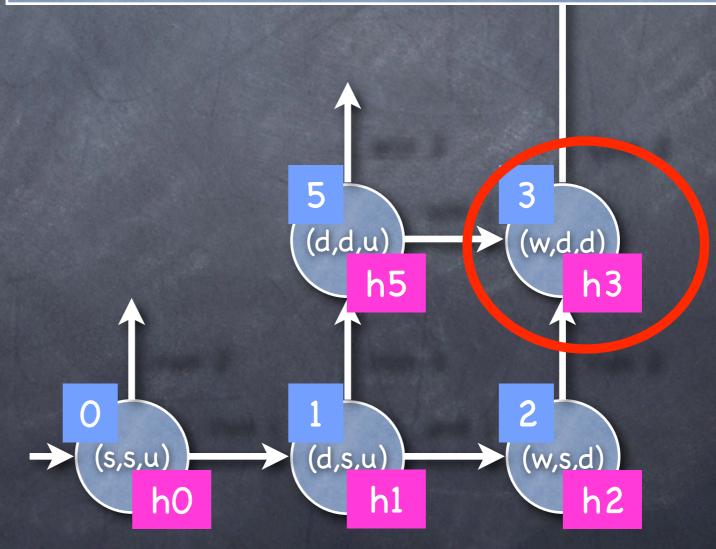




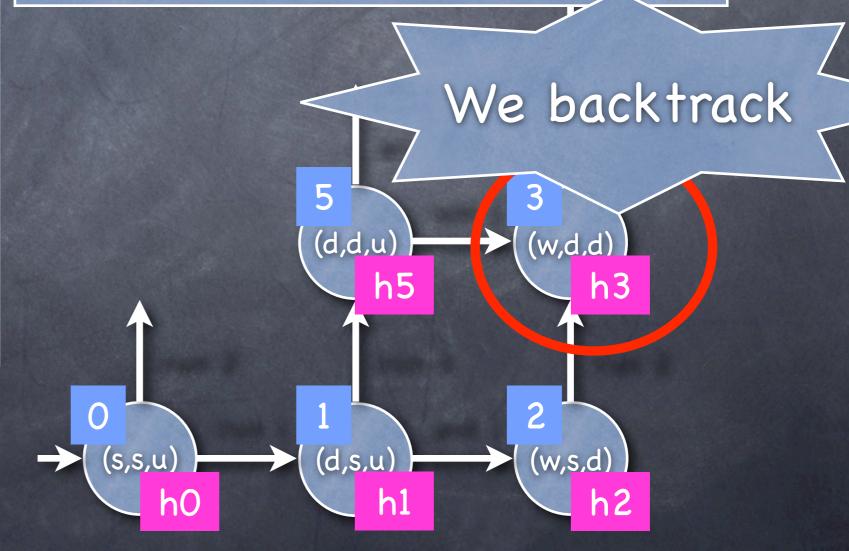




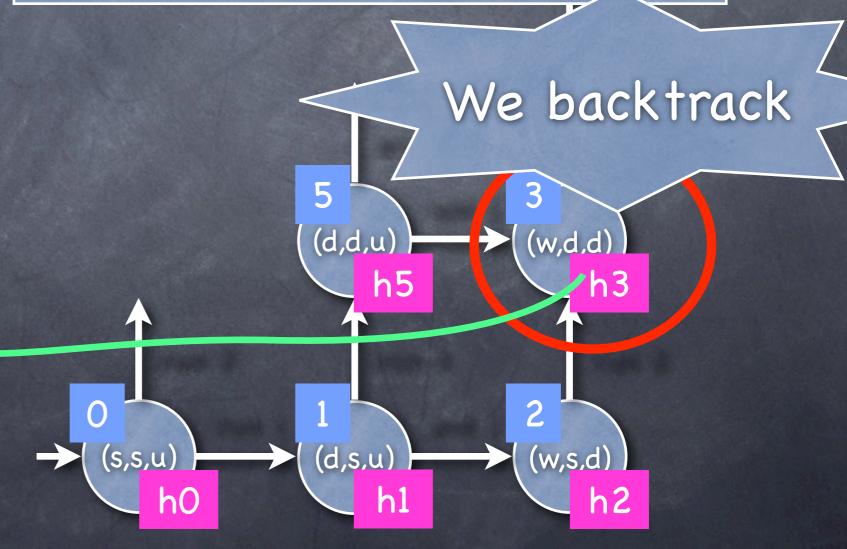
h0		0
h1		1
h2		2
h3		3
h4		4
h5	The second	5



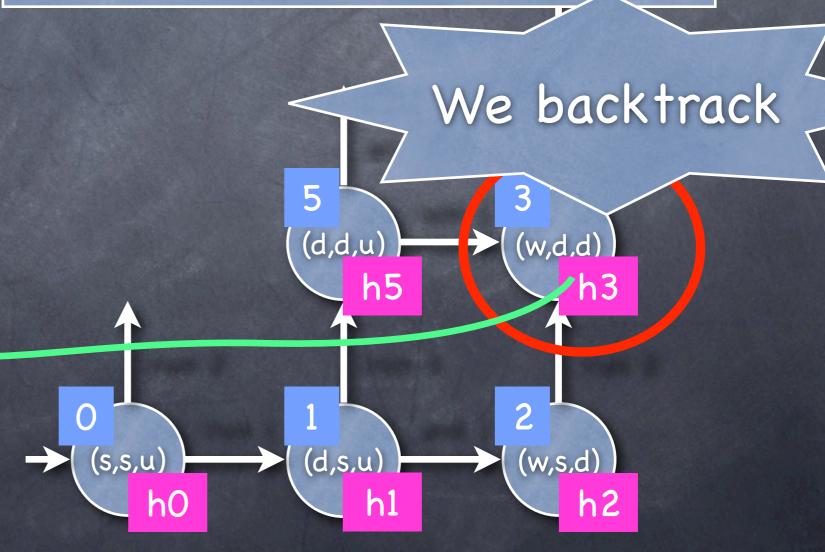
h0		0
h1		1
h2		2
h3		3
h4		4
h5	THE RESERVE	5



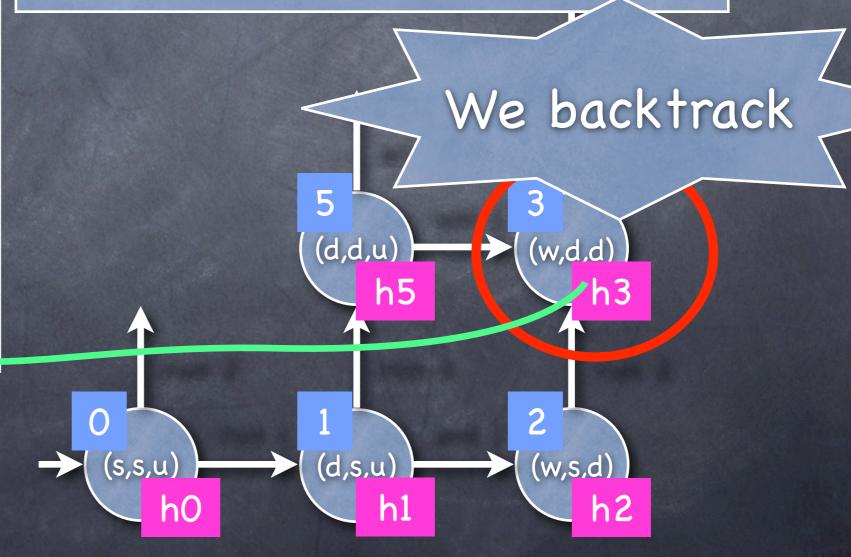
h0 h1 h2 h3 3 h4



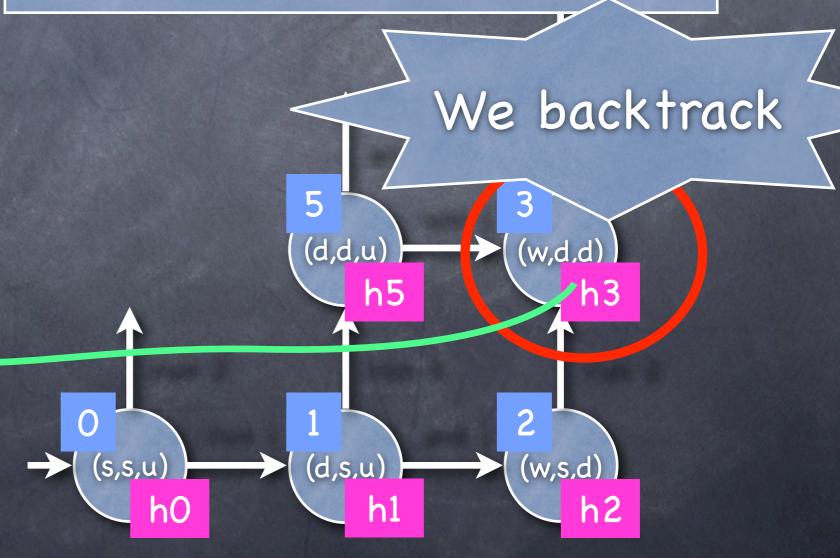
0 run 1 h0 win 1 h1 3 run 2 lose 2 ns run 2 h4



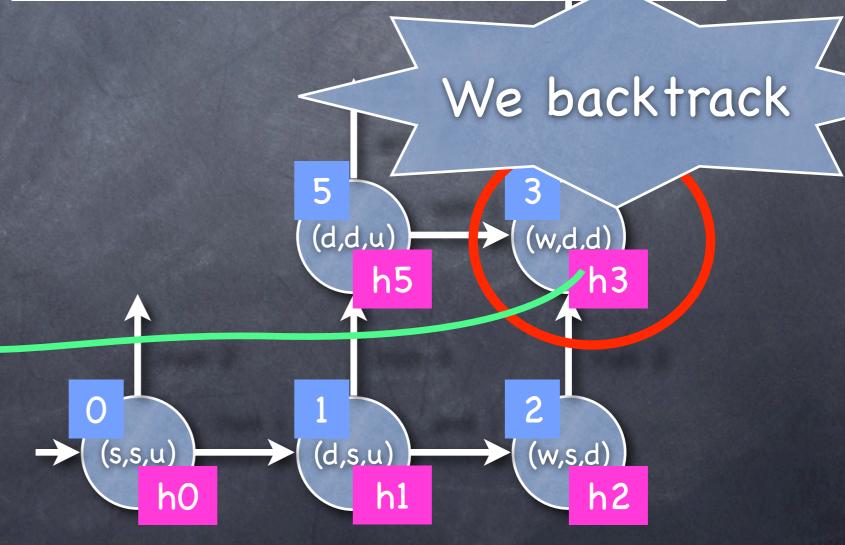
0 run h0 win 1 h1 lose ns run 2 h4

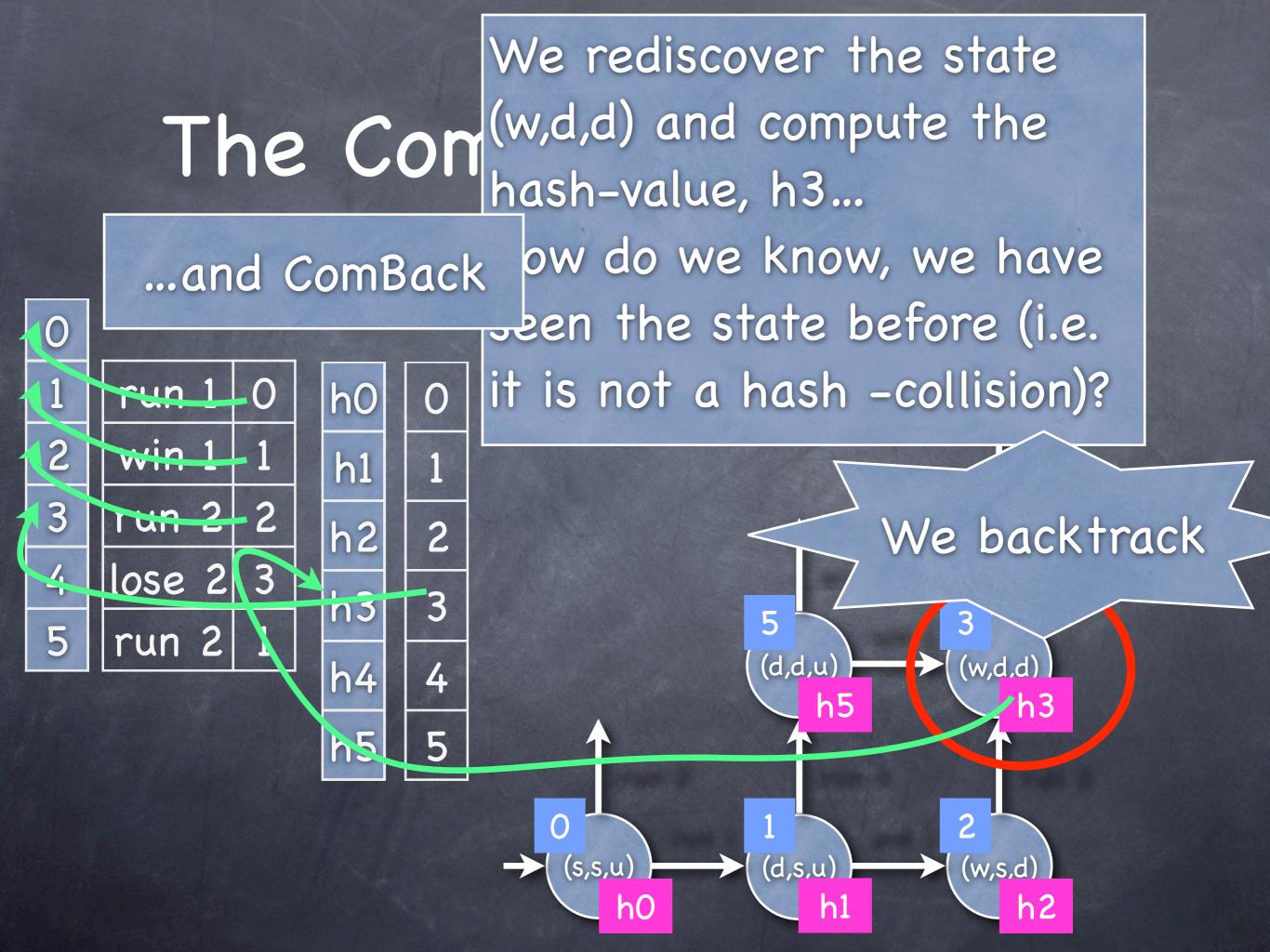


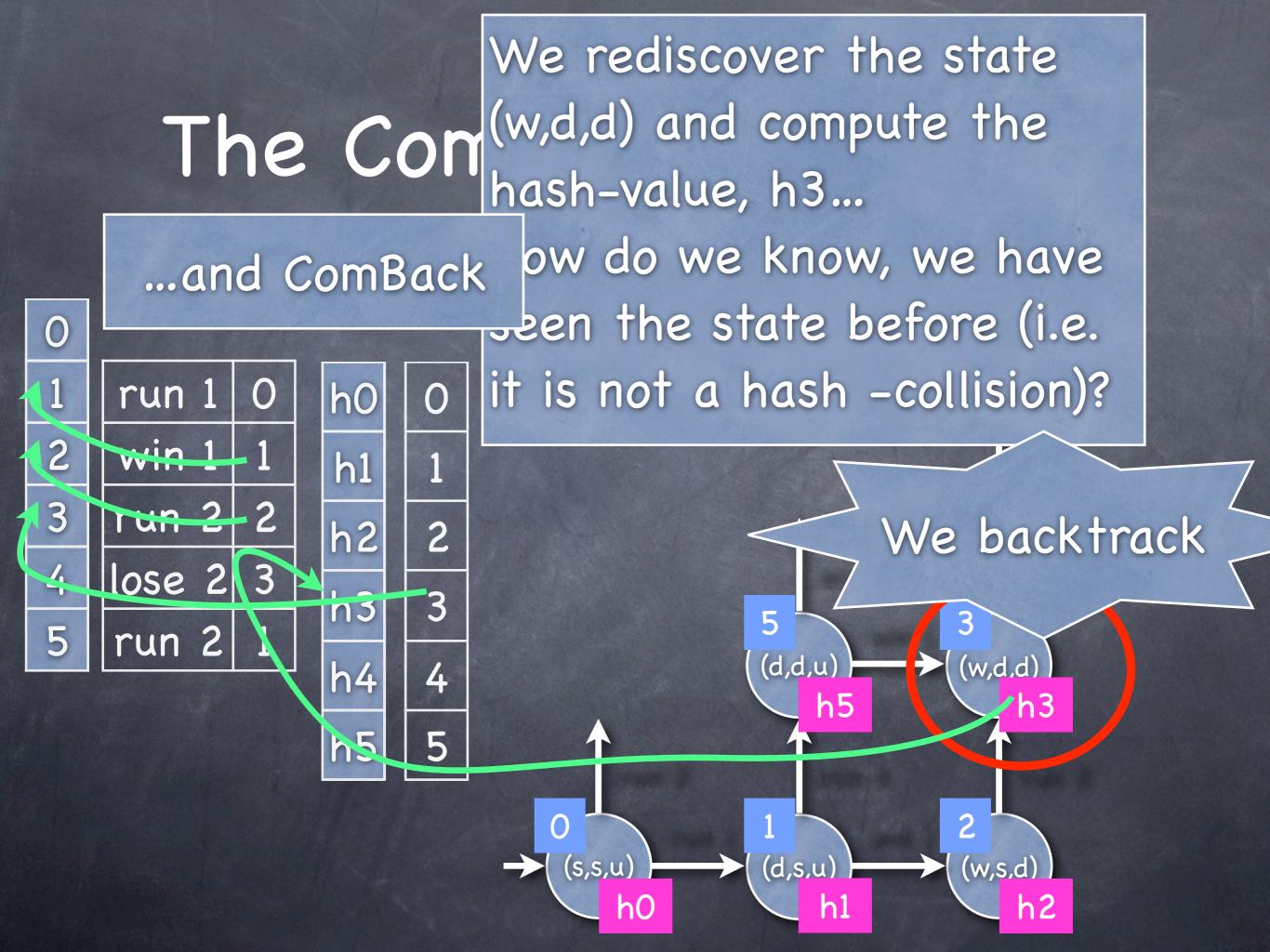
0 h0 h1 ns run 2 h4

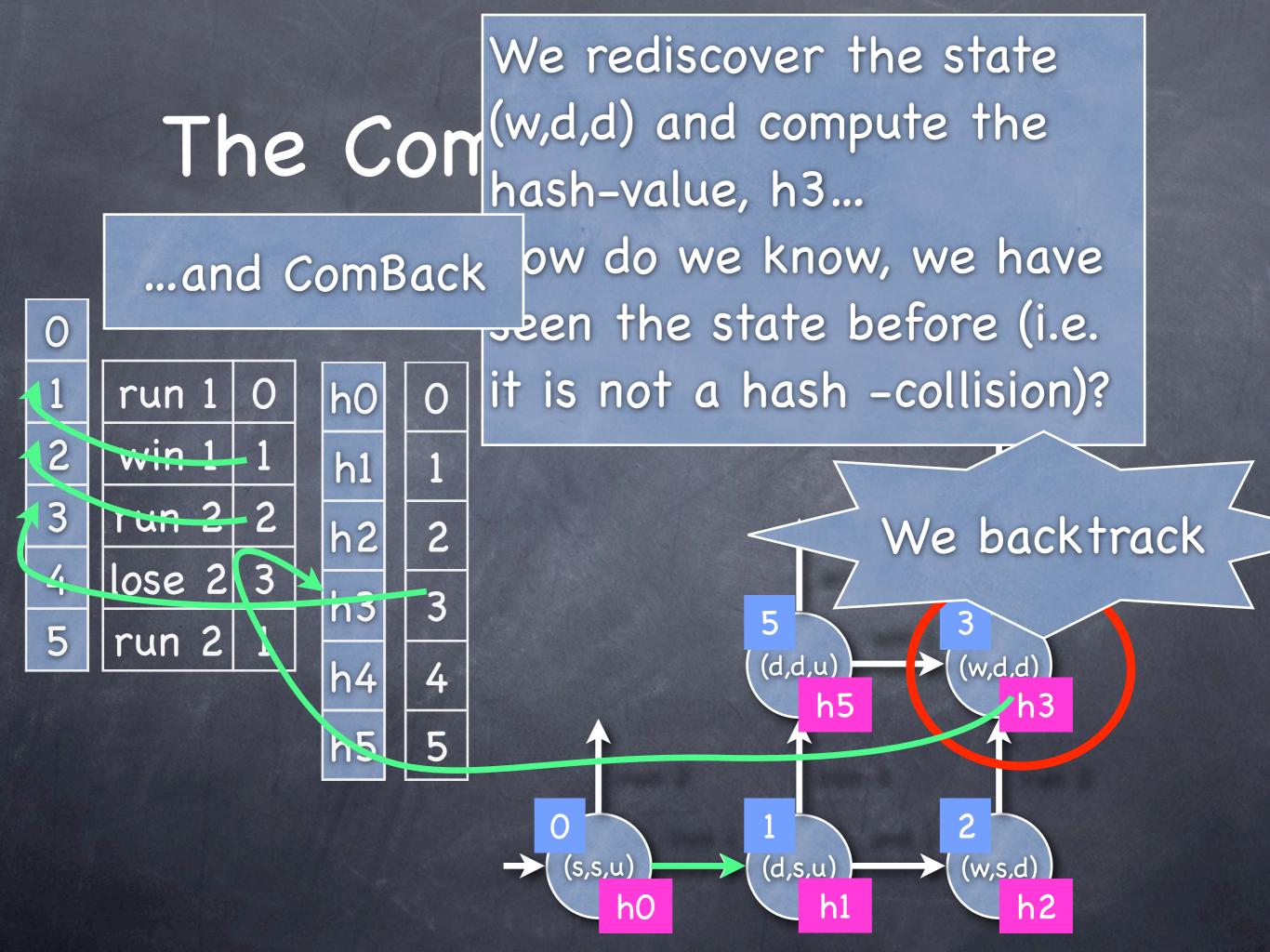


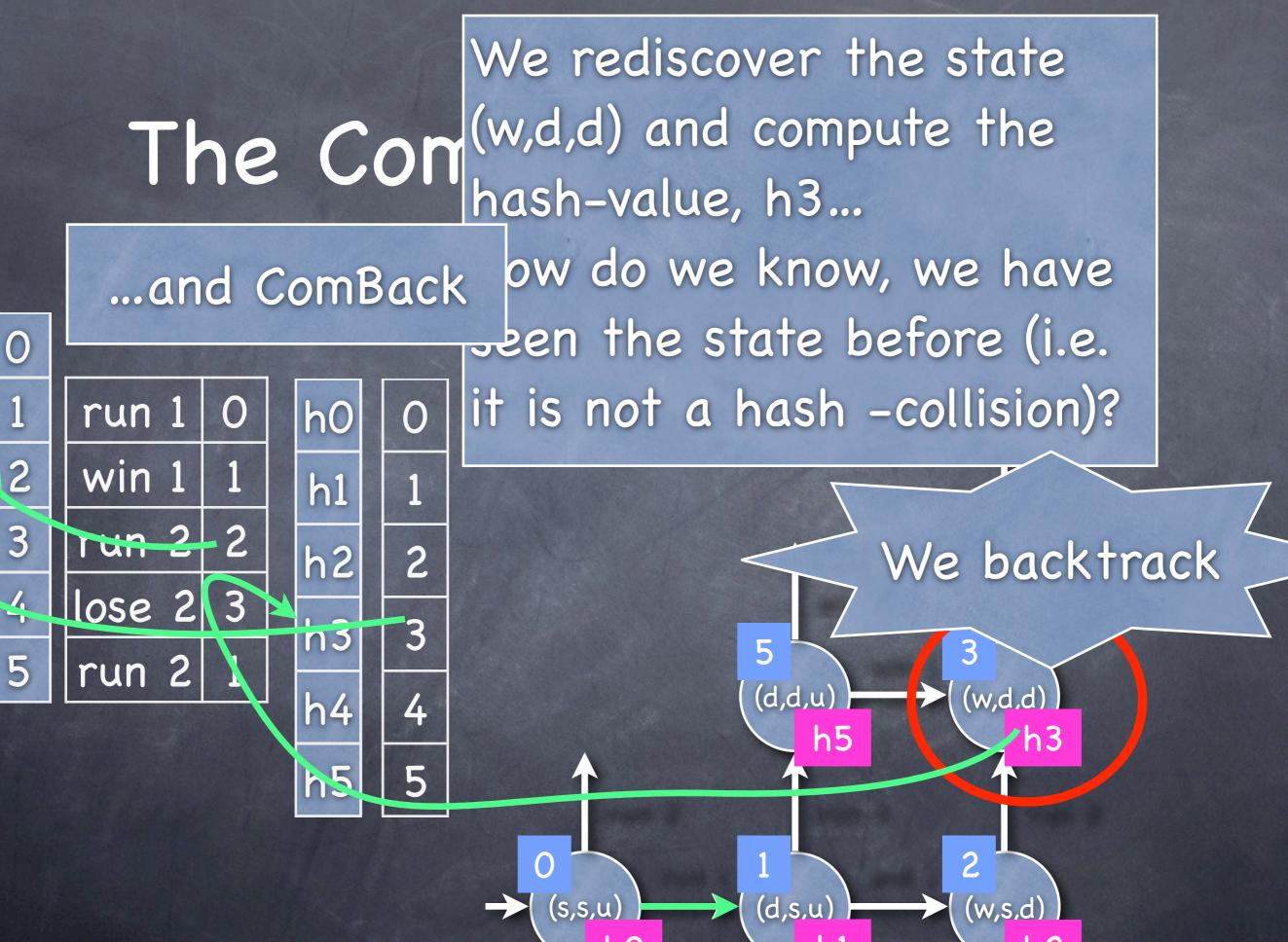
h0 h1 n3 run 2 h4













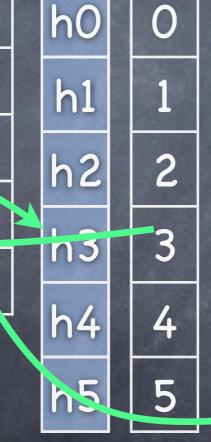
ow do we know, we have ...and ComBack

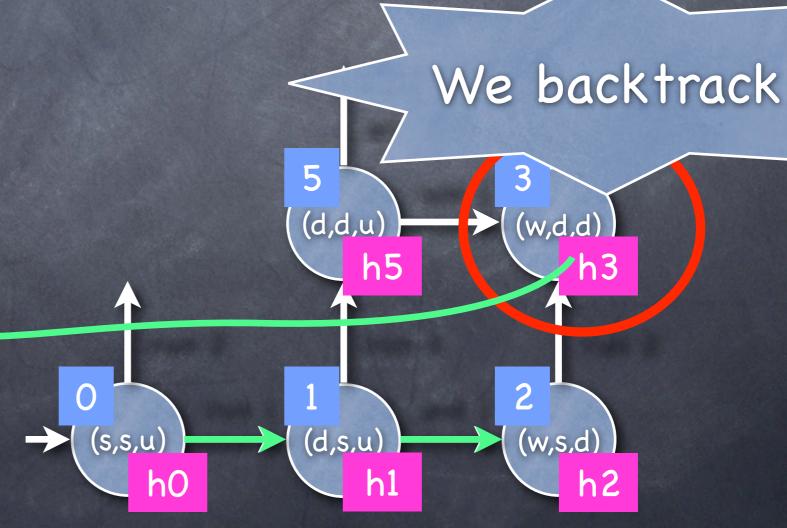
Jeen the state before (i.e.

it is not a hash -collision)?



0



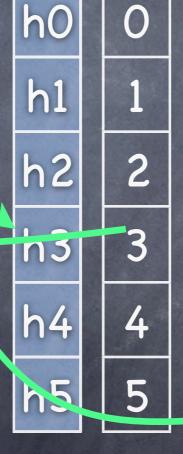


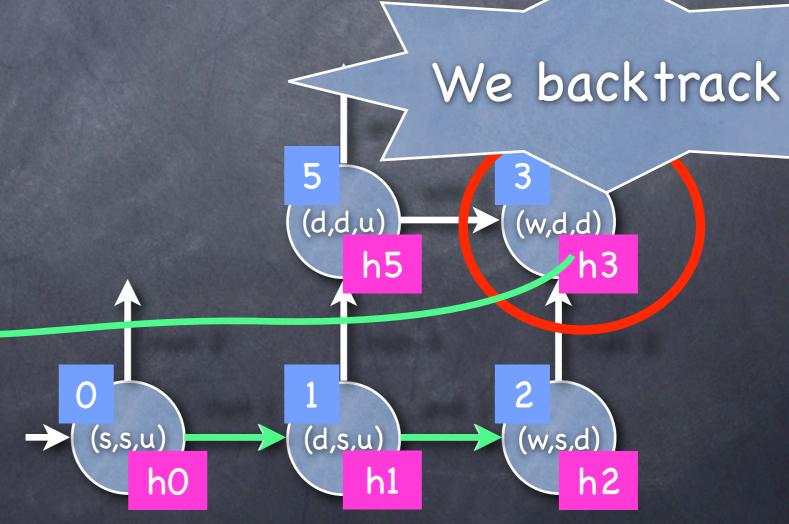


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it is not a hash -collision)?







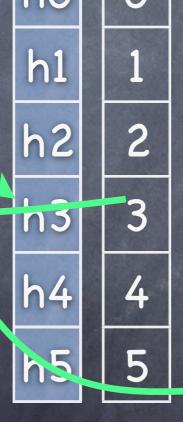


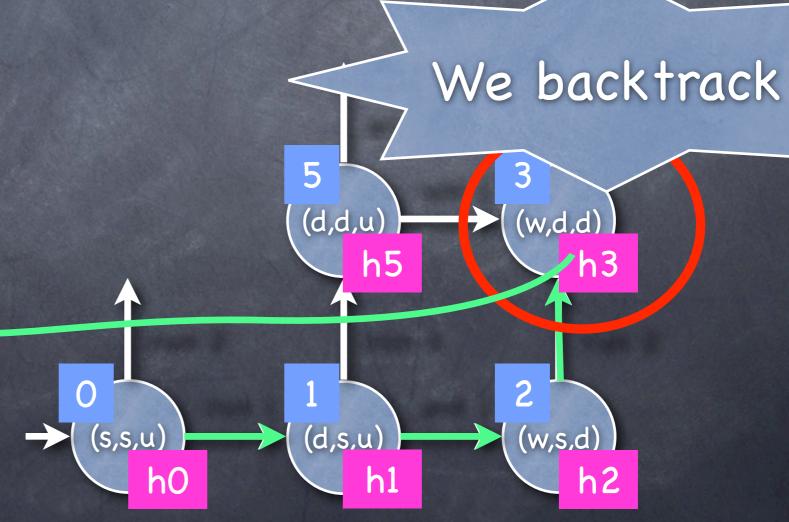
ow do we know, we have ...and ComBack Jeen the state before (i.e.

it is not a hash -collision)?



3



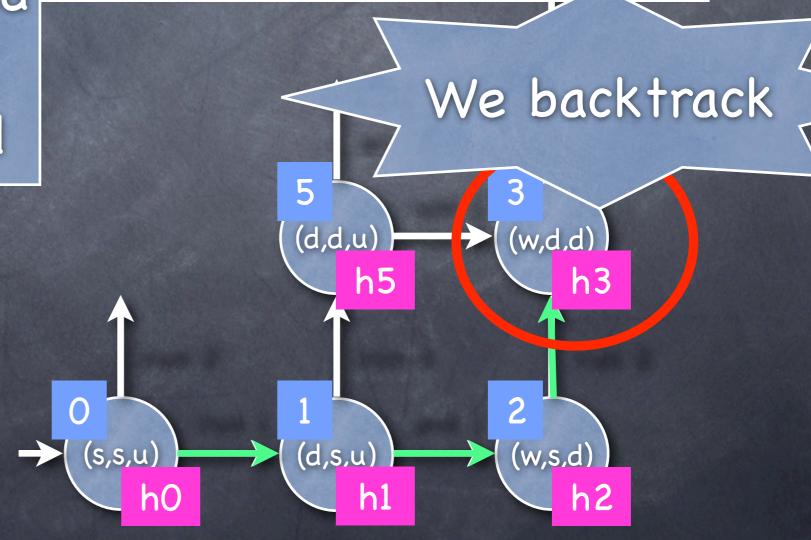


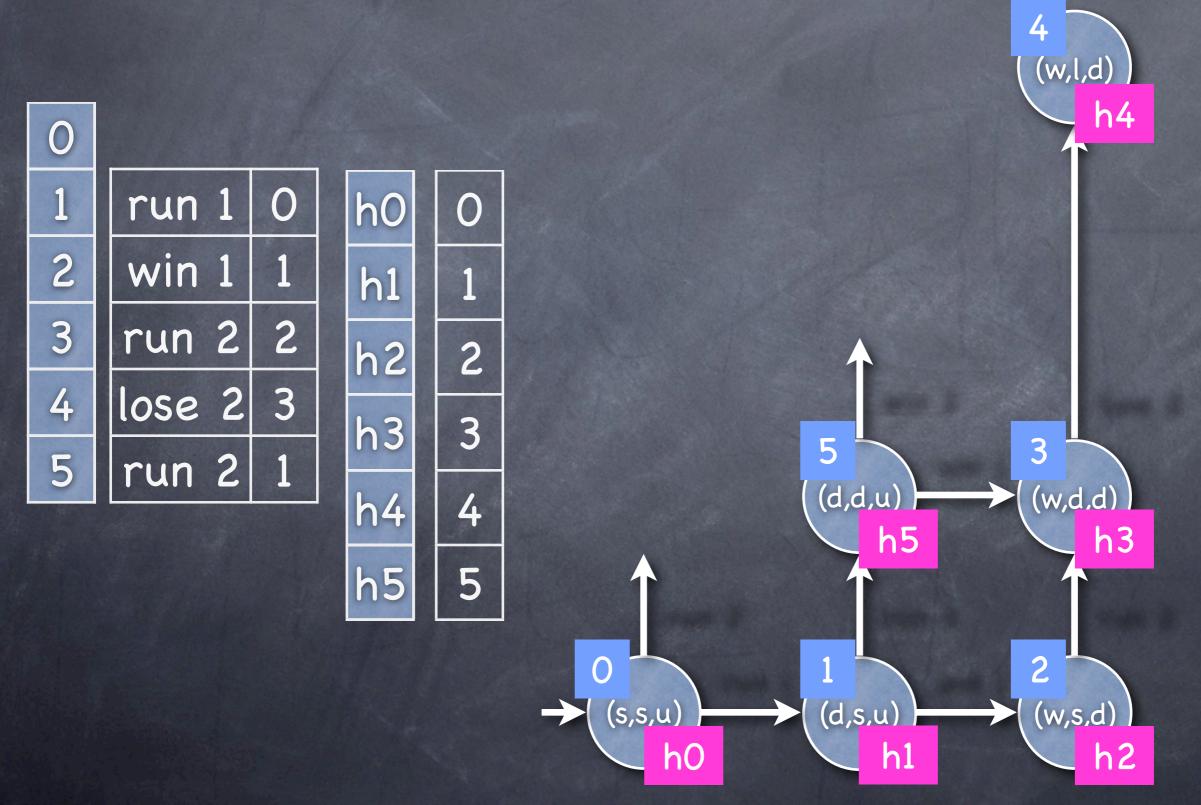
...and ComBack

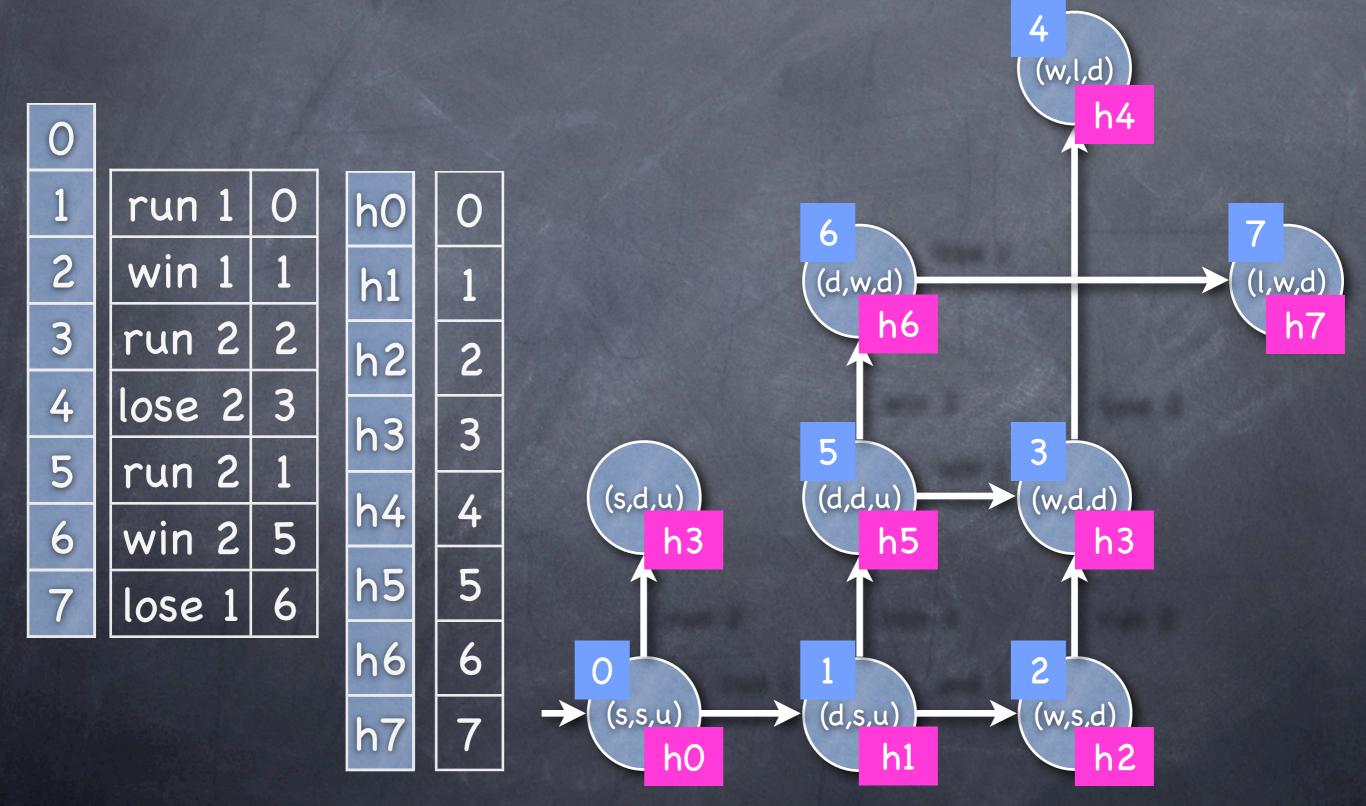
ow do we know, we have en the state before (i.e. is not a hash -collision)?

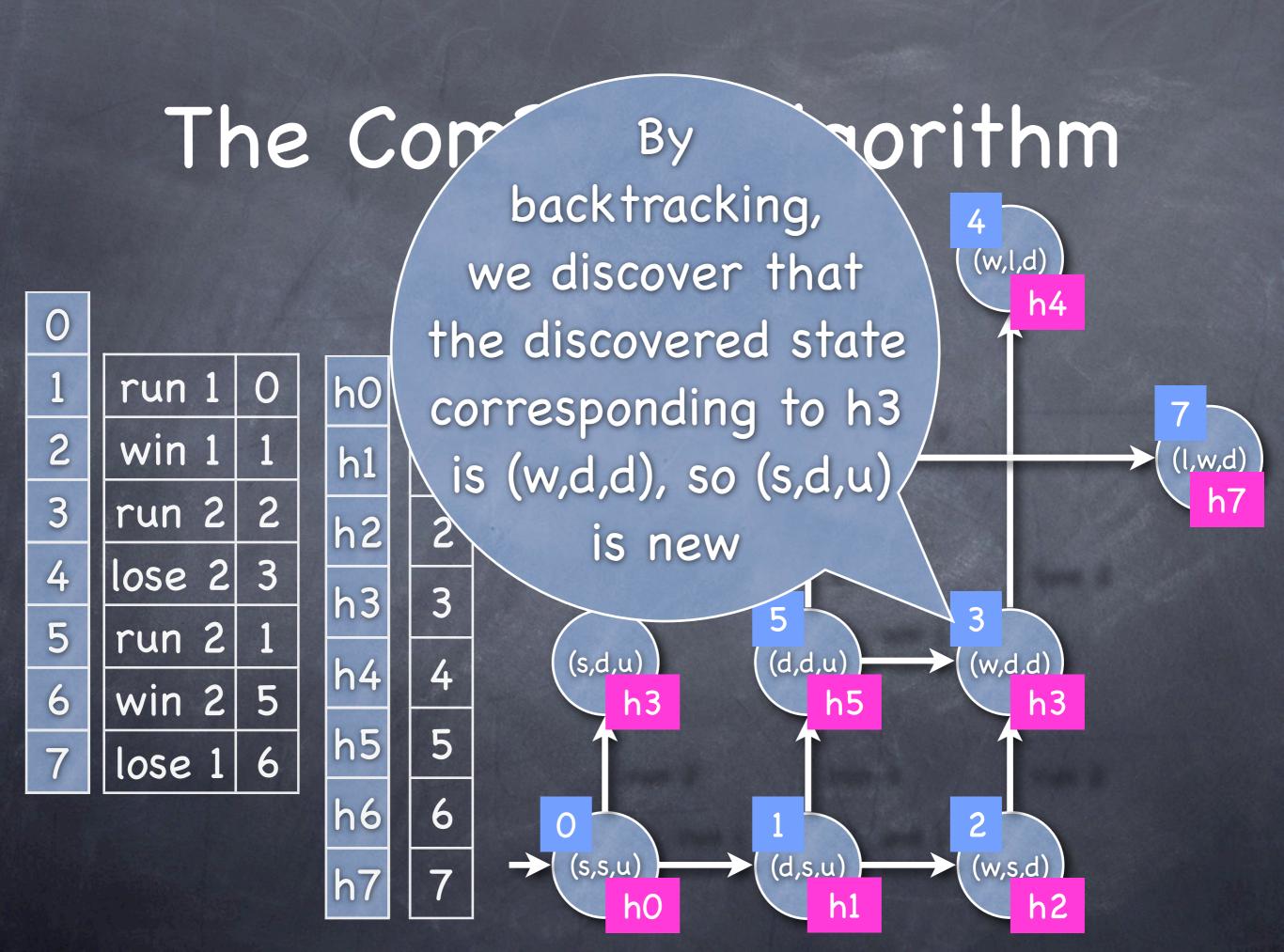
... and notice we have (re-)arrived at (w,d,d). No need to proceed

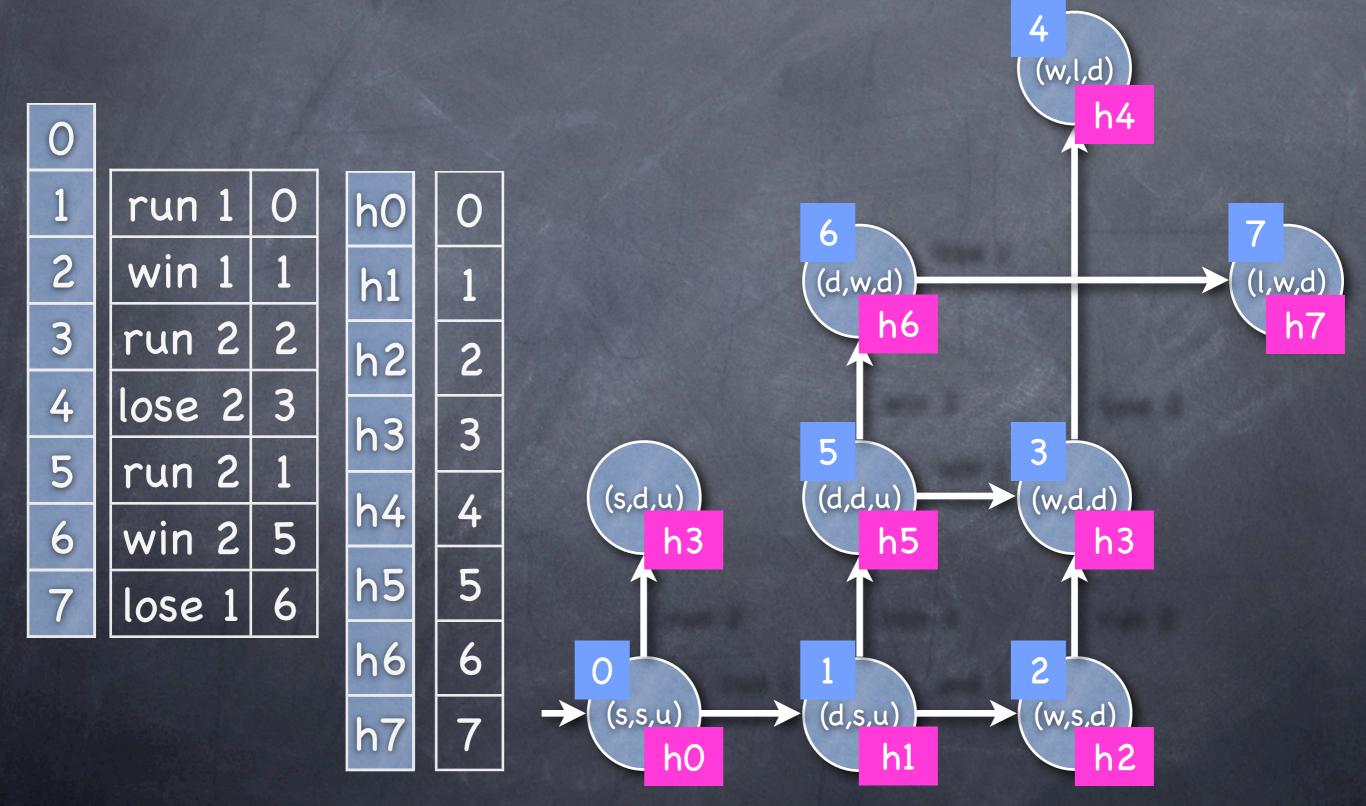
3

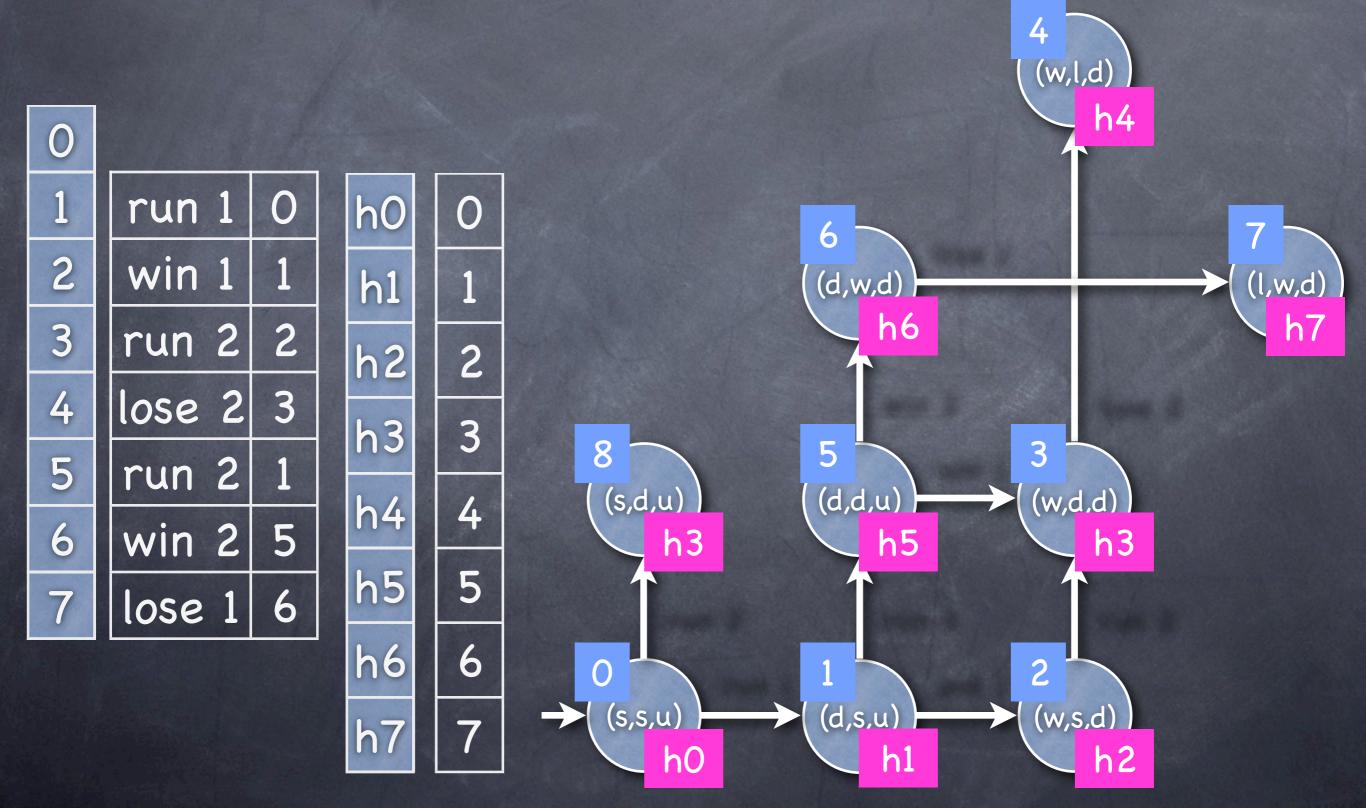


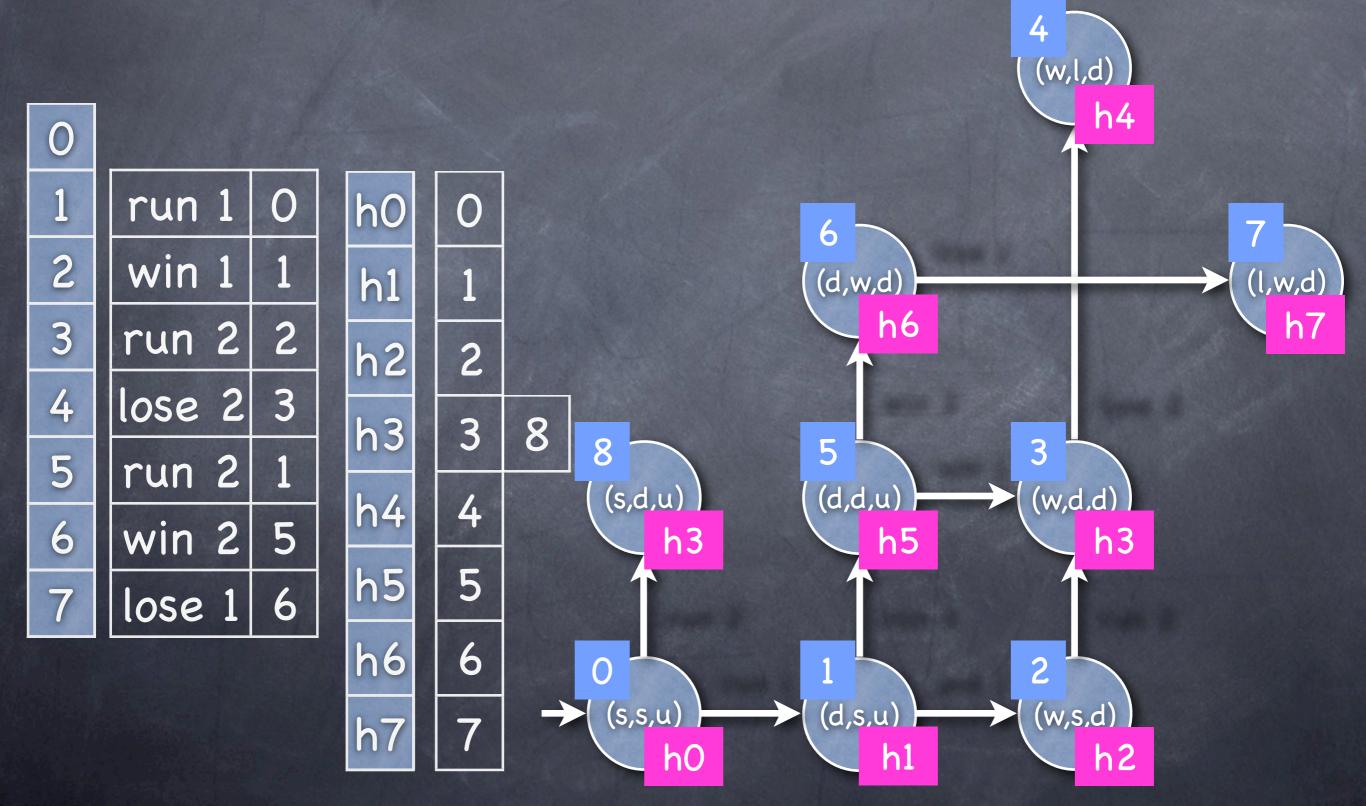


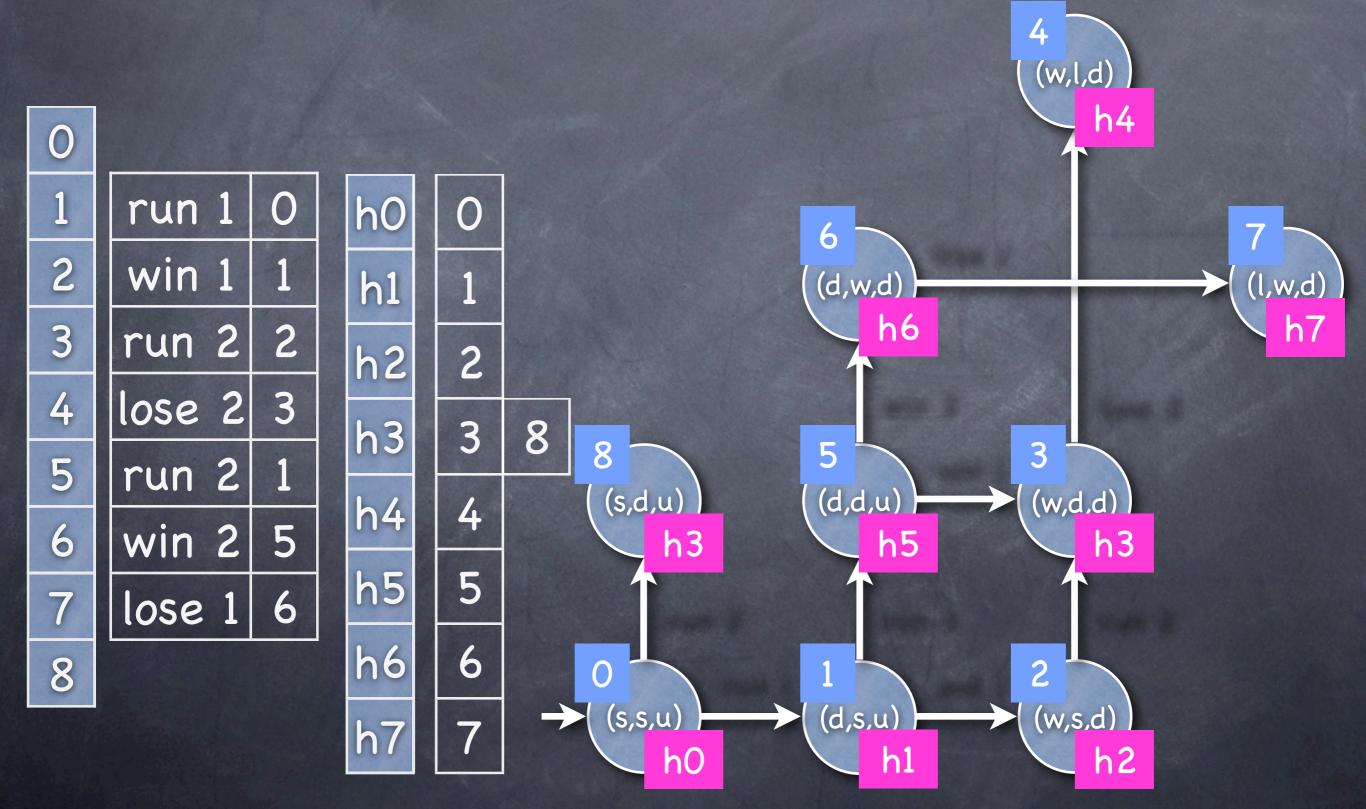


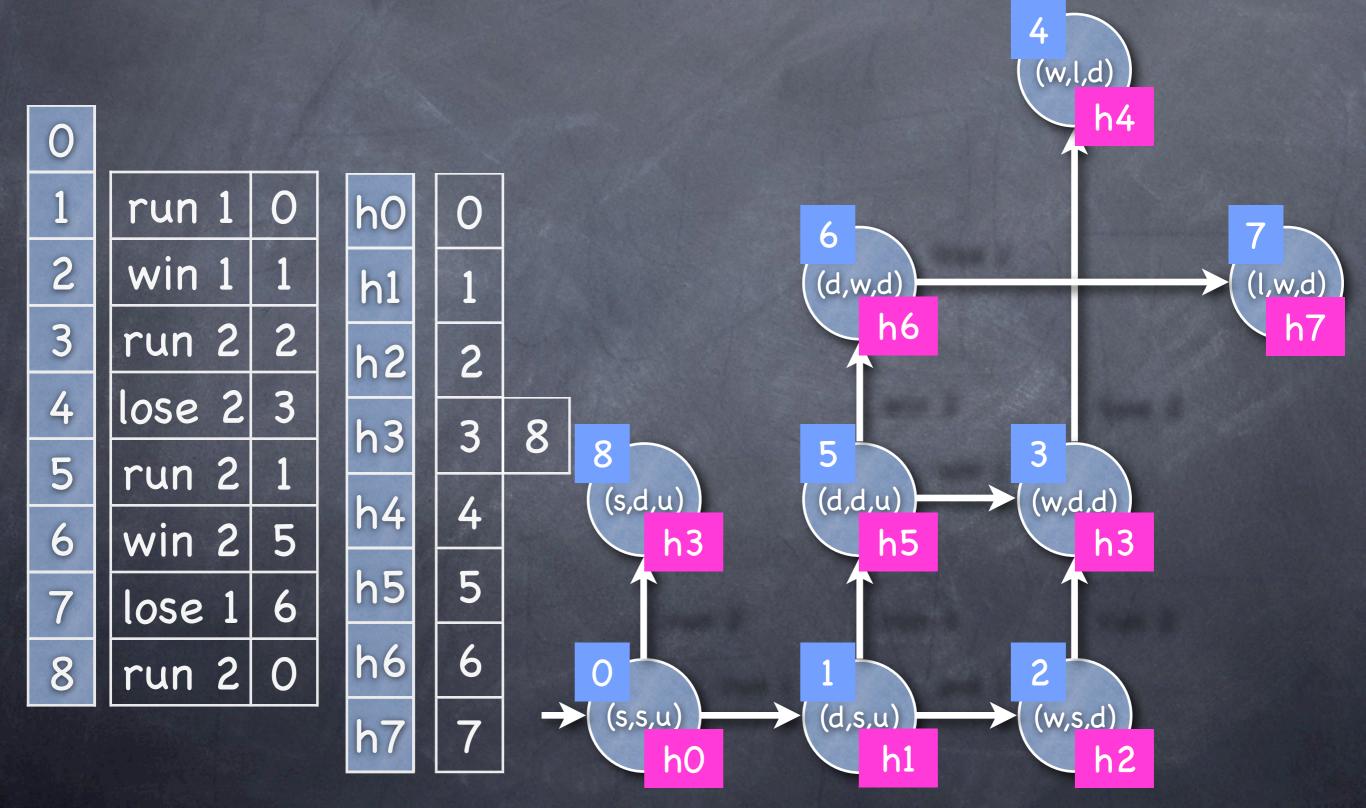


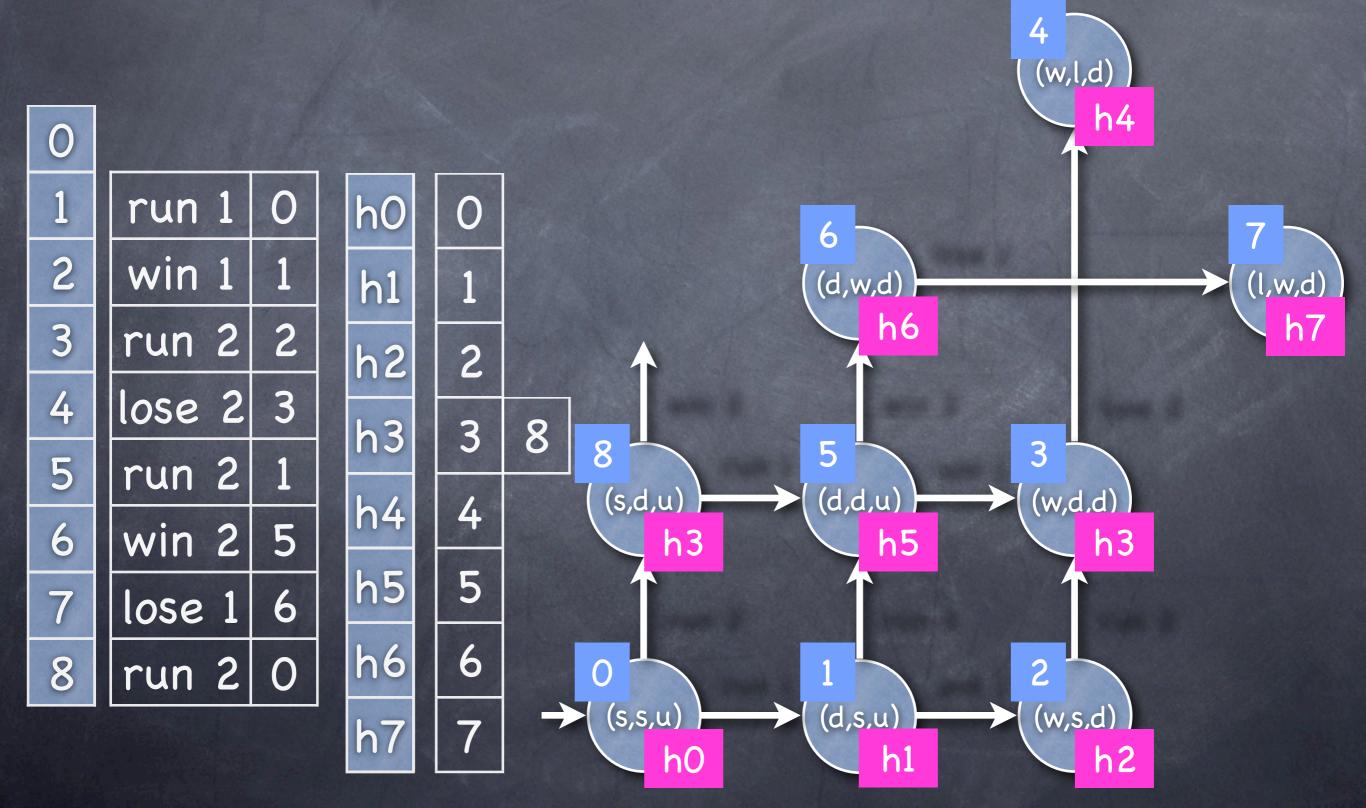


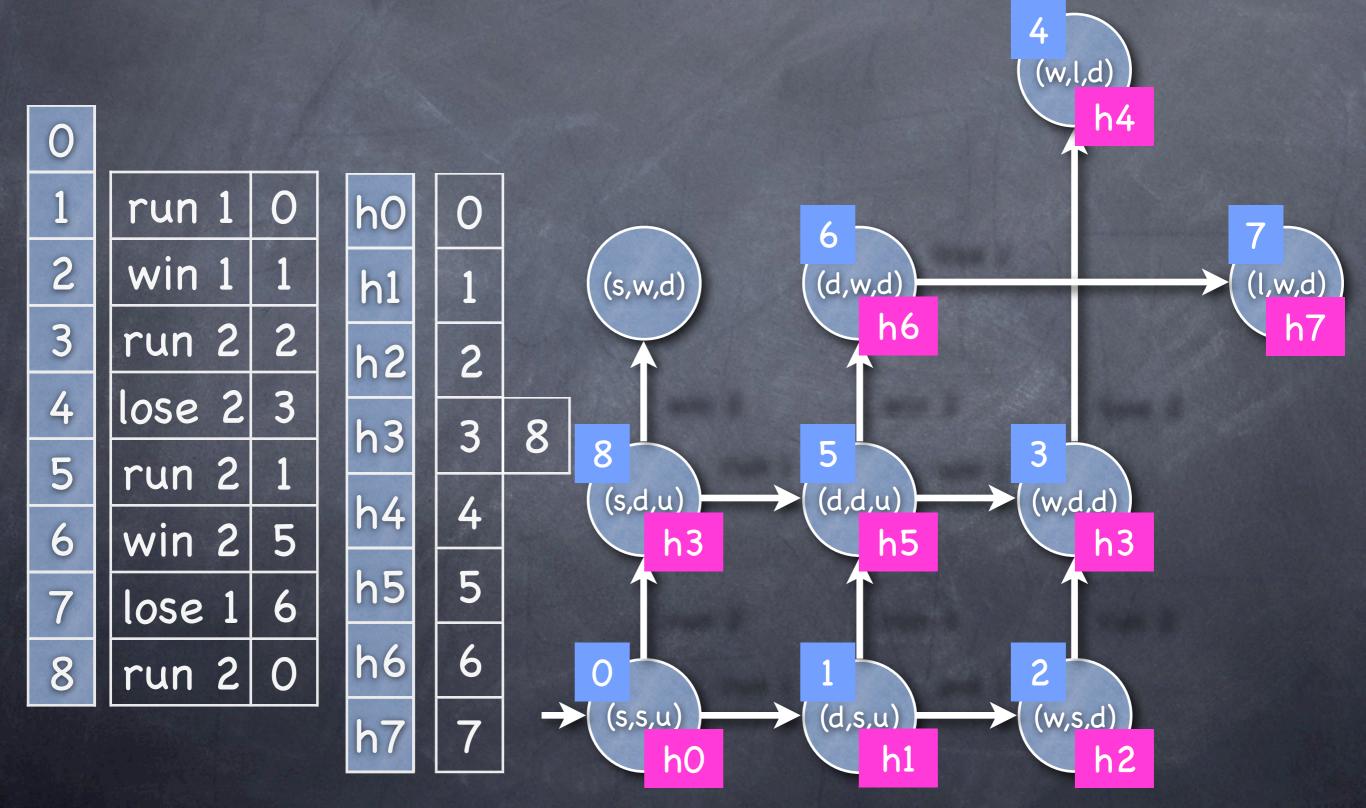


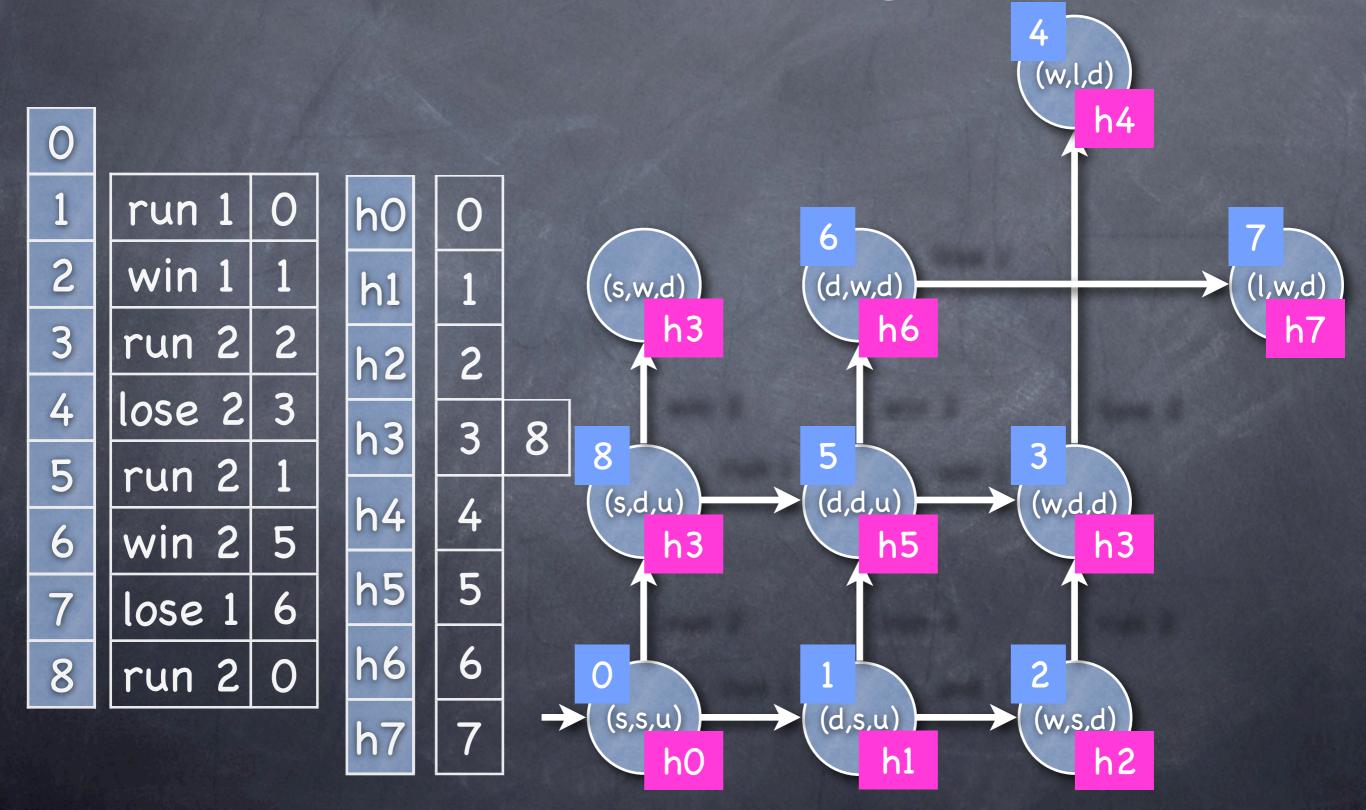


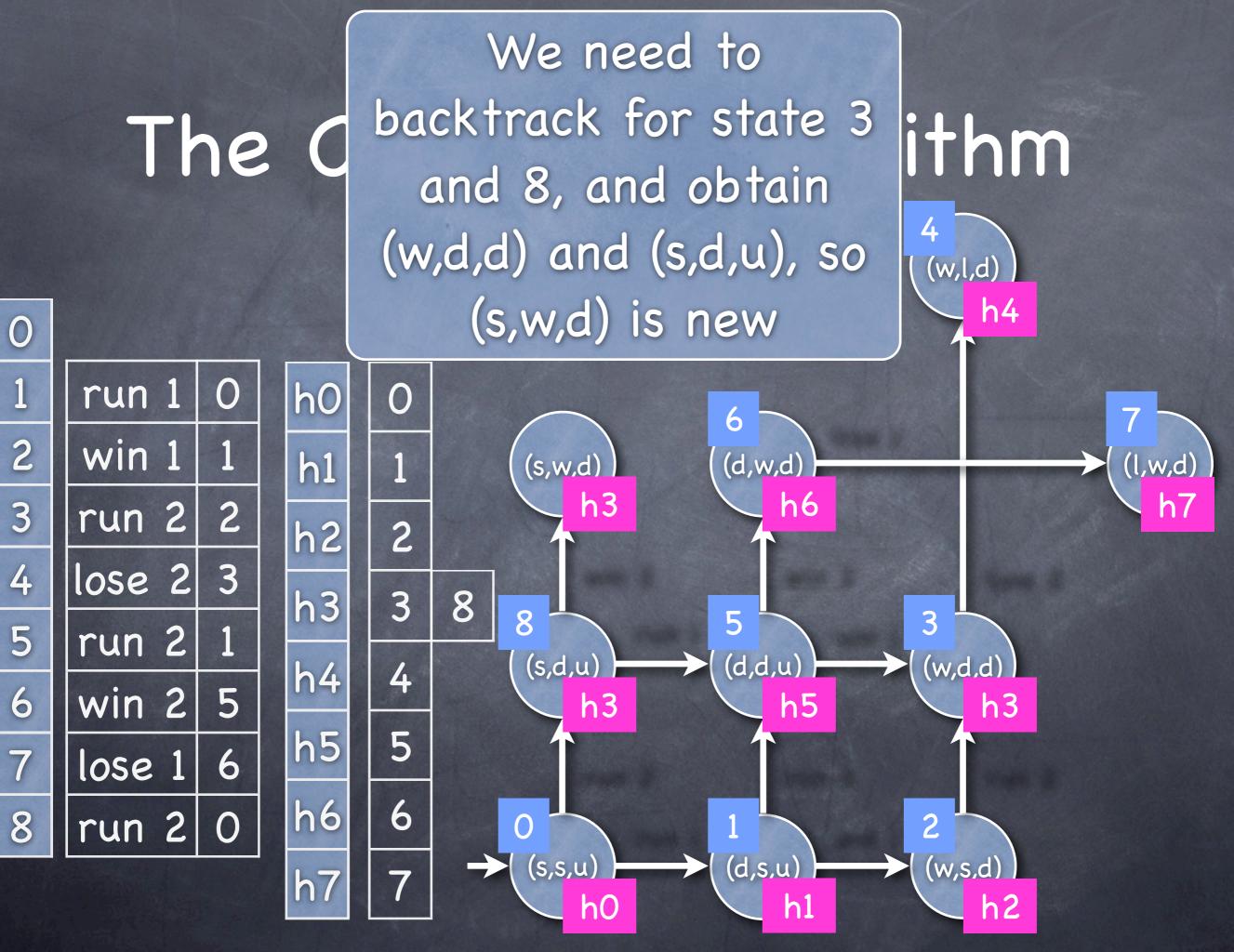


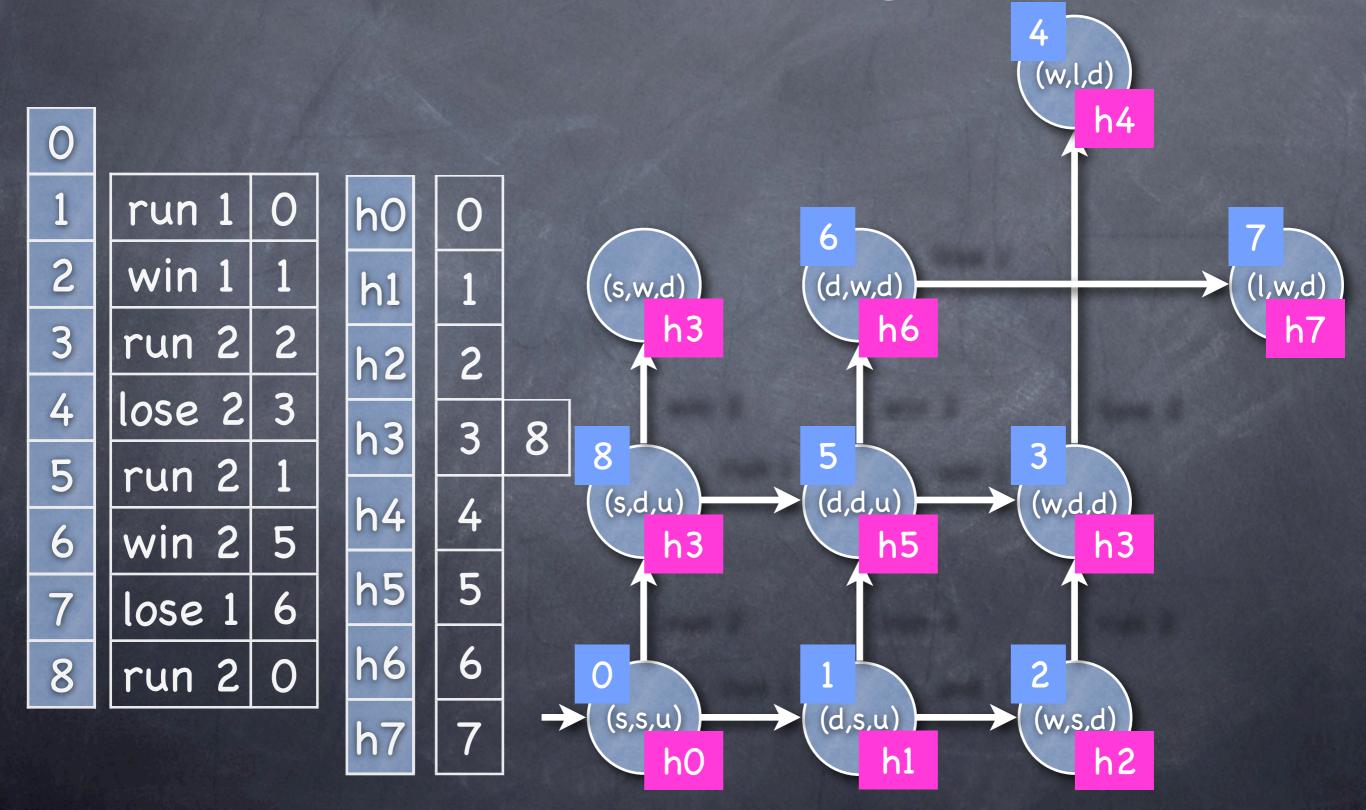


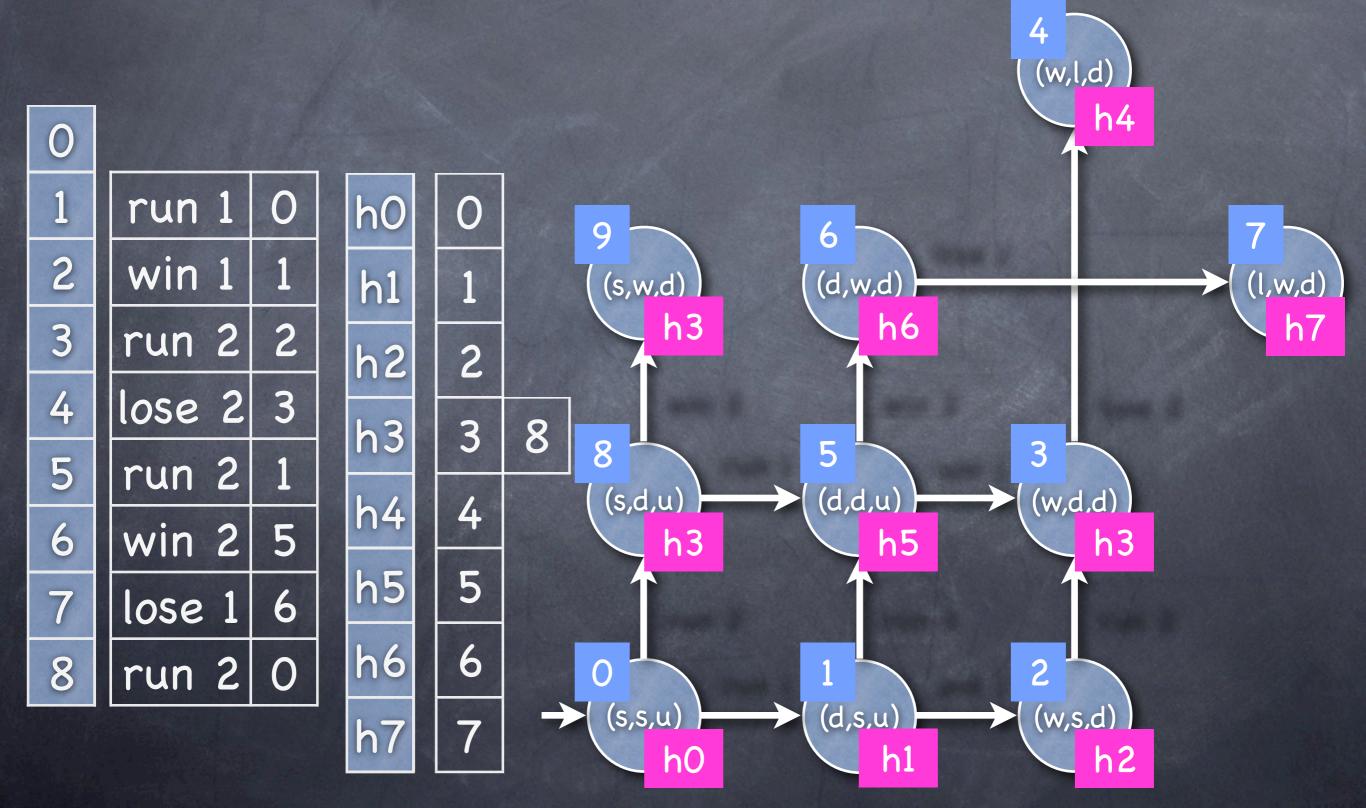


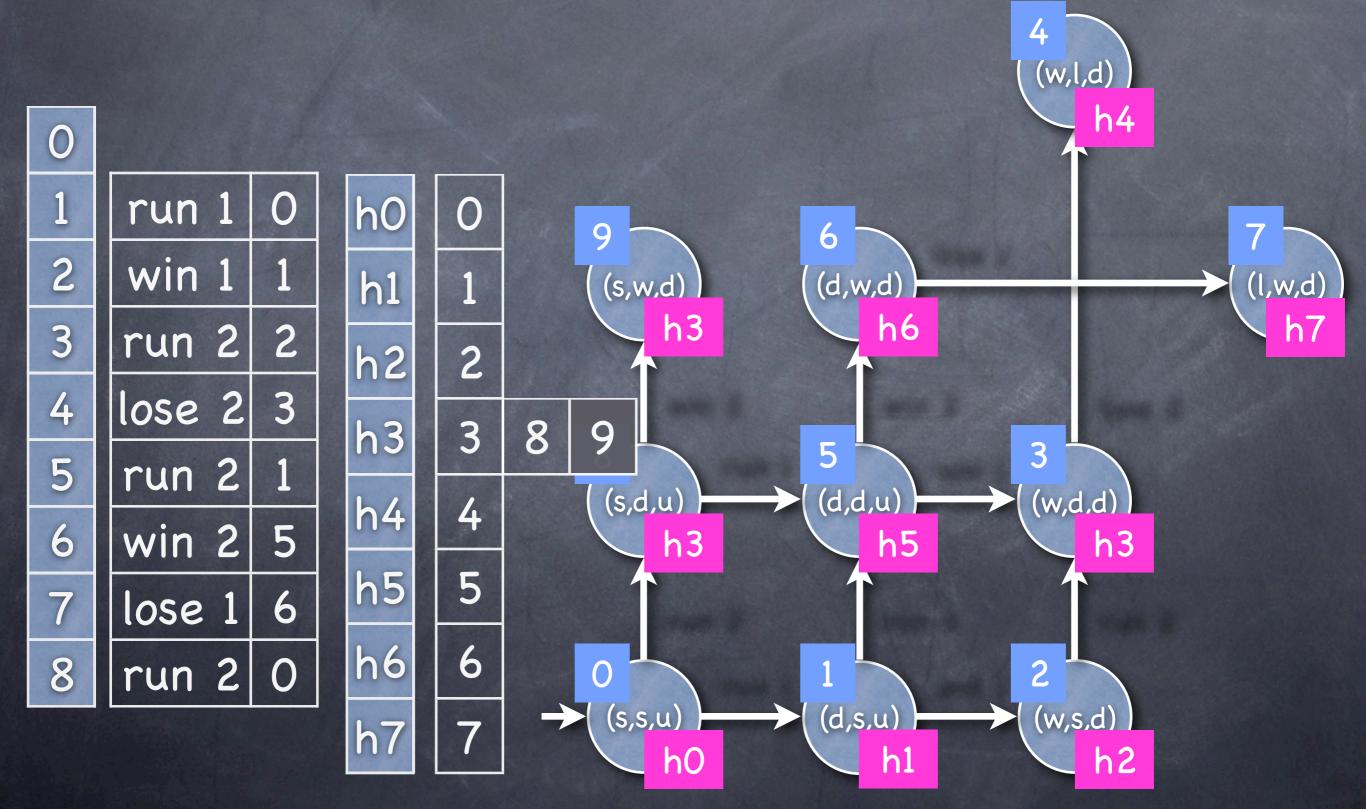


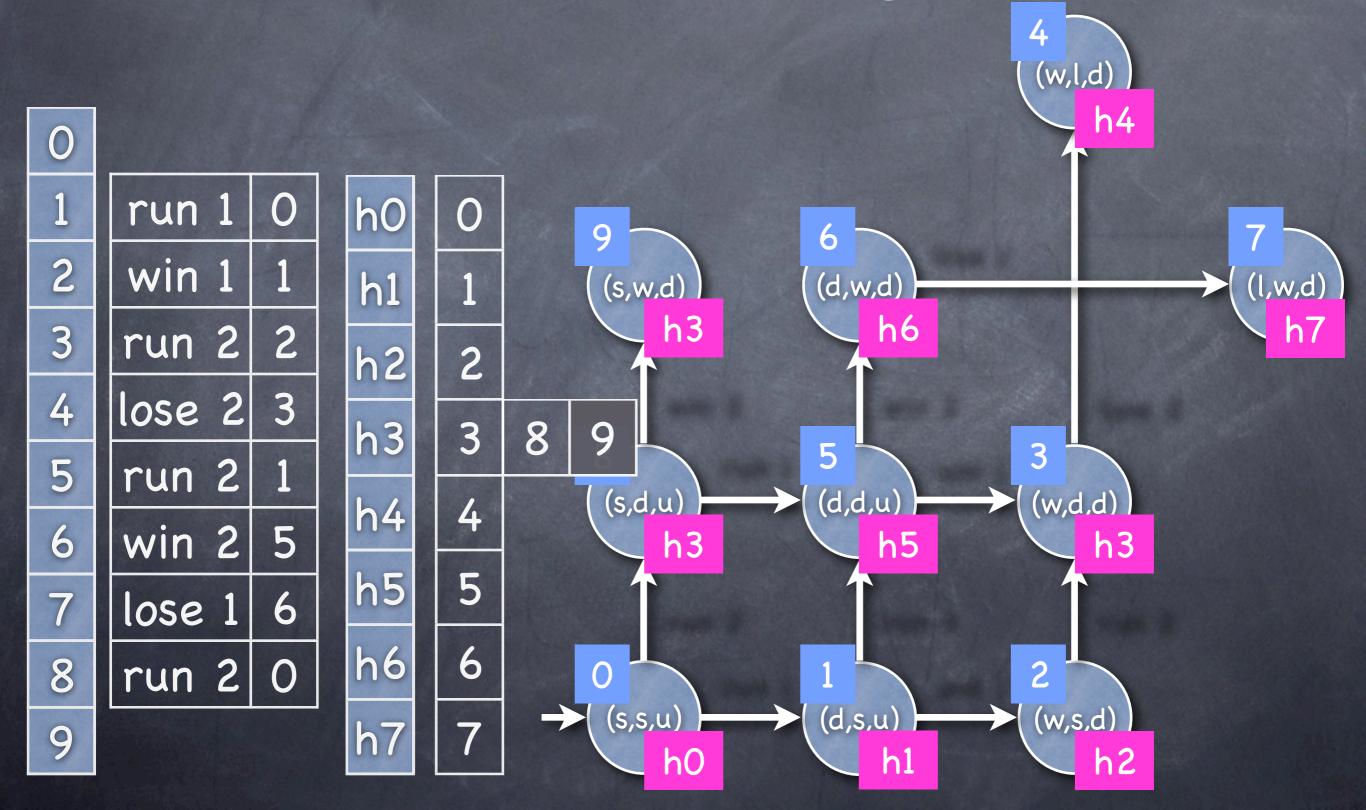


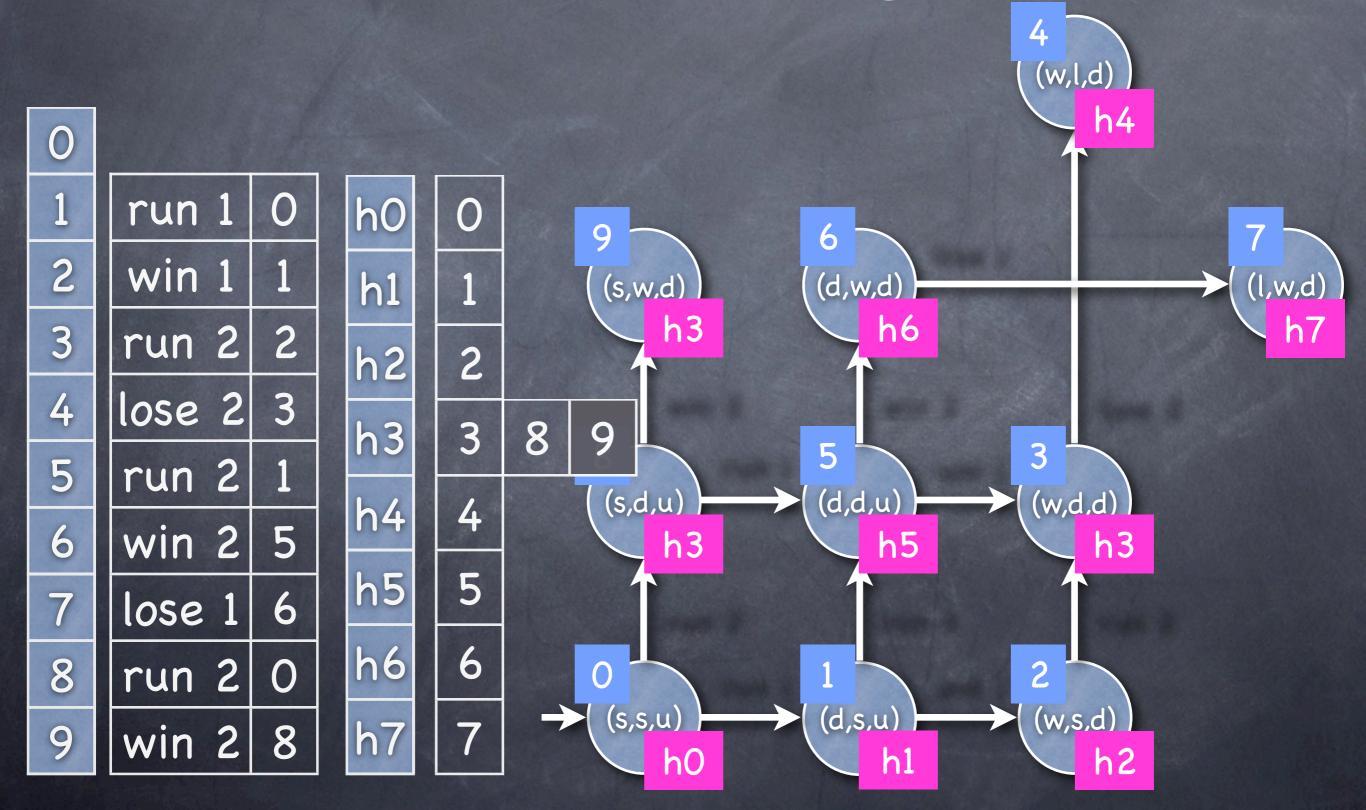


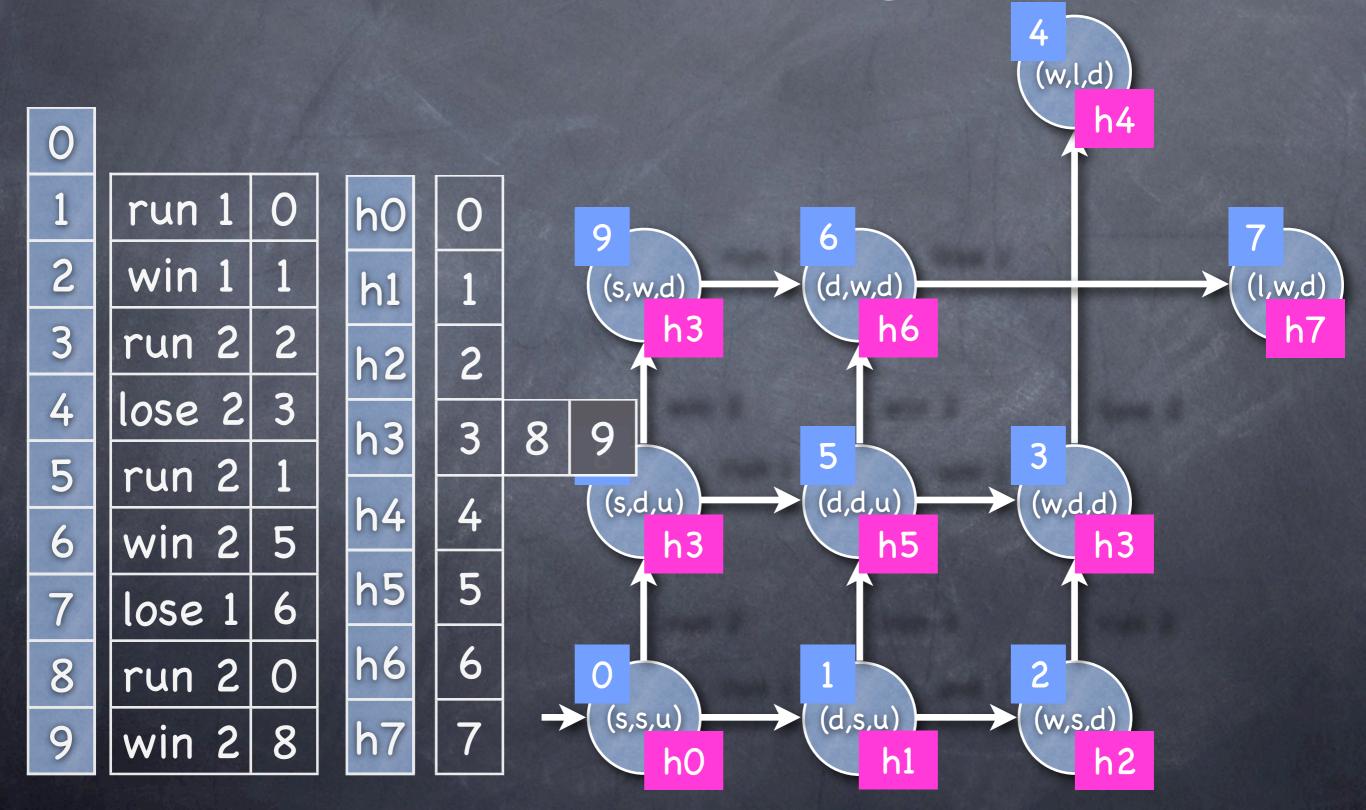












Evaluation of the Algorithm

- The algorithm uses $|R| \cdot (w + (2 \cdot w + h)) + 2 \cdot |R| \cdot (h + \log|T|) = |R| \cdot (3 \cdot (w + h) + \log|T|) <= 7 \cdot w \cdot |R|$
 - R: reachable states
 - w: size of machine word
 - h: size of hash-value
 - T: transitions
- Experiments show that 16·w·|R| is used
 - $oldsymbol{o}$ log|T| = 3·w
 - In SML-NJ everything is a reference

More Information about the Algorithm

L. Arge, G.S. Brodal, S. Christensen, L.M. Kristensen, and M. Westergaard: The ComBack State Space Exploration Method: Combining Hash Compaction with Backtracking, not yet submitted

Conclusion

- BRITNeY animation: formalism-independent tool for making formal models look good
- State space tool: formalism-independent platform for experimenting with multiple reduction methods of state spaces